

# Week 5

# Industrial

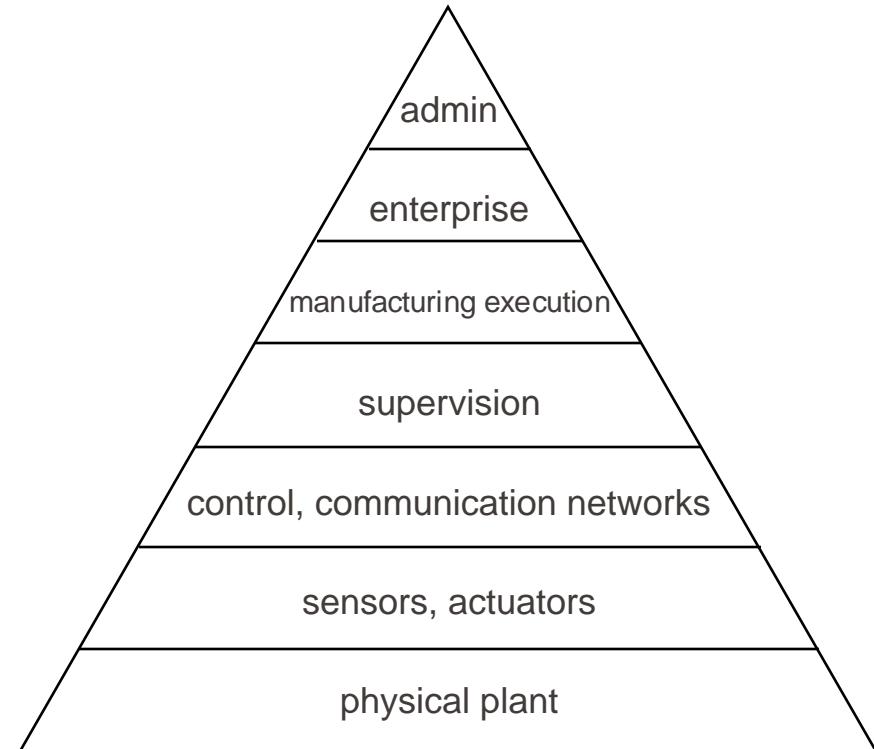
# Communication

# Networks

Dr. Philipp Sommer

# Industrial Communication Networks

- OT vs IT
- Field bus principles
- Field bus operation
- Standard field busses



## Operational Technology (OT):

- Factory floor
- Operation of physical processes and machines
- Examples: PLCs, SCADA, DCS
- Highly specialized equipment and software/protocols
- Stand-alone factory network

## Information Technology (IT):

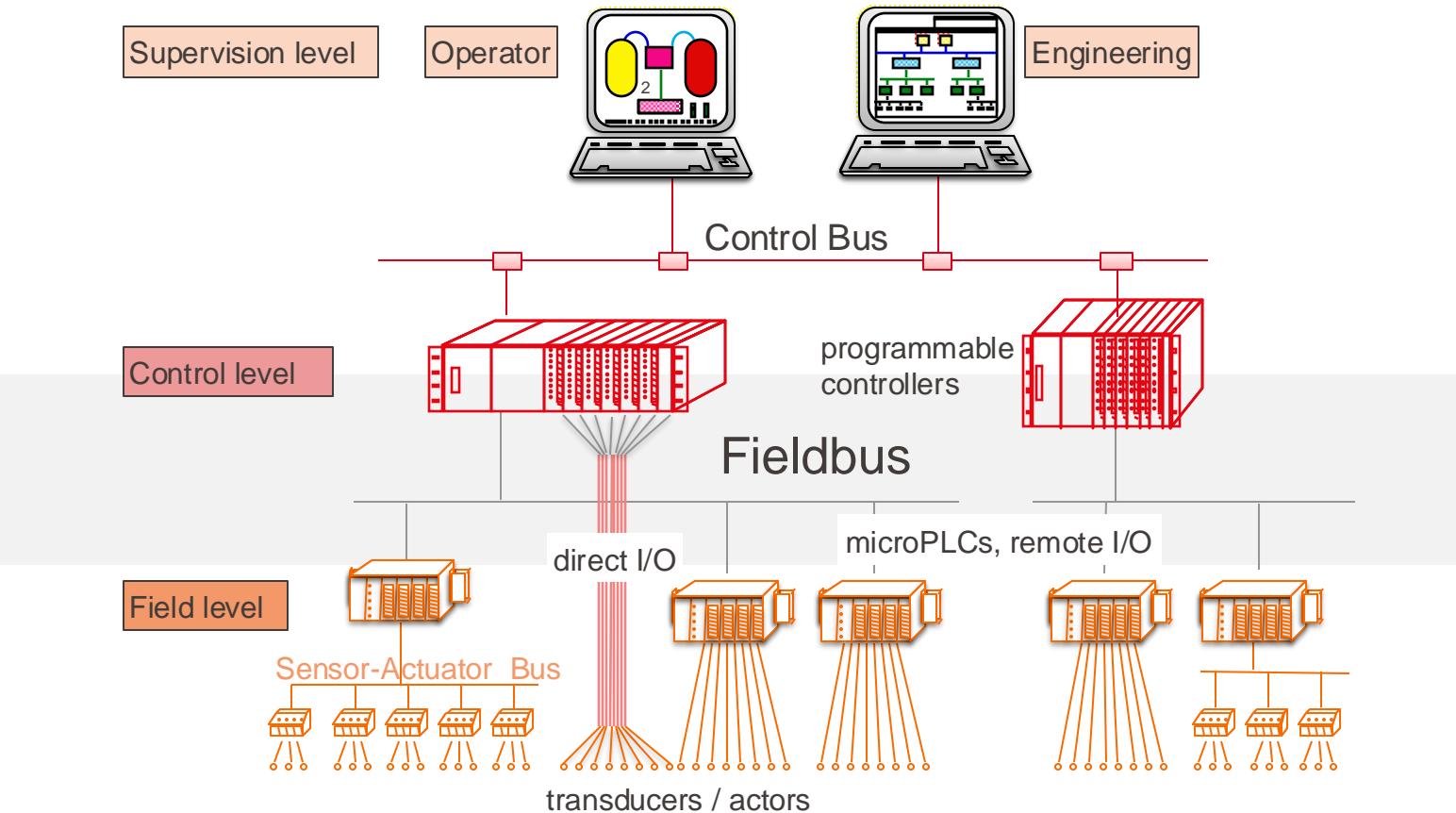
- Office environment
- Digital information (data)
- Examples: PCs, laptops, servers
- Highly standardized software and hardware components
- Connected to the cloud

IT-OT Convergence: IT and OT systems will grow closer together.

- Family of wired computer networking technologies standardized by the IEEE from 1983 onwards
- Widely used for Local-Area Networking (LAN) in the IT domain
- Various physical layer variants (10/100/1000 Mbit/s, ...)
- Gigabit speed over twisted-pair cables (CAT5 or better), 100 meters
- Can supply power to devices: Power-over-Ethernet (PoE)



# Networks in Automation Hierarchy

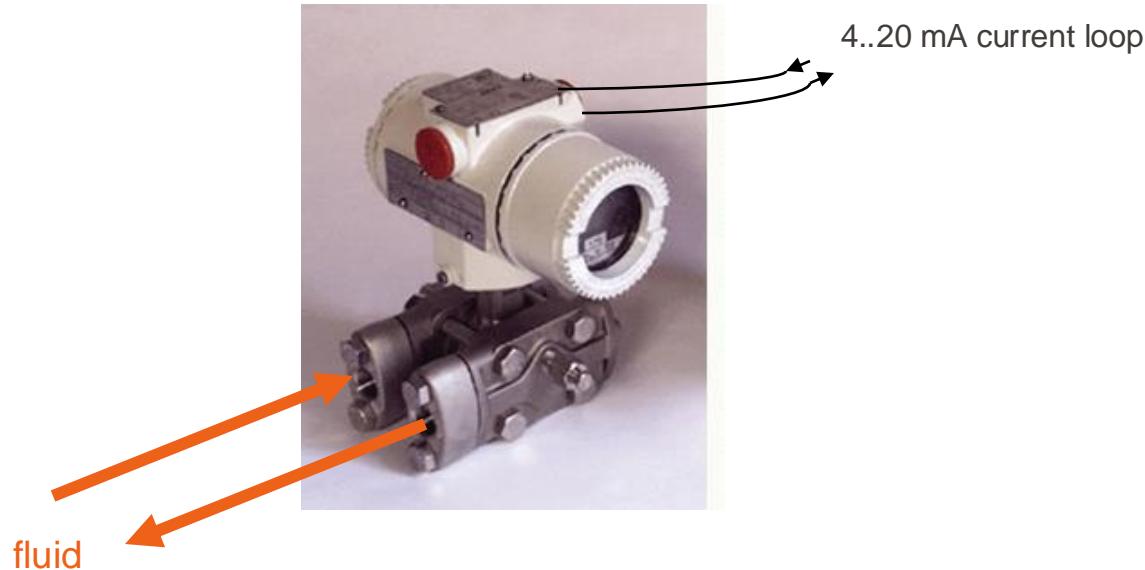


- Traditionally, Ethernet is used for the communication among the PLCs and for communication of the PLCs with the supervisory level and with the engineering tools.
- Fieldbus is in charge of the connection with the decentralized I/O and for time-critical communication among the PLCs.



# Connecting field devices

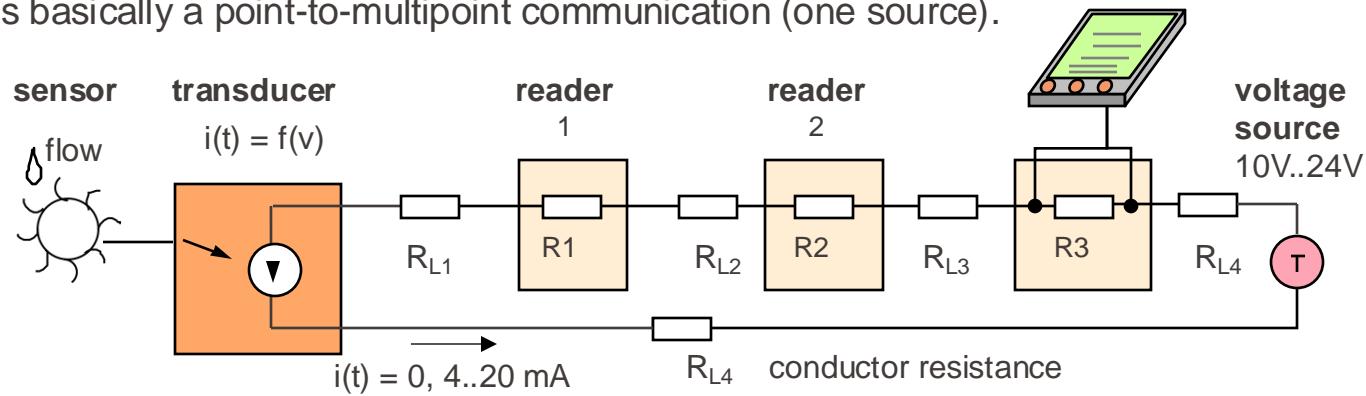
# Field device: example differential pressure transducer



The device transmits its value by means of a current loop.

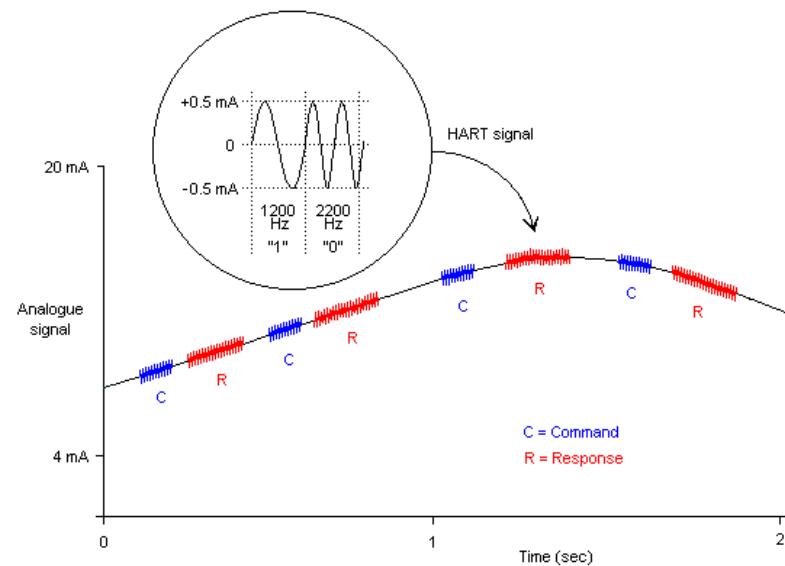
# 4-20 mA loop - the conventional, analog standard

- The 4-20 mA is the most common analog transmission standard in industry.
- The transducer limits the current to a value between 4 mA and 20 mA, proportional to the measured value, while 0 mA signals an error (wire break).
- The voltage drop along the cable and the number of readers induces no error.
- Simple devices are powered directly by the residual current (4 mA), allowing to transmit signal and power through a single pair of wires.
- 4-20mA is basically a point-to-multipoint communication (one source).

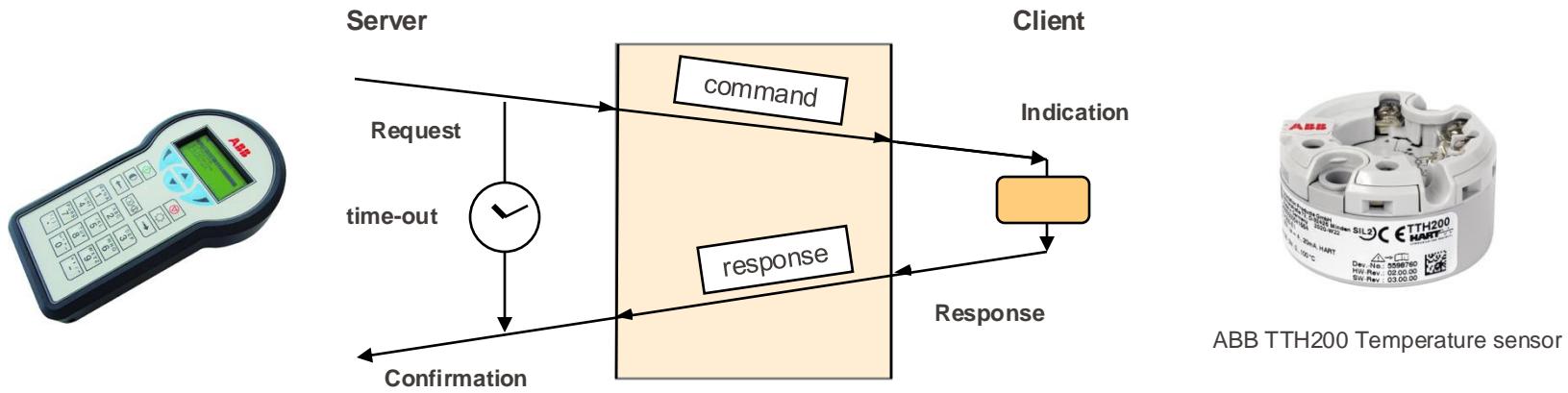


# HART – Principle (1986)

- HART (Highway Addressable Remote Transducer) was developed to retrofit 4-to-20mA current loop transducers with digital data communication
  - HART modulates the 4-20mA current with a low-level frequency-shift-keyed (FSK) sine-wave signal, without affecting the average analogue signal.
  - HART uses low frequencies (1200Hz and 2200 Hz) to deal with poor cabling, slow - but sufficient for its purpose.
  - Transmission of device characteristics is normally not real-time critical



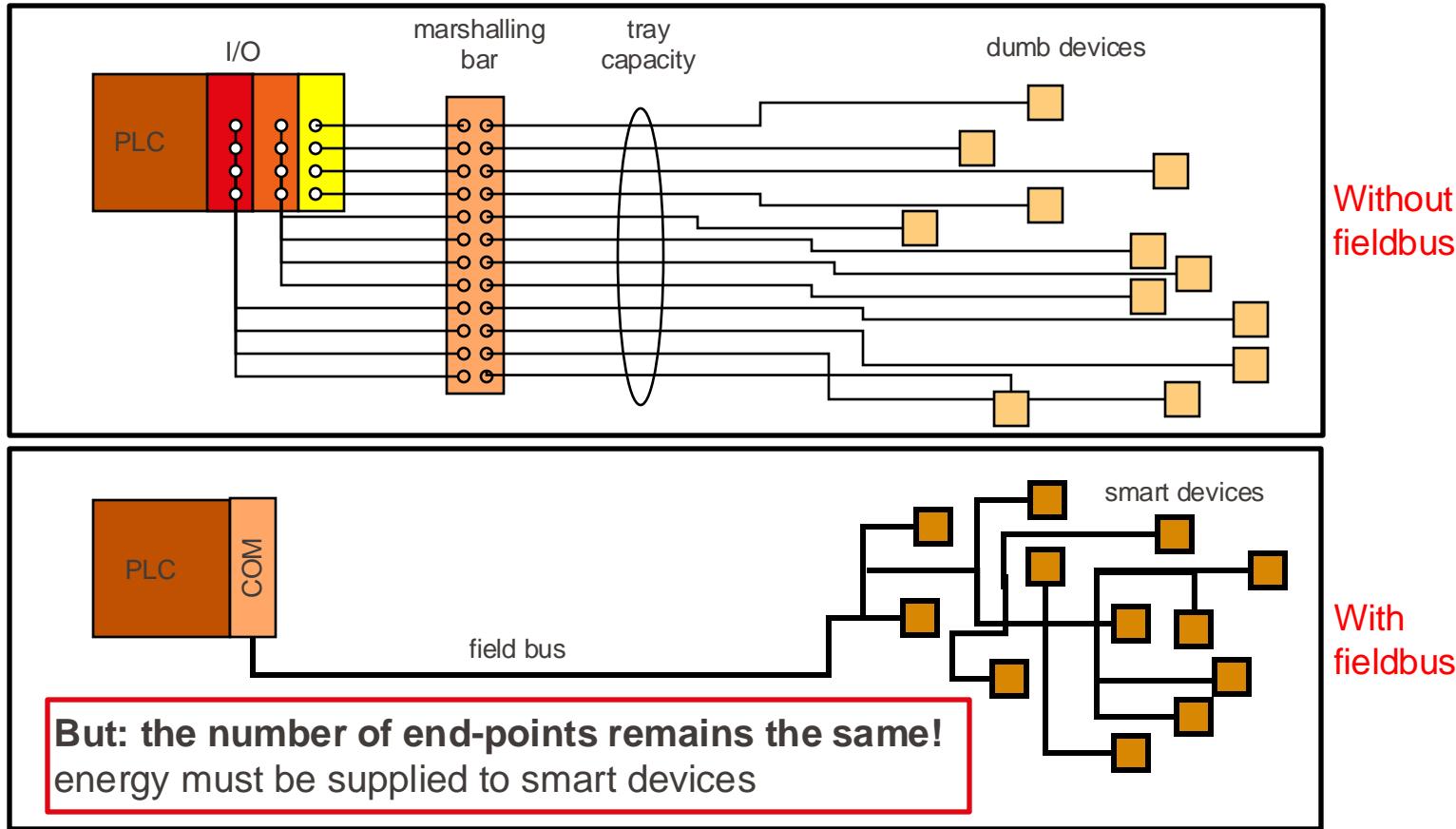
- Practically all 4..20 mA devices come equipped with HART today.
- HART communicates point-to-point, under the control of a server, e.g., a hand-held device.



- Universal commands (mandatory):
  - Identification
  - primary measured variable and unit (floating point format)
  - loop current value (%) = same info as current loop
  - read current and up to four predefined process variables
  - sensor serial number
  - instrument manufacturer, model, tag, serial number, descriptor, range limits, ...
- Total 44 standard commands, plus user-defined commands
- Transducer-specific (vendor-defined):
  - calibration data, trimming, ...

# Fieldbus

# The original fieldbus idea: save wiring



# Marshalling Bar

- The marshalling bar is the interface between the PLC and the instrumentation (sensor/actuators).
- The fieldbus replaces the marshalling bar or rather moves it piecewise to the process
- intelligent concentrator / wiring



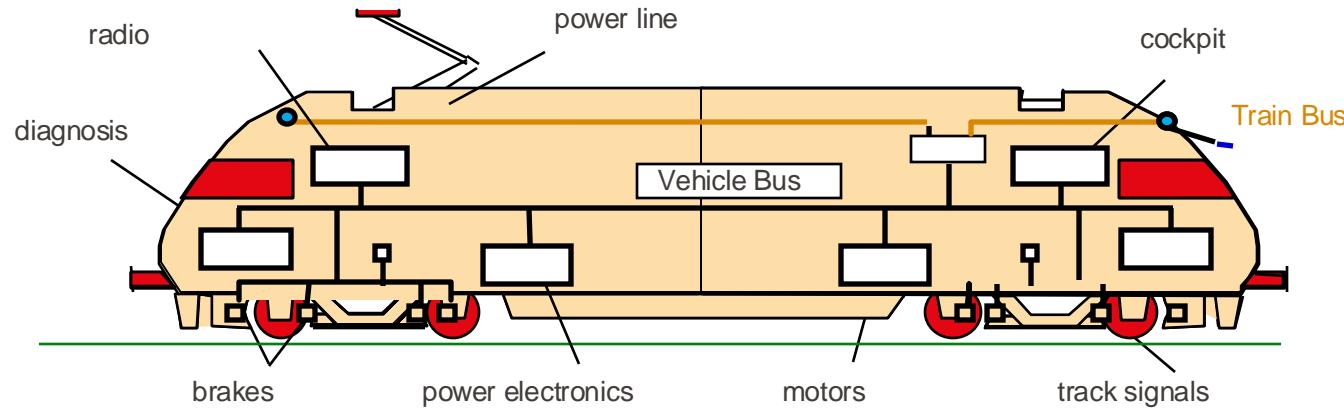
A data network, interconnecting an automation system, characterized by:

- many small data items (process variables) with **bounded delay** (1ms..1s)
- transmission of non-real-time traffic for commissioning and diagnostics
- harsh environment (temperature, vibrations, EM-disturbances, water, salt,...)
- robust and easy installation by skilled people
- high **integrity** (no undetected errors) and high **availability** (redundant layout)
- clock **synchronization** (milliseconds to microseconds)
- low attachment costs (€ 5.- .. €50 / node)
- moderate data rates (50 kbit/s - 5 Mbit/s), large distance range (10m - 20 km)

# Expectations: Why use a fieldbus?

- reduce cabling
- increased modularity of plant (each object comes with a CPU and communication interface)
- easy fault location and maintenance
- simplify commissioning
- simplify extension and retrofit
- off-the-shelf standard products to build “Lego”-control systems

# Fieldbus application example: locomotives and drives



data rate

1.5 Mbit/second

delay

1 ms (16 ms for skip/slip control)

medium

twisted wire pair, optical fibers (EM disturbances)

number of stations

up to 255 programmable stations, 4096 simple I/O

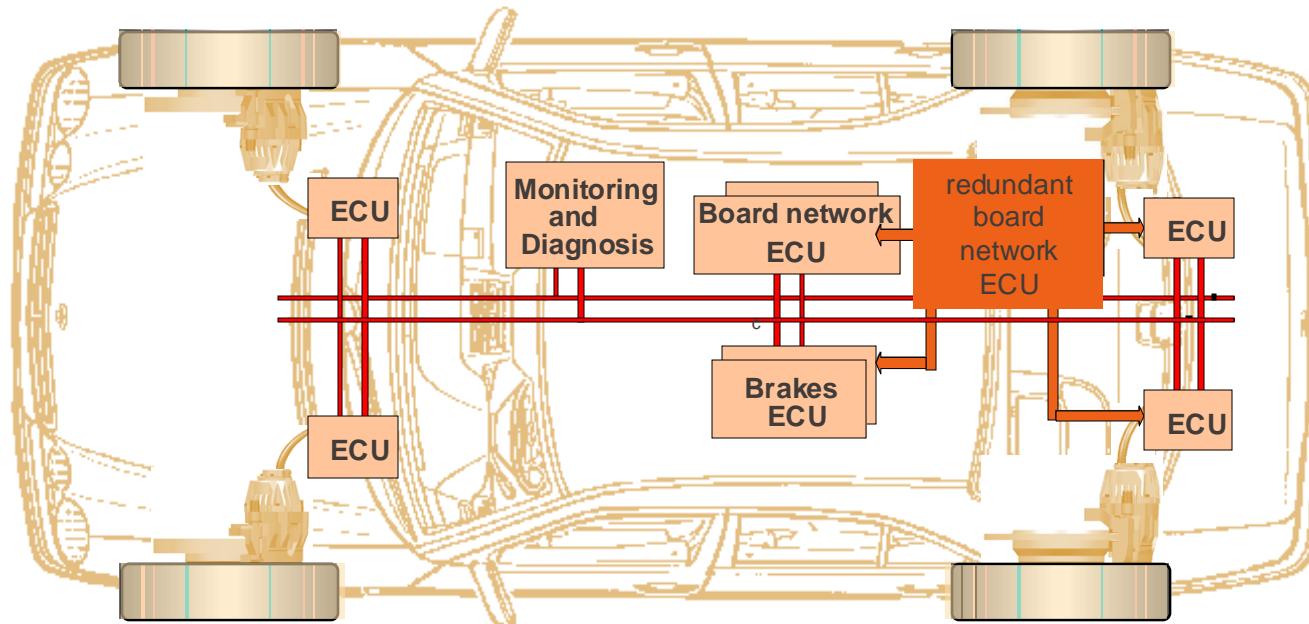
integrity

very high (signaling tasks)

cost

engineering costs dominate

# Fieldbus application example: car



DAIMLERBENZ  
AKTIENGESELLSCHAFT

- Redundant Engine Control Units
- Fault-tolerant 2-voltage on-board power supply
- Diagnostic System

# Fieldbus operation

Distribute **process variables** to all interested parties:

- source identification: requires a naming scheme
- accurate process value and units
- quality indication: {good, bad, substituted}
- time indication: how long ago was the value produced
- (optional description)

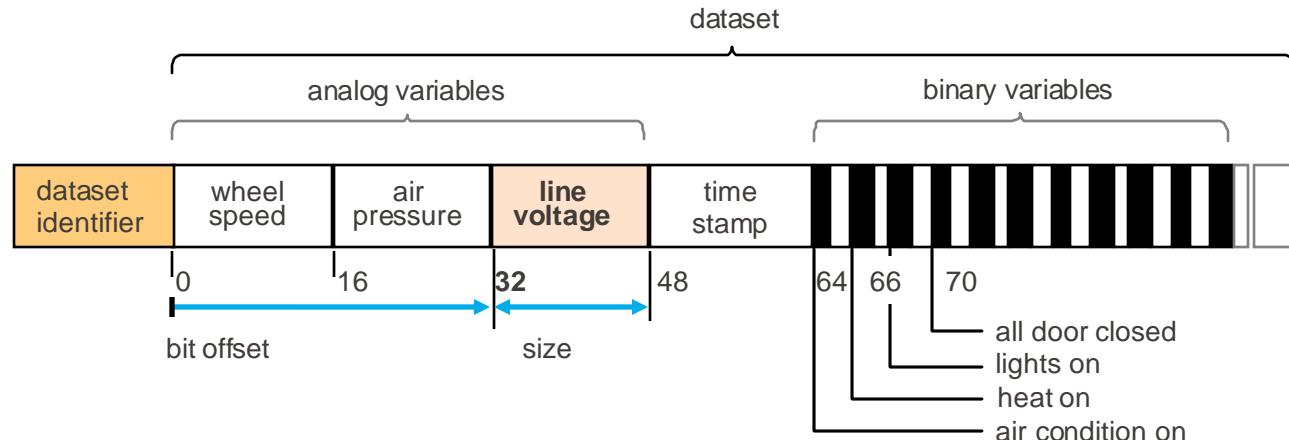
source	value	quality	time	description
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source	value	quality	time	description
minimum				

- In principle, the bus could transmit the process variable in clear text (even using XML, JSON, ..)
- However, this is quite expensive and only considered when the communication network offers some 100 Mbit/s and a powerful processor is available to parse the messages.
- More compact ways such as ASN.1 have been used in the past

ASN.1: (TLV)	type	length	value
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- Fieldbusses are slower (50kbit/s .. 12 Mbits/s) and thus more compact encodings are used.



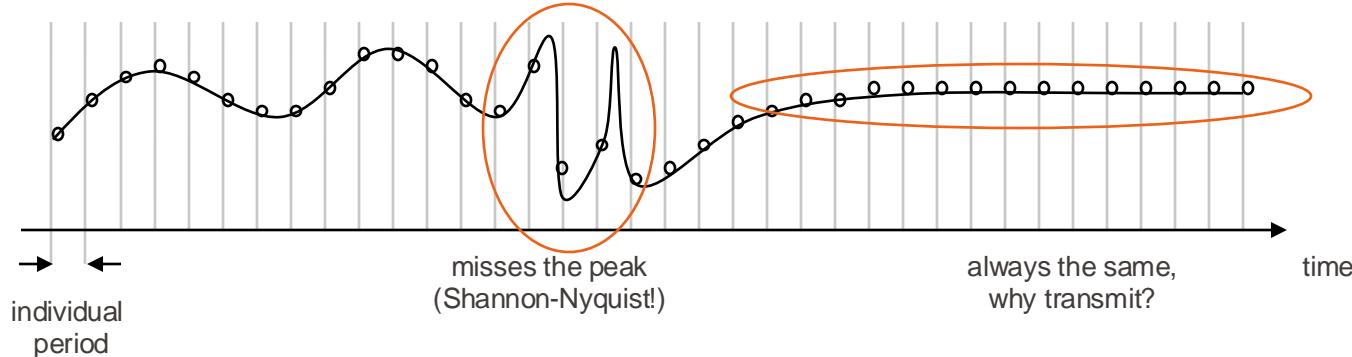
- Fieldbus devices have a low data rate and transmit always the same variables.
- It is economical to group variables of a device in the same frame as a dataset.
- A variable is identified within a dataset by its offset and its size.
- Variables may be of different types, types can be mixed.
- To allow later extension, room is left in the datasets for additional variables.

# Transmission principle: When to transmit

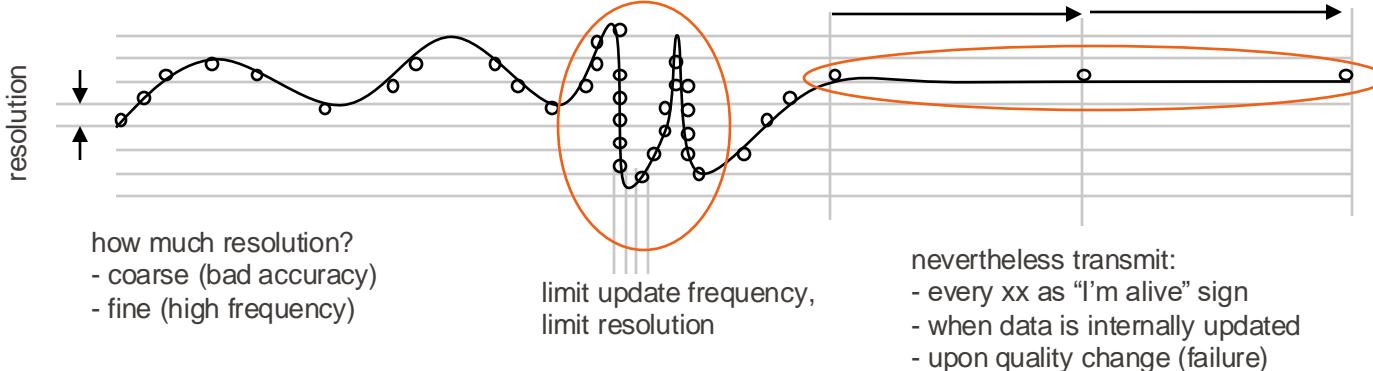
- The previous operation modes made no assumption on when data are transmitted.
- The actual network can transmit data:
  - cyclically (time-driven) or
  - on demand (event-driven),
  - or a combination of both.

# Cyclic versus Event-Driven transmission

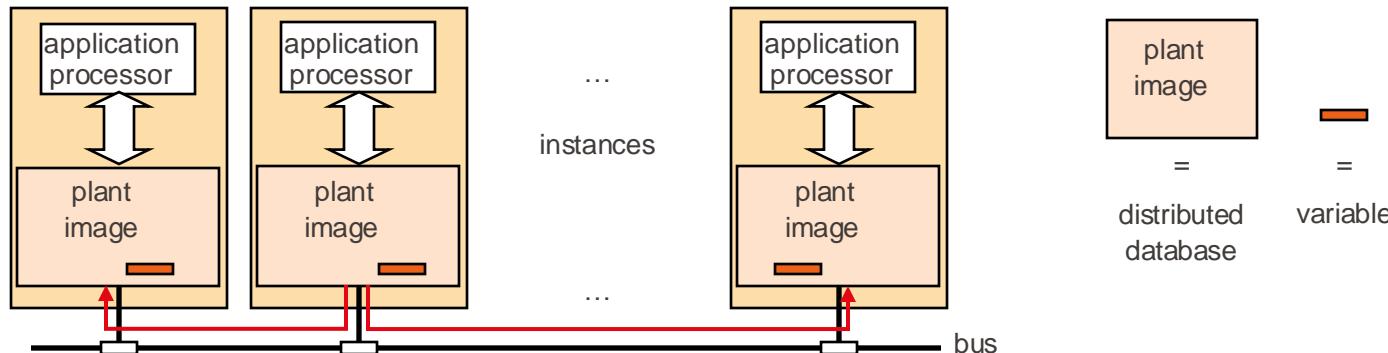
cyclic: send value strictly every xx milliseconds



event-driven: send when value change by more than x% of range



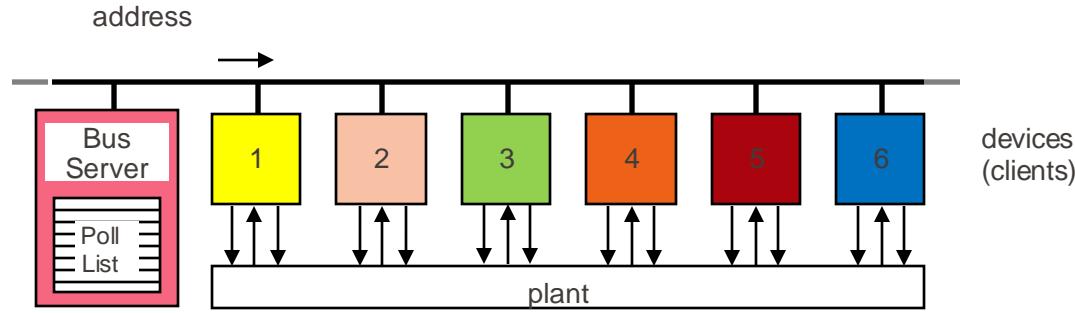
# Cyclic transmission



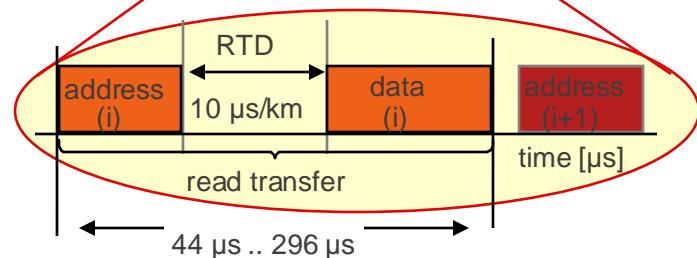
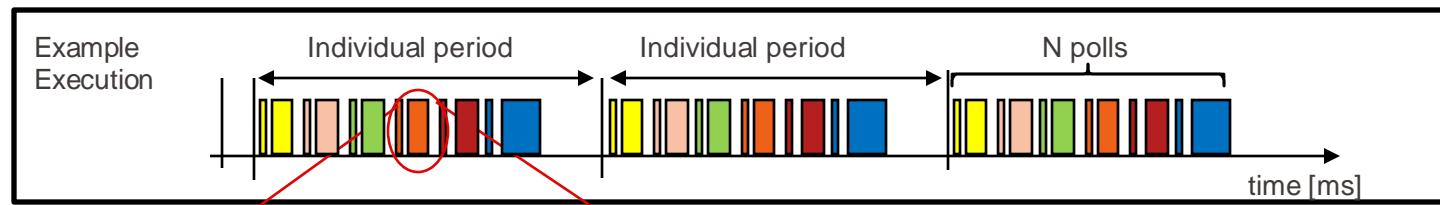
- Most variables are read in multiple devices.
- Broadcasting messages identified by their source (or contents) increases efficiency.
- Only one device is source of a certain process variable (otherwise collision).
- Bus refreshes plant image in the background.
- Replicated traffic memories can be considered as "caches" of the plant state (similar to caches in a multiprocessor system), representing part of the plant image.
- Each station snoops the bus and reads the variables it is interested in.

Each device is source or sink for some process variables

# Cyclic Data Transmission with Source Address Broadcast



Principle: server polls addresses in fixed sequence (poll list)



The duration of each poll is the sum of the transmission time of address and data (bit-rate dependent) and of the reply delay of the signals (round trip delay (RTD), independent of bit-rate).

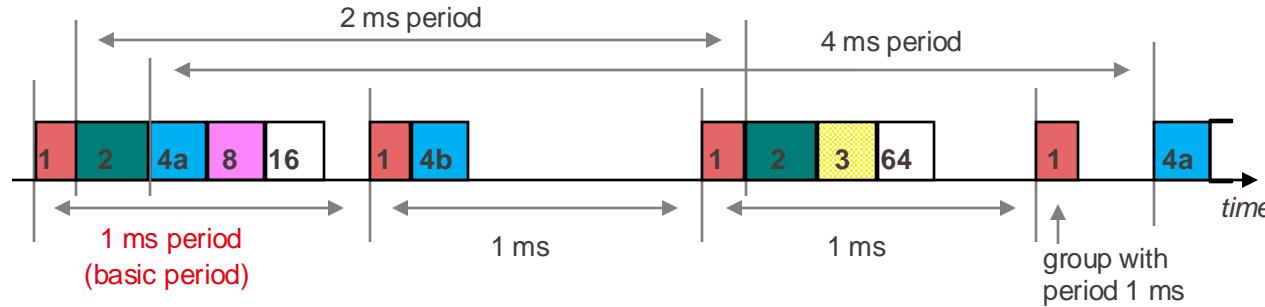
# Cyclic operation characteristics

1. Data are transmitted at **fixed intervals**, whether they changed or not.
2. The delivery delay (refresh rate) is deterministic and constant.
3. The bus is under control of a **central server (or distributed time-triggered algorithm)**.
4. **No explicit error recovery** needed since fresh value will be transmitted in next cycle. Use a freshness counter to detect stale data.
  - Consequence: cycle time **limited** by product of number of data transmitted and the duration of each poll (e.g.  $100 \mu\text{s} / \text{variable} \times 100 \text{ variables} \rightarrow 10 \text{ ms}$ )
  - To keep the poll time low, only **small data items** may be transmitted (< 256 bits)

The bus capacity and usage pattern must be configured beforehand. Determinism gets lost if the cycles are modified at run-time.

# Optimizing Cyclic Operation

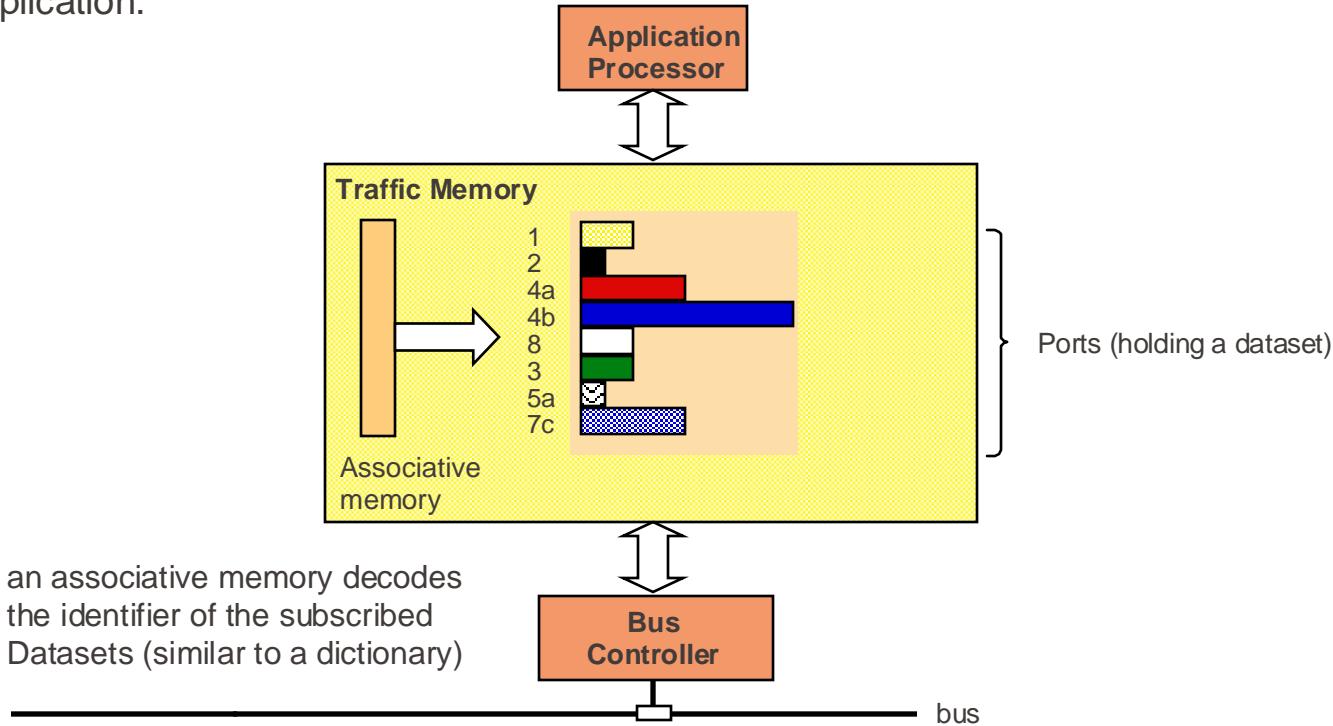
- **Problem:** fixed portion of the bus capacity used
  - load = #bytes transmitted per second
  - utilization = load/capacity
    - poll period increases with number of polled items
    - response time slows down
- **Solution:** introduce **sub-cycles** for less urgent periodic variables
  - length: power of 2 multiple of the base period.



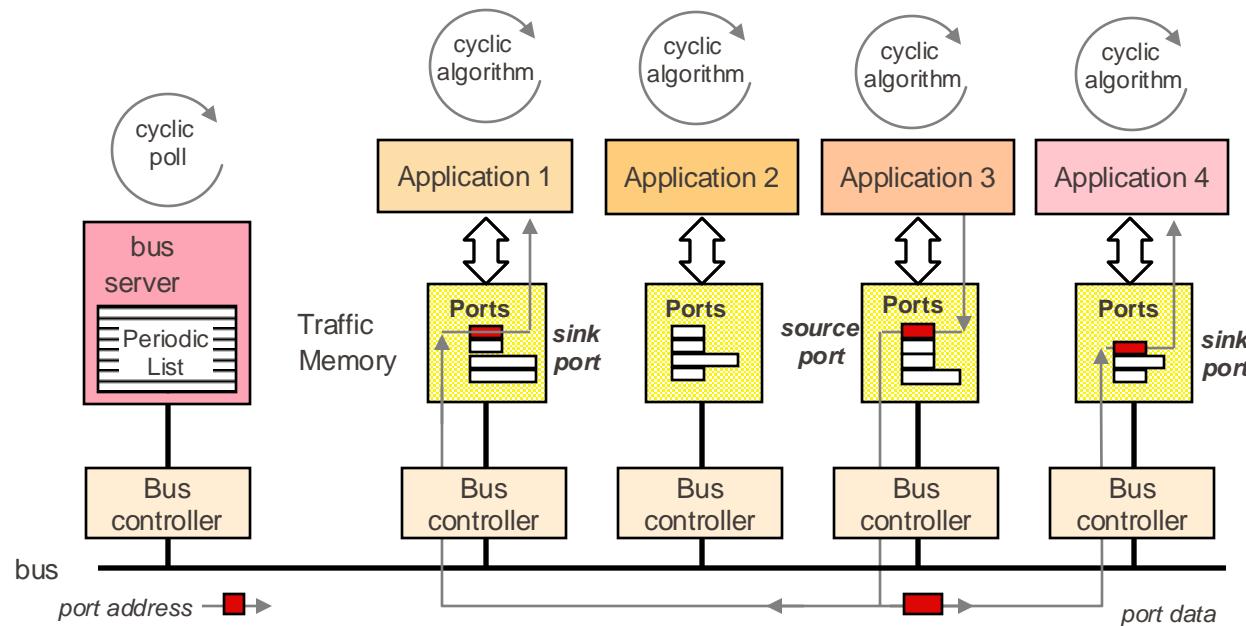
Poll cycles should not be modified at run-time (non-determinism)

# Traffic Memory: Implementation

Communication bus and application processor are decoupled by shared memory, the *Traffic Memory* (content addressable memory). Process variables are directly accessible by the application.



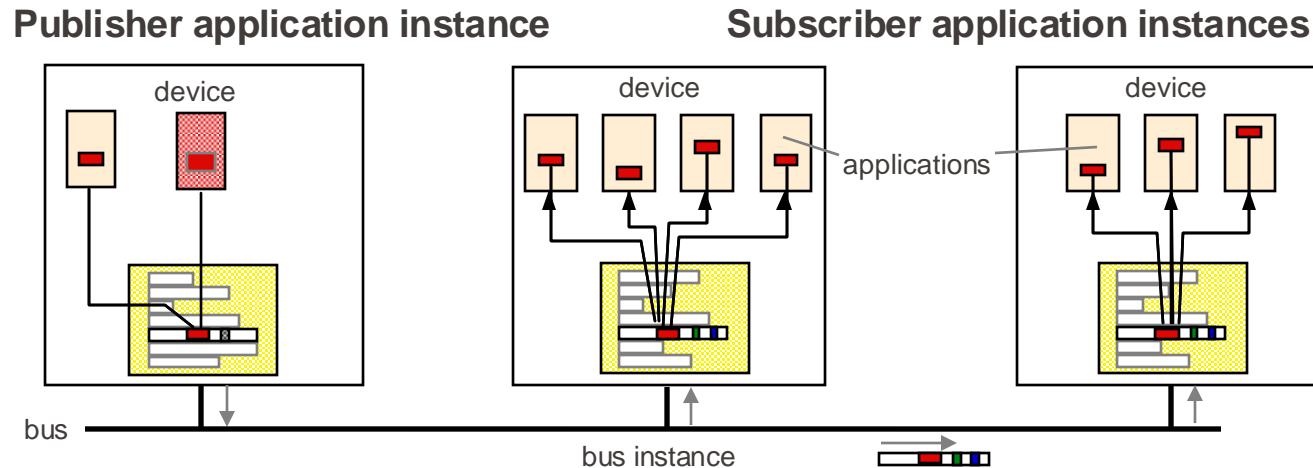
# Cyclic Transmission with Decoupled Application



The bus traffic and the application cycles are **asynchronous** to each other.  
The bus server scans the identifiers at its own pace.

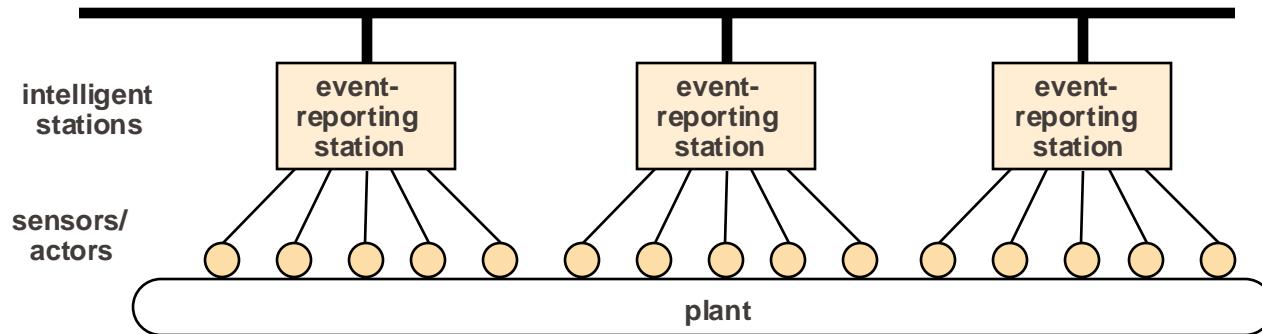
Deterministic behavior, at expense of reduced bandwidth and geographical extension.

# Example: delay requirement



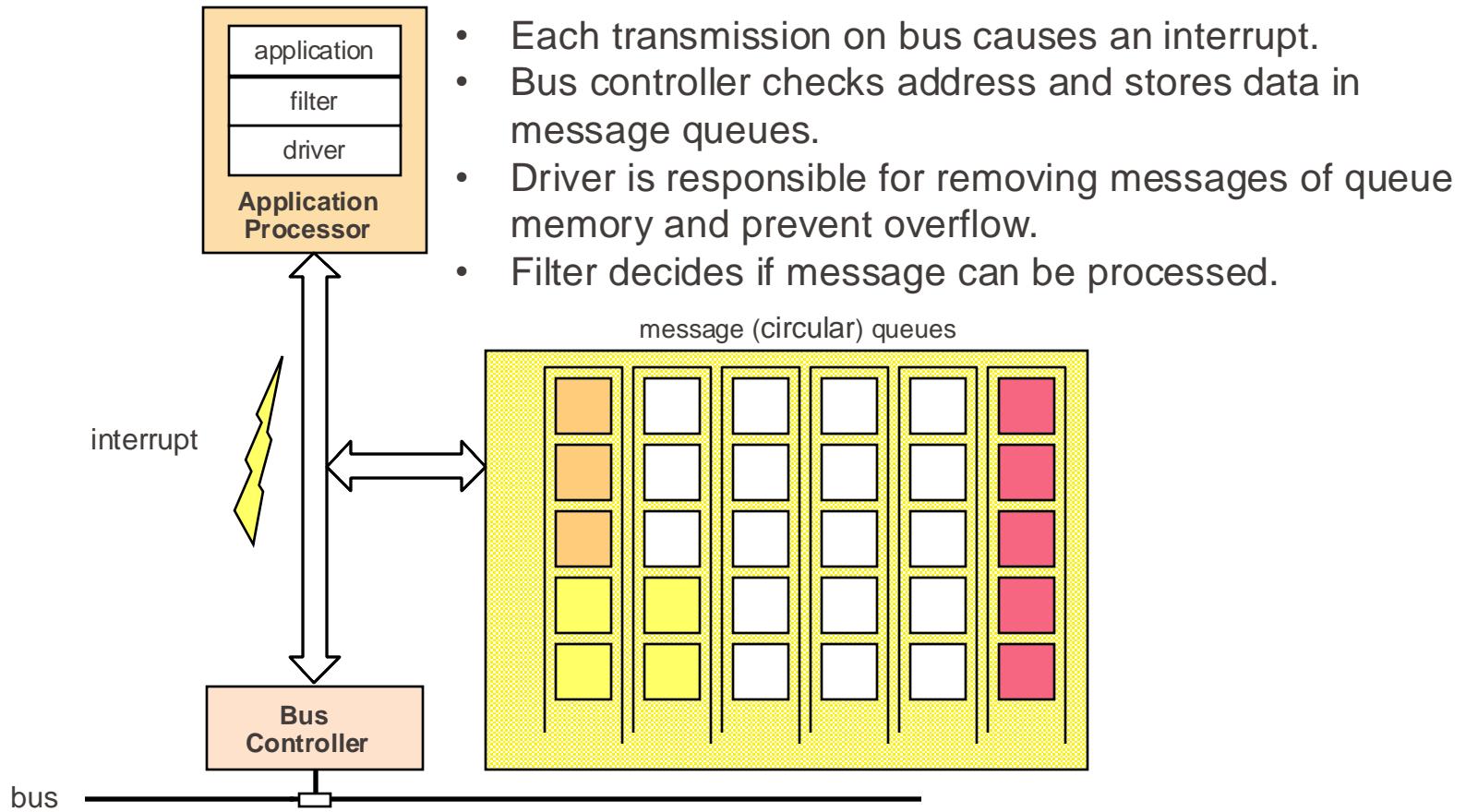
- **Worst-case delay** for transmitting all time critical variables is the sum of:
  - Source application cycle time 8 ms
  - Individual period of the variable on bus 16 ms
  - Sink application cycle time 8 ms
$$= 32 \text{ ms}$$

# Event-driven transmission



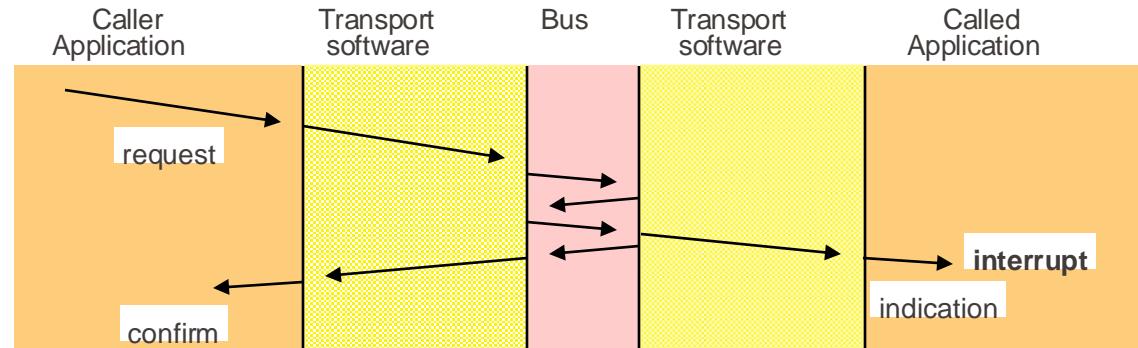
- Events cause transmission only when state changes.
- Bus load very **low on average**, but peaks under exceptional situations since transmissions are correlated by process.
- **Detection** of an event is an intelligent process:
  - Not every change of a variable is an event, even for binary variables.
  - Often, a **combination of changes** builds an event.
  - Only the **application** can **decide** what is an event, since only the application programmer knows the **meaning** of the variables.

# Bus interface for event-driven operation



# Response of Event-driven operation

- Since events can occur anytime on any device, stations communicate by **spontaneous** transmission, leading to possible **collisions**.
- Interruption of server at any instant can disrupt a time-critical task.
- **Buffering** of events can cause unbounded **delays**.
- The time required to transmit the event depends on the **medium access** (arbitration) procedure of the bus.



In an event-driven control system, there is only a transmission or an operation when an event occurs.

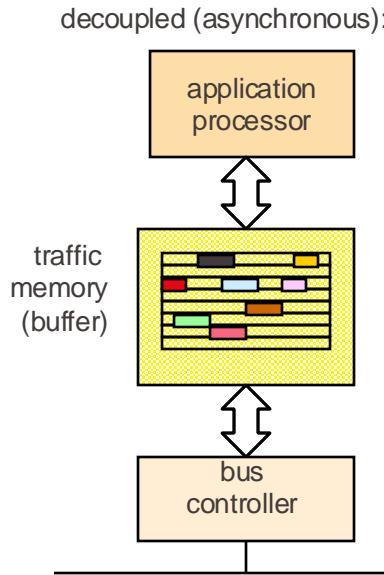
## Advantages:

- Can treat a large number of events, but not all at the same time
- Supports a large number of stations
- System idle under steady - state conditions
- Better use of resources
- Suitable for standard (interrupt-driven) operating systems (Unix, Windows)

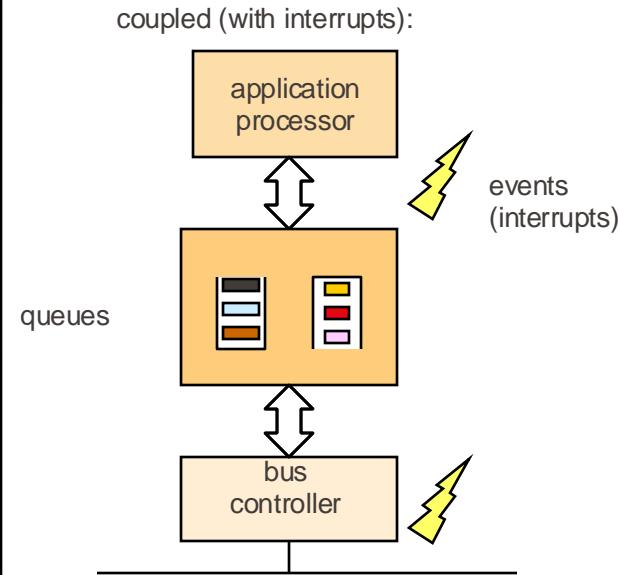
## Drawbacks:

- Needs shared access to resources (arbitration)
- No upper limit to access time if some component is not deterministic
- Response time difficult to estimate, requires analysis
- Limited by congestion effects: process correlated events
- A background cyclic operation is needed to check liveness

# Summary: Cyclic vs Event-Driven Operation



sending: application writes data into memory  
receiving: application reads data from memory  
the bus controller decides when to transmit  
bus and application are not synchronized



sending: application inserts data into queue and triggers transmission, bus controller fetches data from queue  
receiving: bus controller inserts data into queue and interrupts application to fetch them, application retrieves data

# How to combine cyclic and event-based transmissions?

## Cyclic Transmission

represent the **state of the plant**

short and urgent data items:

- motor current
- axle speed
- operator's commands
- emergency stops,...

Since variables are refreshed periodically, no retransmission protocol is needed to recover from transmission errors.

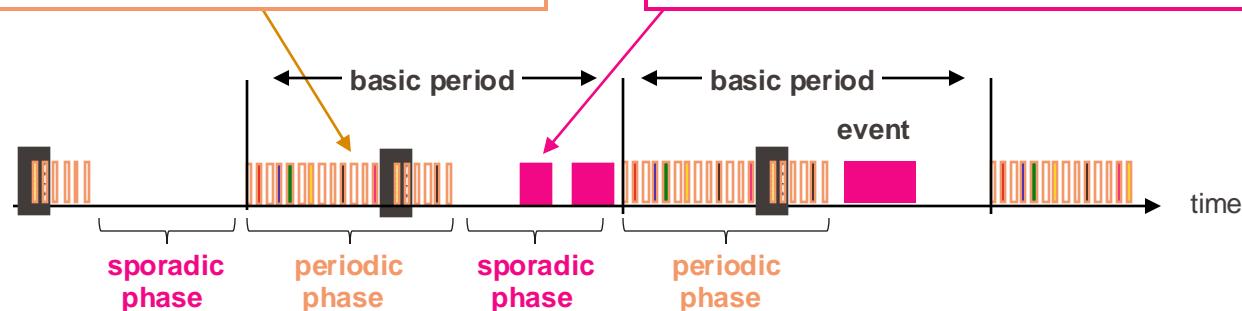
## Event-based Transmission

represent **state changes of the plant**

infrequent, sometimes long messages reporting events, for:

- Users: set points, diagnostics, status
- System: initialisation, down-loading, ...

Since messages represent state changes, a protocol must recover lost data in case of transmission errors.



# Cyclic or Event-driven Operation For Real-time?

## Cyclic operation:

- Data are transmitted at fixed intervals, whether they changed or not.
- Deterministic: delivery time is bound
- Worst case is normal case
- All resources are pre-allocated

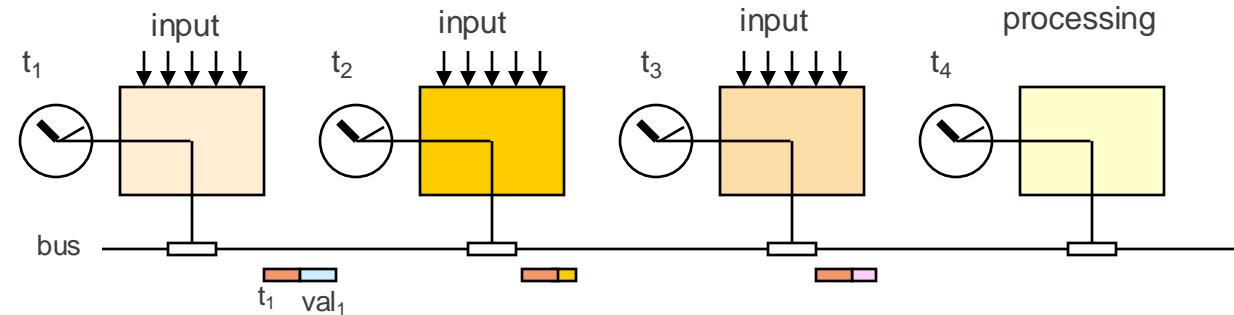
## Event-driven operation:

- Data are only transmitted when they change or upon explicit demand.
- Non-deterministic: delivery time vary widely
- Typical case works most of the time
- Best use of resources

# Time Synchronization

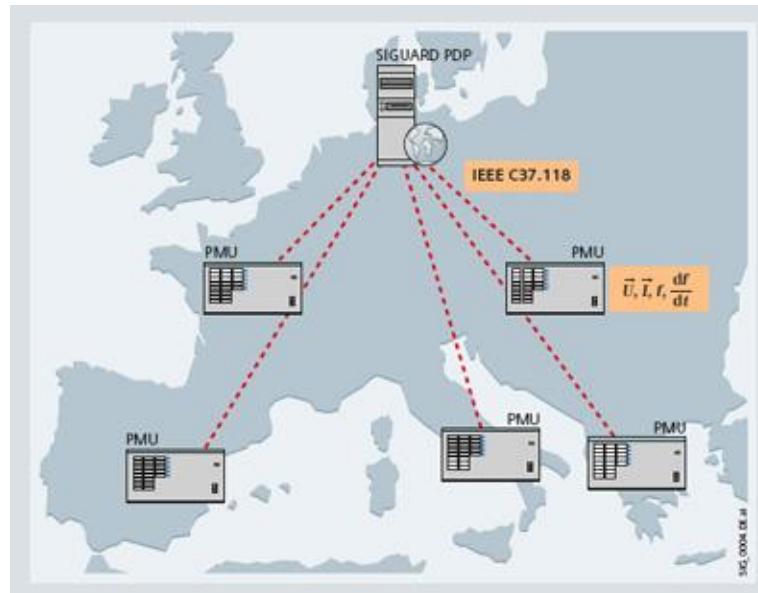
# Time-stamping and synchronisation

- In many applications, e.g. disturbance logging and sequence-of-events, the **exact sampling time of a variable must be transmitted together with its value**.
  - Devices equipped with clock recording creation time of value (not transmission time).
- To **reconstruct events coming from several devices**, clocks must be synchronized, considering transmission delays and failures.



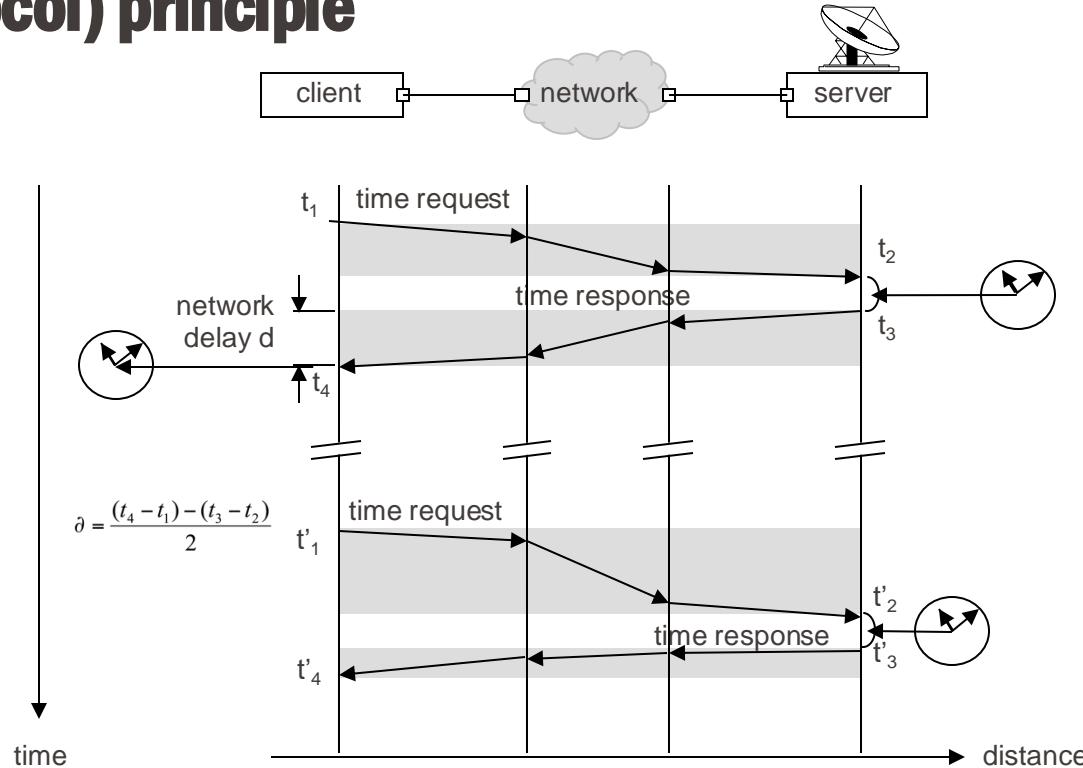
# Example: Phasor information

- Phasor transmission over the European grid: a phase error of 0,01 radian is allowed, corresponding to  $\pm 26 \mu\text{s}$  in a 60 Hz grid or  $31 \mu\text{s}$  in a 50 Hz grid.



- In client-server busses, server distributes time as bus frame.
- Clients can compensate for path delays, **time is relative to server**.
- In **demanding** systems, time is distributed over **separate lines** as **relative time**, e.g. PPS = one pulse per second, or **absolute time**, with accuracy of 1  $\mu$ s.
- In data networks, a **reference clock** (e.g. GPS or atomic clock) distributes time.
- A protocol evaluates **path delays** to compensate them:
  - NTP (Network Time Protocol): about 1 ms is usually achieved.
  - PTP (Precision Time Protocol, IEEE 1588): all network devices collaborate to estimate the delays, an accuracy below 1  $\mu$ s can be achieved without need for separate cables (but hardware support for time stamping required).

# NTP (Network Time Protocol) principle



Measures delay end-to-end over the network (one calculation)  
Problem: asymmetry in the network delays, long network delays

1. What is the difference between a centralized and a decentralized industrial bus?
2. What is the principle of source-addressed broadcast?
3. What is the difference between a time-stamp and a freshness counter?
4. Why is an associative memory used for source-addressed broadcast?
5. What are the advantages / disadvantages of event-driven communication?
6. What are the advantages / disadvantages of cyclic communication?
7. How is time transmitted? Why does it matter?

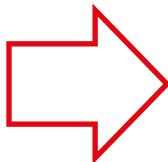
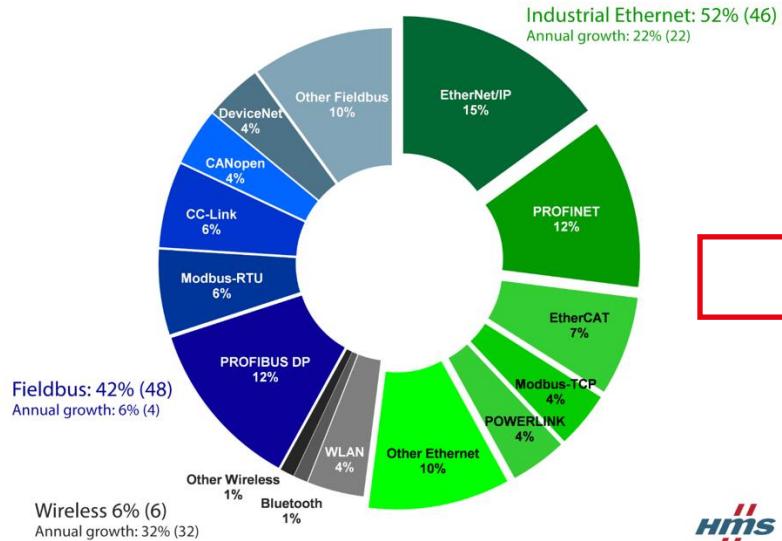
# Field Bus Standards



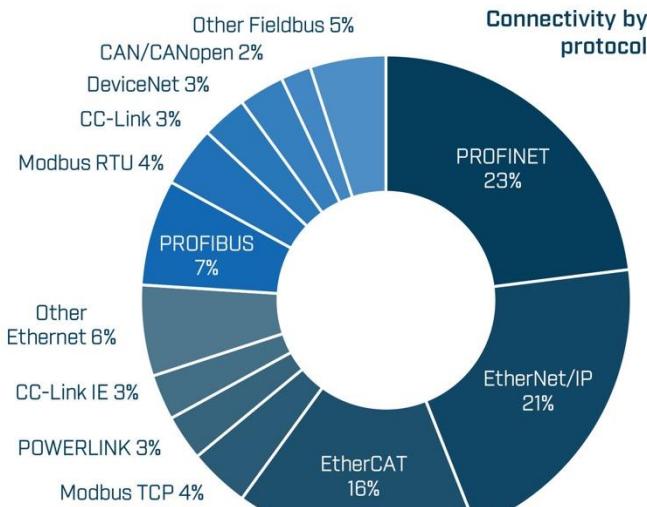
- Installed base, devices availability: processors, input/output
- Interoperability with devices from other vendors
- Topology and wiring technology (layout)
- Power distribution capabilities
- Connection costs per (input-output) point
- Response time
- Deterministic behavior
- Device and network configuration tools
- Bus monitor (baseline and application level) tools

# Worldwide most popular field busses

2018



2024



Market shares held by the leading fieldbus and industrial Ethernet systems  
Source: *HMS Industrial Networks*



- Profibus (Process Field Bus) is a field bus standard developed in the 1990s
- Two variants are in use today:
  - PROFIBUS DP (Decentralised Peripherals)
  - PROFIBUS PA (Process Automation)
- Twisted pair cables, bit rates from 9.6 kbit/s to 12 Mbit/s
- Cable length up to 1200m between repeaters

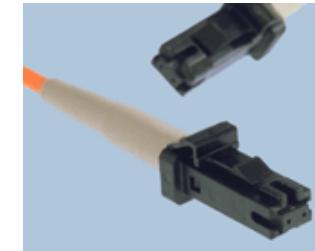


# The Industrial Ethernet „Standards“

- IEC SC65C „standardized“ 22 different, incompatible "Industrial Ethernets", driven by „market demand“.



EtherNet/IP	(Rockwell, OEDA)
Profibus, Profinet	(Siemens, PNO)
P-NET	(Denmark)
INTERBUS	(Phoenix)
Vnet/IP	(Yokogawa, Japan)
TCnet	(Toshiba, Japan)
EtherCAT	(Beckhoff, Baumüller)
Powerlink	(BR, AMK)
EPA	(China)
Modbus-RTPS	(Schneider, IDA)
SERCOS	(Bosch-Rexroth / Indramat)
...	

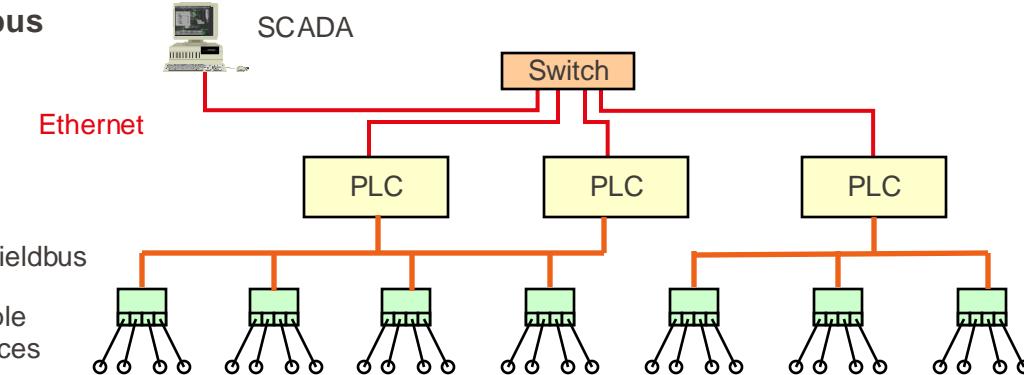


**EtherNet/IP**



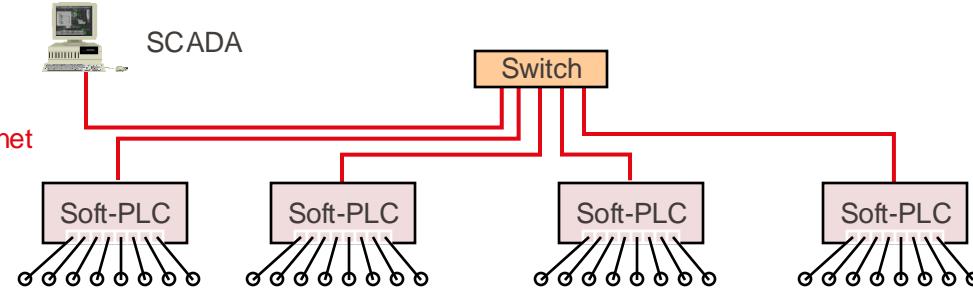
### Classical Ethernet + Fieldbus

cheap field devices  
decentralized I/O  
cyclic operation



### Ethernet as Fieldbus

costlier field devices  
Soft-PLC as concentrators  
Event-driven operation



# Ethernet & Time-Sensitive Networking (TSN)

- Set of standards developed by the IEEE 802 TSN task group
- Goal: Allow **time-sensitive** transmissions of data using **deterministic** Ethernet networks
- TSN standards suite provides:
  - Clock synchronization (PTP)
  - Deterministic communication (OSI Layer 2)
  - Traffic scheduling
  - Path control and reservation
  - Redundancy
- Classical Industrial Ethernet protocols (e.g. PROFINET, EtherNet/IP) will be built upon TSN-enabled hardware in the future.

# Ethernet-APL (Advanced Physical Layer)



- Ethernet 10BASE-T1L (10 MBit/s, full-duplex) over two-wire cables (reuse cables for 4-20mA bus)
- Loop-powered (devices supplied with power over the same cables)
- Long cable lengths (up to 1000 meters)
- Intrinsic safety
- Industry standard support by major manufacturers (Siemens, ABB, Emerson, Rockwell Automation, etc.)

- Non-time critical busses are being displaced by LANs (Ethernet) and cheap peripheral busses (USB, ...)
- The cabling objective of field busses (more than 32 devices over 400 m) is out of reach for cheap peripheral busses such as USB.
- Field busses tend to live very long (10-20 years), contrarily to office products.
- There is no real incentive from the control system manufacturers to reduce the fieldbus diversity, since the fieldbus binds customers.
- TSN-based hardware will be widely available leading to lower costs for industrial networking equipment.
- Wireless plays a niche role, but will enable new applications (e.g. mobile robots connected through 5G)