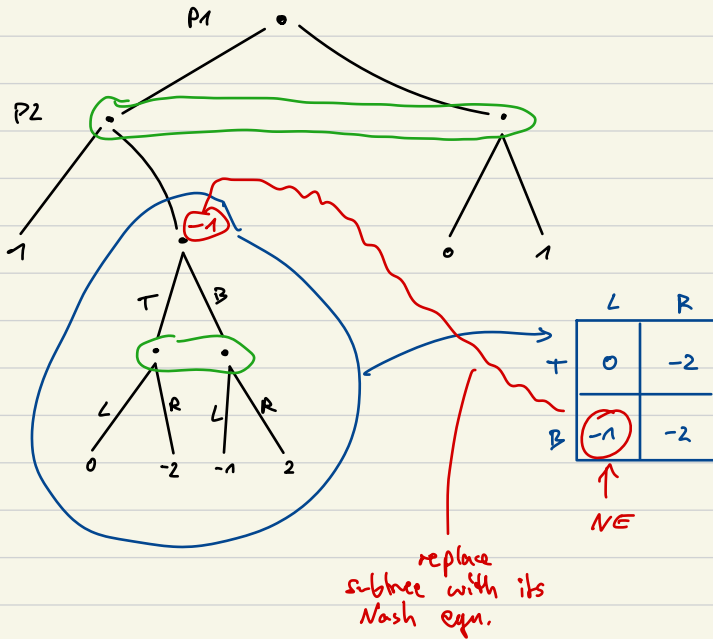


LECTURE 5

PART 1: EXTENSIVE FORM GAMES

Example for backward induction:

Consider the zero-sum game: (player 1: min, player 2: max)



We can conclude that the overall value of the game is 0.

Subgame Perfect Equilibria

Consider a feedback game (i.e. inf. sets don't span over stages) with K stages. Let I/J be the set of all inf. sets of player 1 / player 2.

Feedback game $\Rightarrow I = \bigcup_{i=1}^k I_i$ where I_i are inf. set at stage i

Define $g: I \rightarrow A$ with $g = \{g_1, \dots, g_k\}$
where $g_i: I_i \rightarrow A$

and similarly define $\sigma: J \rightarrow A$ for player 2.

A pair of strategies (g^*, σ^*) is a subgame perfect eq. if for all stages $t = 1, \dots, k$ and every g_1, g_2, \dots, g_{t-1} and $\sigma_1, \sigma_2, \dots, \sigma_{t-1}$ prior to stage t , for all g_t, σ_t , we have

$$\begin{aligned} J(g_{1:t-1}, g_t^*, g_{t+1:k}^*, \sigma_{1:t-1}, \sigma_t, \sigma_{t+1:k}^*) \\ &= J(g_{1:t-1}, g_t^*, g_{t+1:k}^*, \sigma_{1:t-1}, \sigma_t^*, \sigma_{t+1:k}^*) \\ &= J(g_{1:t-1}, g_t, g_{t+1:k}^*, \sigma_{1:t-1}, \sigma_t^*, \sigma_{t+1:k}^*) \end{aligned}$$

where we use the notation $g_{1:t-1} := (g_1, \dots, g_{t-1})$, meaning

$$\begin{aligned} J(g_{1:t-1}, g_t^*, g_{t+1:k}^*, \sigma_{1:t-1}, \sigma_t^*, \sigma_{t+1:k}^*) &:= \\ J(g_1, g_2, \dots, g_{t-1}, g_t^*, g_{t+1}^*, \dots, g_k^*, & \\ \sigma_1, \sigma_2, \dots, \sigma_{t-1}, \sigma_t^*, \sigma_{t+1}^*, \dots, \sigma_k^*) & \end{aligned}$$

Nash Equilibria

(g^*, σ^*) is a Nash Equ. if $\forall g_{1:k}, \sigma_{1:k}$ we have

$$J(g_{1:k}^*, \sigma_{1:k}^*) \leq J(g_{1:k}^*, \sigma_{1:k}) \leq J(g_{1:k}, \sigma_{1:k}^*)$$

Observe that:
(see Lemma 7.1
in Hespánha book)

subgame perfect equilibrium
 \Downarrow
Nash equilibrium

Proof idea:

For stage $t=1$ by definition of subgame perfect equilibrium,

$$\forall y_1, J(y_1^*, y_2^{*:K}, \sigma^*) \leq J(y_1, y_2^{*:K}, \sigma^*)$$

by applying def. for stage 2 \rightarrow

$$\begin{aligned} & \leq J(y_1, y_2, y_3^{*:K}, \sigma^*) \\ & \vdots \\ & \leq J(y_1^{*:K}, \sigma^*) \end{aligned}$$

PART 2: Review (for Quiz)

① General-sum games with finite action spaces

Ex 1 (see Thm. 10.1 in Hespánha book)

Consider a 2-pl. general-sum game with costs $A, B \in \mathbb{R}^{m \times n}$

To find mixed-strategy NE $y^* \in Y, z^* \in Z$ we can solve the quadratic optimization problem

quadratic
objective fct. $\left\{ \begin{array}{l} \min \\ y \in \mathbb{R}^m, z \in \mathbb{R}^n, \\ q, p \in \mathbb{R} \end{array} \right.$

linear
constraints $\left\{ \right.$

s.t.

$$y^T (A+B) z - p - q$$

$$Az \geq 1 \cdot p, \quad 1^T z = 1, \quad z \geq 0$$

$$B^T y \geq q \cdot 1, \quad 1^T y = 1, \quad y \geq 0$$

coordinate-wise

To get an idea for why this is hard to solve, we want to show that the objective function is non-convex.

Recall: For $f: \mathbb{R}^d \rightarrow \mathbb{R}$, $f(x) = x^T Q x$,

f is convex \iff all eigenvalues of $Q + Q^T$ are ≥ 0

How to show this?

\hookrightarrow let $x_1, x_2 \in \mathbb{R}^n$, $\lambda \in [0, 1]$; we need to show that

$$f(\lambda x_1 + (1-\lambda)x_2) \leq \lambda f(x_1) + (1-\lambda)f(x_2) \iff \begin{matrix} \text{all ev.'s of } Q+Q^T \\ \text{are } \geq 0 \end{matrix}$$

\hookrightarrow some algebra ...

$$\begin{aligned} & (\lambda x_1 + (1-\lambda)x_2)^T Q (\lambda x_1 + (1-\lambda)x_2) \\ &= \lambda x_1^T Q x_1 + \dots \\ & \vdots \end{aligned}$$

more calculations yield: $0 \leq (x_1 - x_2)^T \left(\frac{Q + Q^T}{2} \right) (x_1 - x_2) \quad \forall x_1, x_2$

$$\iff 0 \leq y^T \left(\frac{Q + Q^T}{2} \right) y, \quad \forall y$$

$$\iff \text{all ev.'s of } Q + Q^T \text{ are } \geq 0$$

In our case we want to show non-convexity:

$$f: \mathbb{R}^{n+m} \rightarrow \mathbb{R}$$

$$f(y, z) := y^T (A + B) z = \begin{bmatrix} y \\ z \end{bmatrix}^T \underbrace{\begin{bmatrix} 0 & \frac{A+B}{2} \\ \frac{(A+B)^T}{2} & 0 \end{bmatrix}}_{:= Q} \begin{bmatrix} y \\ z \end{bmatrix}$$

For any z , set $y = (A+B)z$. Then $z^T (A+B)^T (A+B) z \geq 0$.

But if we instead pick $y = -(A+B)z$, then $-z^T (A+B)^T (A+B) z \leq 0$.

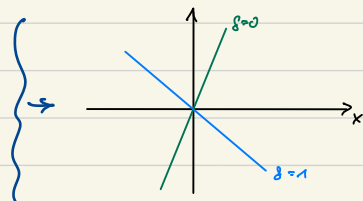
\Rightarrow eigenvalues of $Q+Q^T$ not all ≥ 0

\Rightarrow non-convexity.

② Review of 0-sum games

Ex 2: Consider $f: [-1, 1] \times \{0, 1\} \rightarrow \mathbb{R}$

$$f(x, \delta) = \begin{cases} 2x & \text{if } \delta = 0 \\ -x & \text{if } \delta = 1. \end{cases}$$

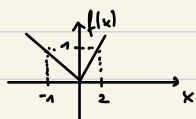


Does it hold that

$$\min_{x \in [-1, 1]} \max_{\delta \in \{0, 1\}} f(x, \delta) = \max_{\delta \in \{0, 1\}} \min_{x \in [-1, 1]} f(x, \delta)$$

(Recall: We know that LHS \geq RHS from inf-sup theorem.)

LHS: $\max_{\delta \in \{0, 1\}} f(x, \delta)$ looks like



$$\text{RHS: } \min_{x \in [-1, 1]} f(x, \delta) = \begin{cases} -2 & \text{if } \delta = 0 \\ 1 & \text{if } \delta = 1 \end{cases}$$

We conclude that LHS \neq RHS!

(recall that equality is only guaranteed to hold in certain cases, e.g. convex-concave objective)

③ Potential games

Ex 3

Let $J_i: T_1 \times T_2 \times \dots \times T_N \rightarrow \mathbb{R}$. Suppose the game with payoffs $\{J_i\}_{i=1}^N$ admits a potential function.

TRUE or FALSE?

- 1) $p^* \in T$ is a NE $\Rightarrow p^*$ is a maximizer of the potential function
 \rightarrow FALSE (e.g. stag hunt: only one of NEs maximizes potential function)
- 2) $p^* \in T$ is maximizer of potential $\Rightarrow p^*$ is NE
 \rightarrow TRUE (proven in lecture)
- 3) best reply dynamics converges \Rightarrow game is potential
 \rightarrow FALSE (we proved " \Leftarrow " direction; as counterexample consider
$$\begin{bmatrix} (1,0) & \textcircled{(2,0)} \\ (2,0) & (0,1) \end{bmatrix} \text{NE}$$
 (players are maximizers)
 \downarrow
best reply converges to NE but game is not potential \rightarrow verify this!)