


Mock Exam 2025		
	<b>Course:</b>	Numerical Analysis – EL/MX/CH
	<b>Lecturer:</b>	Prof. Michael Herbst
	<b>Date:</b>	Spring Session 2025 <b>Duration:</b> 3h 00

## Mock exam disclaimer

- This mock exam gives you an example how the final exam will look like. Questions 1–4, 7, & 8 have been taken from the **2024 exam**, question 5 & 6 from the **2024 mock exam**. As the course has changed compared to last year, some questions may use terminology we did not employ this year.
- In 2025, the final exam will only contain **pen and paper** questions, which moreover will cover the contents of **the entire class**. Furthermore you can expect this year's **final exam** to be a little **more involved** than this mock exam.

## Exam instructions

This exam has **8 exercises** with in total **45 points**. The material consists of **two** parts:

- This **question sheet** with the questions. All questions are pen and paper questions.
- A personalised **answer sheet** (with your name and sciper number).

### Answering the questions

- For each question on this question sheet, provide the **answers in the corresponding section** of the **answer sheet**. **Do not write onto the question sheet itself**. Only these answer sheets will be marked.
- On the **answer sheet** only **write within the black boxes**. If you need extra space, additional blank pages are given in the back. **Clearly identify for which question** you provide additional answers. **Also add a remark in the original answer box** where you run out of space that additional text can be found in the appendix.
- **Please write with a pen (no pencil or erasable ballpen)**.
- For each of your answers, **outline the reasoning** and **justify your answer**.
- At the end of the exam both the question and answer sheet will be collected.

## Authorised material

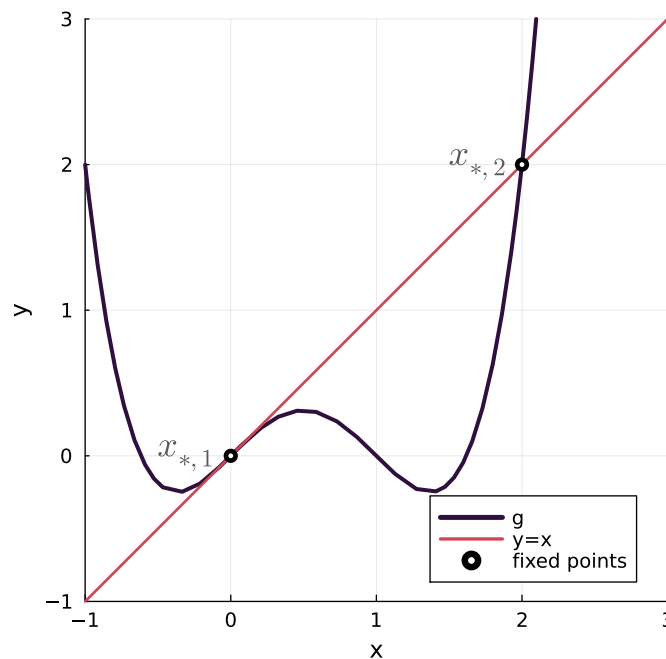
- You are allowed a **two-sided handwritten A4 cheatsheet** (hand-written on paper, no print-outs).
- No other notes, sheets or books are allowed. No calculator, mobile phone, tablet, laptop or any other electronic device.

## Exercise 1 (7P) — exam 2024

Given a parameter  $\theta \in \mathbb{R}$ , consider the function

$$g(x) = x^4 - 2\theta x^3 + x.$$

- (a) **(2P)** Show that the only fixed points of  $g$  are  $x_{*,1} = 0$  and  $x_{*,2} = 2\theta$ .
- (b) **(1P)** Figure 1 depicts  $g(x)$  for  $\theta = 1$  together with its two fixed points  $x_{*,1}$ ,  $x_{*,2}$ . By visual inspection determine for which of the two fixed points  $x_{*,1}$  and  $x_{*,2}$  the fixed-point iterations  $x^{(k+1)} = g(x^{(k)})$  converge, provided that a starting point  $x^{(0)}$  sufficiently close to the respective fixed point has been chosen. Justify your answer.



- (c) **(3P)** We return to the general case where  $\theta$  is a parameter of the problem.
- For which values of  $\theta$  do fixed-point iterations converge to  $x_{*,2}$  provided a good starting point is chosen?
  - For which values of  $\theta$  is the fastest convergence rate to  $x_{*,2}$  achieved?
- (d) We consider the case where  $\theta$  is chosen such that the fastest convergence rate in the fixed-point iterations is achieved (the value you determined in (c) (ii)).
- (1P)** What is the convergence order of Newton's method for these value(s) of  $\theta$ ? Does Newton provide any advantage over fixed point iterations in this case?

## Exercise 2 (8P) — exam 2024

- (a) **(1.5P)** We are given  $n + 1$  nodes  $x_1, \dots, x_{n+1} \in \mathbb{R}$ . Define the Lagrange basis associated to  $x_1, \dots, x_{n+1}$  and specify how this basis can be used to find the  $n$ -th degree interpolating polynomial through the data points  $(x_1, y_1), (x_2, y_2), \dots, (x_{n+1}, y_{n+1})$ .
- (b) **(1.5P)** Using the Lagrange basis, find the interpolating polynomial through the points  $(x_1 = -1, y_1 = -2)$ ,  $(x_2 = 1, y_2 = 0)$ ,  $(x_3 = 4, y_3 = 6)$ .

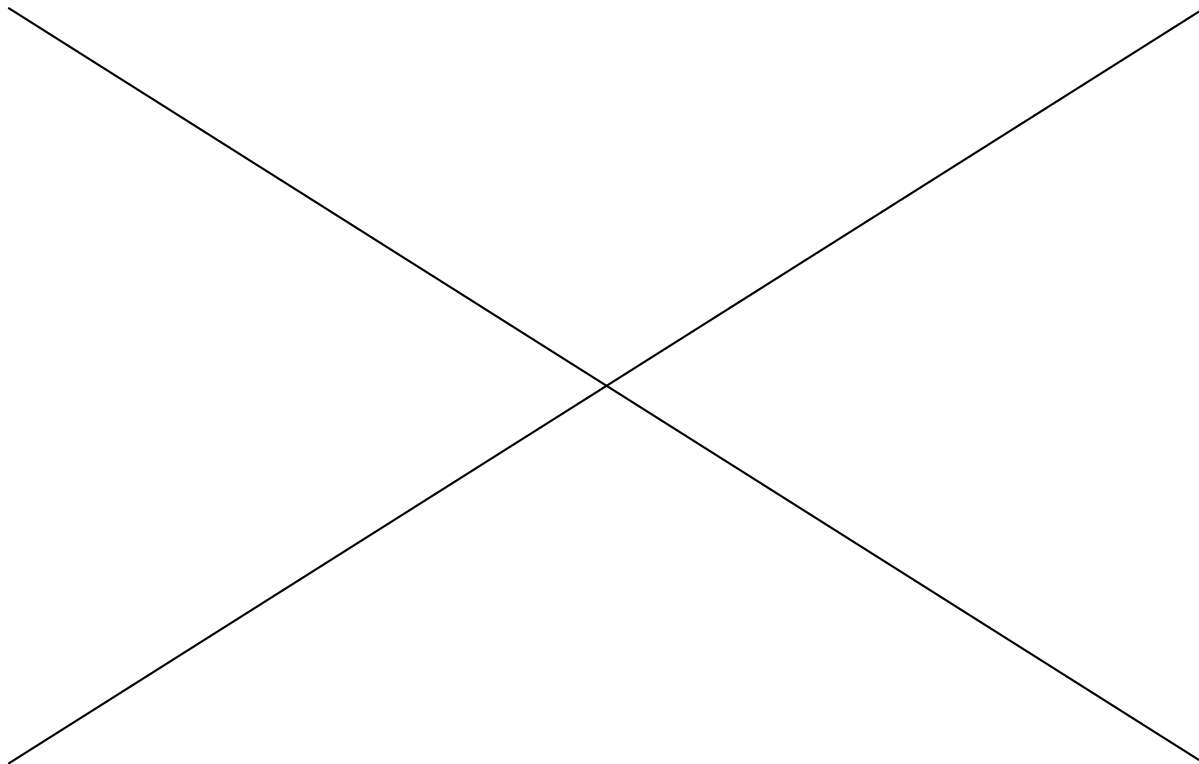
For polynomial interpolation with equally spaced nodes, the interpolation error is governed by the following theorem:

**Theorem.** For a  $C^{n+1}$  function  $f : [a, b] \rightarrow \mathbb{R}$  and  $a = x_1 < x_2 < \dots < x_{n+1} = b$  equally distributed nodes in  $[a, b]$  the  $n$ -th degree polynomial interpolant  $p_n$  of the data  $(x_i, f(x_i))$  with  $i = 1, 2, \dots, n + 1$  satisfies the following bound on the interpolation error

$$\max_{x \in [a, b]} |f(x) - p_n(x)| \leq \frac{1}{4(n+1)} \left( \frac{b-a}{n} \right)^{n+1} \max_{x \in [a, b]} |f^{(n+1)}(x)|. \quad (1)$$

We consider polynomial interpolation with  $n + 1$  equally distributed nodes over the interval  $[-1, 1]$  for the functions  $f_1(x) = \sin(x)$  and  $f_2(x) = \frac{1}{1+20x^2}$ .

- (c) **(3P)** Show that for  $f_1$  the interpolation error goes to 0 as  $n \rightarrow \infty$ .
- (d) **(2P)** For  $f_2$  we have  $\max_{x \in [0, 1]} |f_2^{(n+1)}(x)| \approx 20^n n!$  such that the right-hand side of (1) grows to infinity as  $n \rightarrow \infty$ . What happens to the polynomial interpolation in this case? What needs to be changed in the polynomial interpolation procedure to achieve exponential convergence for such functions  $f_2$ ?



**Exercise 3 (7P) — exam 2024**

We consider the matrix

$$\mathbf{A} = \begin{pmatrix} 1 & \alpha \\ \alpha & 1 \end{pmatrix}$$

for  $\alpha > 0$  and an associated linear system

$$\mathbf{Ax} = \mathbf{b} \quad \text{with} \quad \mathbf{b} = \begin{pmatrix} 1 \\ 1 \end{pmatrix} \quad (2)$$

where we seek the solution  $\mathbf{x} \in \mathbb{R}^2$ .

When representing (2) on a computer we assume that the available floating-point precision is unable to represent  $\mathbf{b}$  exactly introducing a small error  $\varepsilon > 0$ : the computer is only able to solve the approximate linear system

$$\mathbf{A}\tilde{\mathbf{x}} = \tilde{\mathbf{b}} \quad \text{with} \quad \tilde{\mathbf{b}} = \begin{pmatrix} 1 + \varepsilon \\ 1 - \varepsilon \end{pmatrix} \quad (3)$$

and thus only able to obtain an approximate solution  $\tilde{\mathbf{x}} \in \mathbb{R}^2$ .

- (a) **(2P)** For a *general* square and invertible matrix  $\mathbf{M} \in \mathbb{R}^{n \times n}$  define the condition number  $\kappa(\mathbf{M})$  in terms of matrix norms. Also provide an expression to compute the condition number using eigenvalues of appropriate matrices.
- (b) **(2P)** Show that the condition number of  $\mathbf{A}$  is

$$\kappa(\mathbf{A}) = \left| \frac{1 + \alpha}{1 - \alpha} \right|. \quad (4)$$

You may use that a matrix  $\begin{pmatrix} a & c \\ c & b \end{pmatrix} \in \mathbb{R}^{2 \times 2}$  has eigenvalues

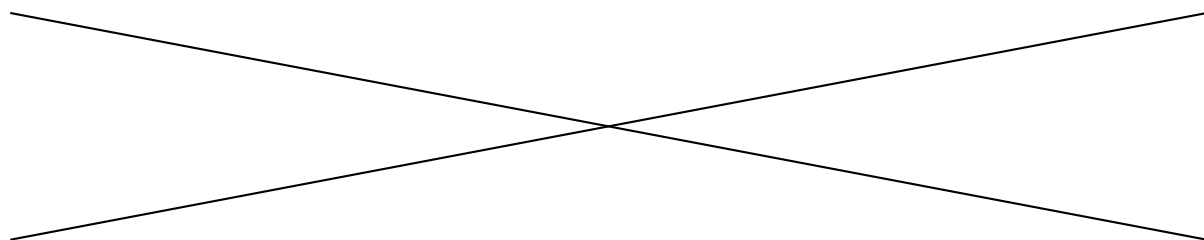
$$\frac{a+b}{2} - \sqrt{\frac{(a-b)^2}{4} + c^2} \quad \text{and} \quad \frac{a+b}{2} + \sqrt{\frac{(a-b)^2}{4} + c^2}.$$

- (c) **(1P)** The solution to the perturbed system (3) is given by

$$\tilde{\mathbf{x}} = \frac{1}{1+\alpha} \begin{pmatrix} 1 \\ 1 \end{pmatrix} + \frac{\varepsilon}{1-\alpha} \begin{pmatrix} 1 \\ -1 \end{pmatrix}$$

and similarly  $\mathbf{x} = \frac{1}{1+\alpha} \begin{pmatrix} 1 \\ 1 \end{pmatrix}$ . Compute the relative error in the solution  $\frac{\|\tilde{\mathbf{x}} - \mathbf{x}\|}{\|\mathbf{x}\|}$  as well as the relative error in the right-hand side  $\frac{\|\tilde{\mathbf{b}} - \mathbf{b}\|}{\|\mathbf{b}\|}$ .

- (d) **(2P)** Describe in one sentence what the condition number measures for a linear system (2). Use this to explain your results in (c).



## Exercise 4 (5P) — exam 2024

Consider the algorithm for **LU factorization** given below.

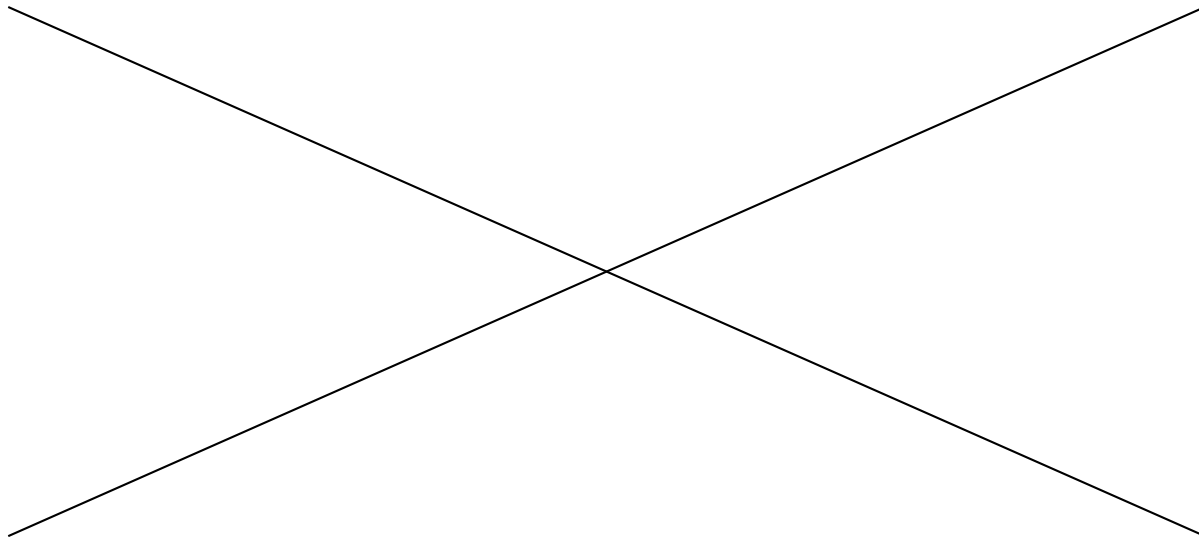
**Algorithm (LU factorisation).**

**Input:**  $\mathbf{A} \in \mathbb{R}^{n \times n}$ ,

**Output:**  $\mathbf{U} \in \mathbb{R}^{n \times n}$ ,  $\mathbf{L} \in \mathbb{R}^{n \times n}$

1.  $\mathbf{A}^{(1)} = \mathbf{A}$
2. for  $k = 1, \dots, n - 1$  (*algorithm steps*)
  1.  $L_{kk} = 1$
  2. for  $i = k + 1, \dots, n$  (*Loop over rows*)
    1.  $L_{ik} = \frac{A_{ik}^{(k)}}{A_{kk}^{(k)}}$
    2. for  $j = k + 1, \dots, n$  (*Loop over columns*)
      1.  $A_{ij}^{(k+1)} = A_{ij}^{(k)} - L_{ik}A_{kj}^{(k)}$
3.  $\mathbf{U} = \mathbf{A}^{(n)}$

- (a) **(1.5P)** Making reference to the algorithm explain why LU factorisation is said to have a computational cost of  $\mathcal{O}(n^3)$ .
- (b) **(1.5P)** Let  $\mathbf{A} \in \mathbb{R}^{n \times n}$  be an  $n$  by  $n$  square matrix. If  $\mathbf{A}$  has no special structure what is the memory usage for storing  $\mathbf{A}$ ? What is the memory usage for storing the  $\mathbf{L}$  and  $\mathbf{U}$  factors once LU factorisation has been performed? Specify your answer using big  $\mathcal{O}$  notation.
- (c) Now assume that  $\mathbf{A}$  is sparse.
  - **(1P)** If we only store the non-zero elements of the matrix explicitly, what is the memory cost of  $\mathbf{A}$  in this case?
  - **(1P)** Explain the phenomenon called fill-in and its consequences for the memory cost of the  $\mathbf{L}$  and  $\mathbf{U}$  factors of sparse matrices.



**Exercise 5 (7P) — mock 2024**

Let  $\mathbf{Ax} = \mathbf{b}$  be a linear system with given  $\mathbf{A} \in \mathbb{R}^{n \times n}$  and  $\mathbf{b} \in \mathbb{R}^n$ . We consider the fixed-point map

$$g(\mathbf{x}) = \mathbf{x} + \mathbf{P}^{-1}(\mathbf{b} - \mathbf{Ax}) \quad (5)$$

for an invertible preconditioner matrix  $\mathbf{P} \in \mathbb{R}^{n \times n}$ . The iterative procedure  $\mathbf{x}^{(k+1)} = g(\mathbf{x}^{(k)})$  starting from an initial vector  $\mathbf{x}^{(0)} \in \mathbb{R}^n$  is the Richardson iteration.

(a) **(1P)** Show that if  $\mathbf{x} \in \mathbb{R}^n$  is a fixed point of  $g$  then it is also a solution to the linear system  $\mathbf{Ax} = \mathbf{b}$ .

(b) **(2P)** Show that if  $\mathbf{x}^{(k+1)} = g(\mathbf{x}^{(k)})$  and if  $\mathbf{x}$  is a fixed point of  $g$ , then

$$(\mathbf{x}^{(k+1)} - \mathbf{x}) = (\mathbf{I} - \mathbf{P}^{-1}\mathbf{A})(\mathbf{x}^{(k)} - \mathbf{x}). \quad (6)$$

(c) **(1.5P)** Based on (6) give the conditions for Richardson iterations  $\mathbf{x}^{(k+1)} = g(\mathbf{x}^{(k)})$  to converge to a fixed point  $\mathbf{x}$  independent of the chosen initial vector  $\mathbf{x}^{(0)}$  and right-hand side  $\mathbf{b}$ .

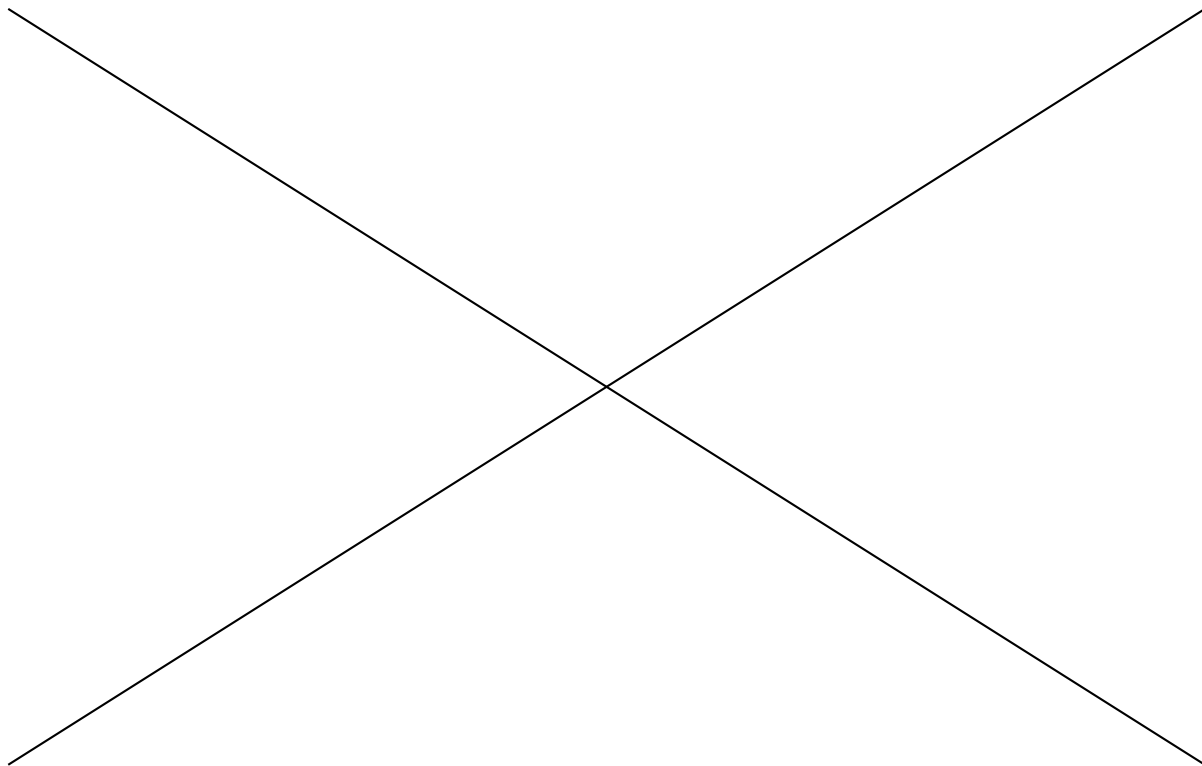
Consider the case

$$\mathbf{A} = \begin{pmatrix} 2 & 0 \\ 0 & \alpha \end{pmatrix} \quad \mathbf{P} = \begin{pmatrix} \beta & 0 \\ 0 & 1 \end{pmatrix} \quad \mathbf{b} = \begin{pmatrix} 1 \\ 1 \end{pmatrix} \quad \mathbf{x}^{(0)} = \begin{pmatrix} 0 \\ 0 \end{pmatrix}$$

where  $\alpha \in \mathbb{R}$  and  $\beta \in \mathbb{R}$  and consider the fixed-point iterations  $\mathbf{x}^{(k+1)} = g(\mathbf{x}^{(k)})$  for  $k = 0, 1, \dots$

(d) **(1.5P)** For which choice of  $\alpha$  and  $\beta$  do the fixed-point iterations converge.

(e) **(1P)** For which choice of  $\alpha$  and  $\beta$  do the fixed-point iterations converge at the fastest possible rate. How many iteration steps are at most needed for convergence?



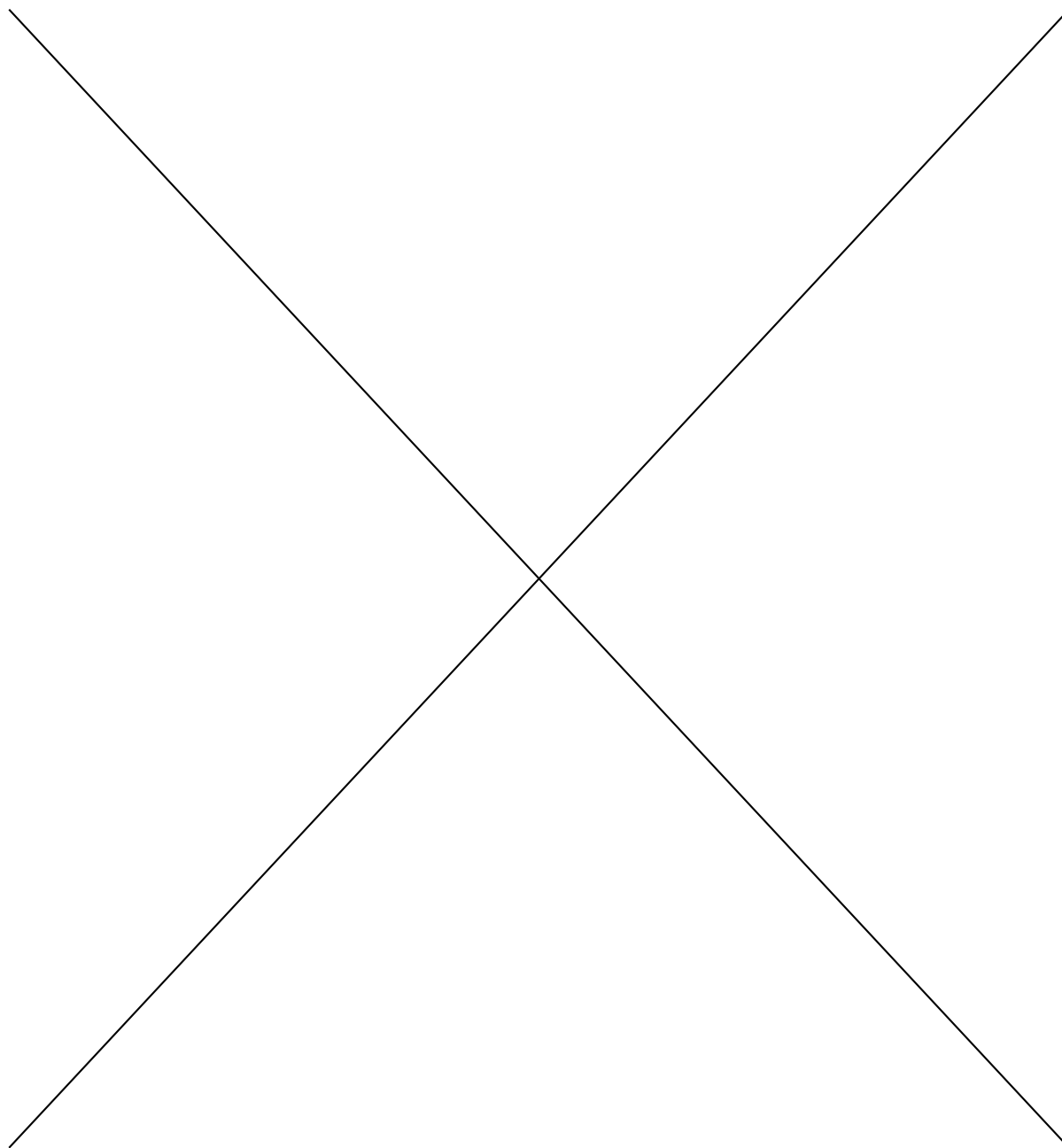
**Exercise 6 (5P) — mock 2024**

Consider  $\mathbf{b} \in \mathbb{R}^2$  and the following  $2 \times 2$  matrix

$$\mathbf{A} = \begin{pmatrix} 1 & -\beta \\ -\frac{1}{2} & 1 \end{pmatrix},$$

for a real number  $\beta \in \mathbb{R}$ . Consider Jacobi's method to solve the linear system  $\mathbf{Ax} = \mathbf{b}$ .

- (a) **(2.5P)** State the conditions for Jacobi's method to converge for any right-hand side  $\mathbf{b}$  and initial vector  $\mathbf{x}^{(0)}$ .
- (b) **(2.5P)** Deduce conditions for the parameter  $\beta$ , which ensure convergence for any right-hand side  $\mathbf{b}$  and initial vector  $\mathbf{x}^{(0)}$ .



## Exercise 7 (6P) — exam 2024

We consider the family of triangular matrices

$$\mathbf{A} = \begin{pmatrix} \lambda_1 & 1 & 1 \\ 0 & \lambda_2 & 1 \\ 0 & 0 & \lambda_3 \end{pmatrix}$$

where  $\lambda_1 \geq \lambda_2 \geq \lambda_3 \geq 0$ .

- (a) **(3P)** We consider the power iterations on the matrix  $\mathbf{A}$  starting from an initial guess  $\mathbf{x}^{(0)} = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}$ . State conditions on the parameters  $\lambda_1$ ,  $\lambda_2$  and  $\lambda_3$  for the power iterations to converge. To which eigenvalue will they converge? Specify the convergence order and provide the convergence rate in terms of  $\lambda_1$ ,  $\lambda_2$  and  $\lambda_3$ .

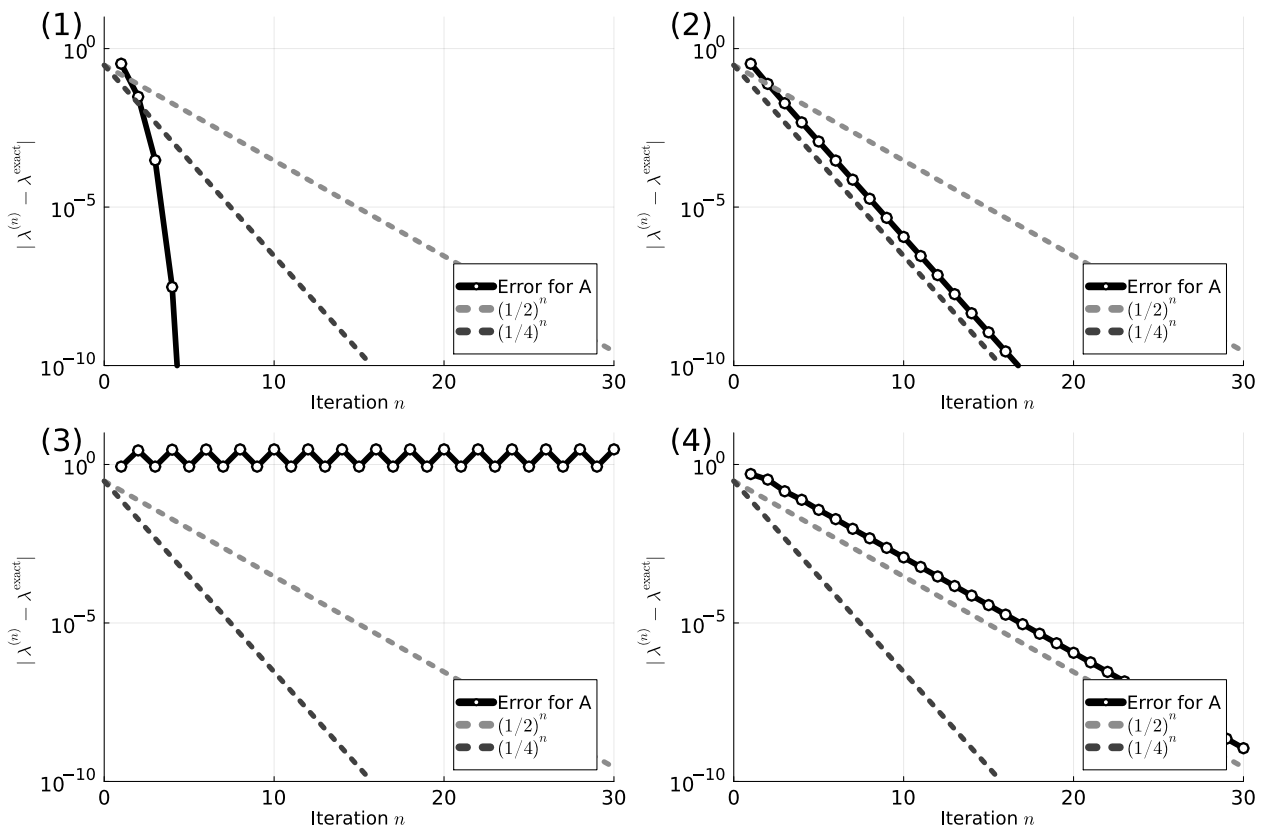
- (b) **(1P)** Prove for a general matrix  $\mathbf{M} \in \mathbb{R}^{n \times n}$ :

If  $(\alpha, \mathbf{x})$  is an eigenpair of  $\mathbf{M}$ , i.e.  $\mathbf{M}\mathbf{x} = \alpha\mathbf{x}$ , and  $\mathbf{M}$  is invertible, then  $(\frac{1}{\alpha}, \mathbf{x})$  is an eigenpair of  $\mathbf{M}^{-1}$ .

- (c) **(2P)** Consider the matrix

$$\mathbf{B} = \begin{pmatrix} 5 & 1 & 1 \\ 0 & 2 & 1 \\ 0 & 0 & 1 \end{pmatrix},$$

which is a special case of  $\mathbf{A}$  for  $\lambda_1 = 5$ ,  $\lambda_2 = 2$ ,  $\lambda_3 = 1$ . We perform *shifted inverse iterations* with shift  $\sigma = 4$ . Which eigenvalue  $\lambda_{\text{exact}}$  is targeted? Which of the following four plots is obtained? Justify your choice making reference to the discussion in the previous parts of the exercise.



## Exercise 8 (4P) — exam 2024

Let  $f : [a, b] \rightarrow \mathbb{R}$  be a real-valued function with  $0 < a < b$ . We consider a numerical integration formula

$$Q(f) = h \sum_{i=0}^n w_i f(t_i)$$

with  $n + 1$  equispaced quadrature nodes

$$t_i = a + ih \quad \text{for } i = 0, \dots, n \quad \text{and} \quad h = \frac{b - a}{n}$$

as well as weights  $w_i$  for  $i = 0, \dots, n$ .  $Q(f)$  approximates  $\int_a^b f(x) dx$ .

- (0.5P)** Define the **degree of exactness** of  $Q$ .
- (1P)** State the trapezoid formula for computing  $\int_a^b f(x) dx$  and provide its degree of exactness.
- (2.5P)** In the lecture we discussed

**Theorem.** If a numerical integration formula  $Q$  has a degree of exactness  $r$  then the formula is of order  $r + 1$ , i.e.

$$\left| \int_a^b f(x) dx - Q(f) \right| \leq C h^{r+1}$$

where  $C$  is a constant independent of  $h$ .

Inspect the following convergence graphs and apply this theorem to determine the degree of exactness of the two quadrature formulae (**Formula A** and **Formula B**). Which of the two formulae behaves like the trapezoid method?

