

# Principles of Digital Communication

May 28th, 2021

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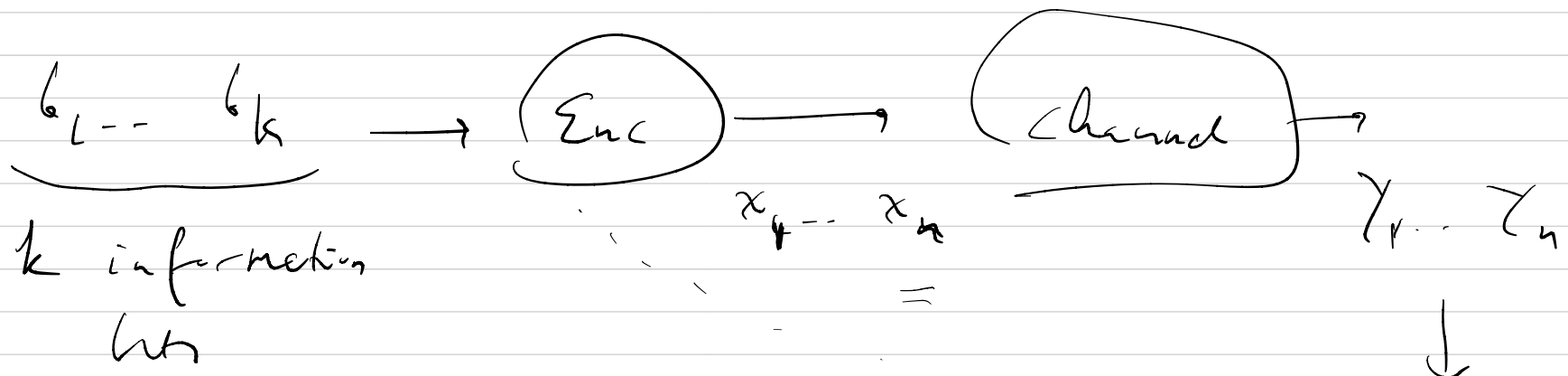
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Question "How to design a convolutional code"?

- General question: How to design a code?

- Problem in the abstract:



$$\vec{y} = \vec{x} + \vec{z}$$

$$\vec{z} \sim N(0, 1)$$

We need to design Enc/Dec

subject to constraints:

$$k/n \text{ large} \Rightarrow R$$

$$\frac{1}{n} \sum_{i=1}^n x_i^2 \text{ small} \Rightarrow P$$

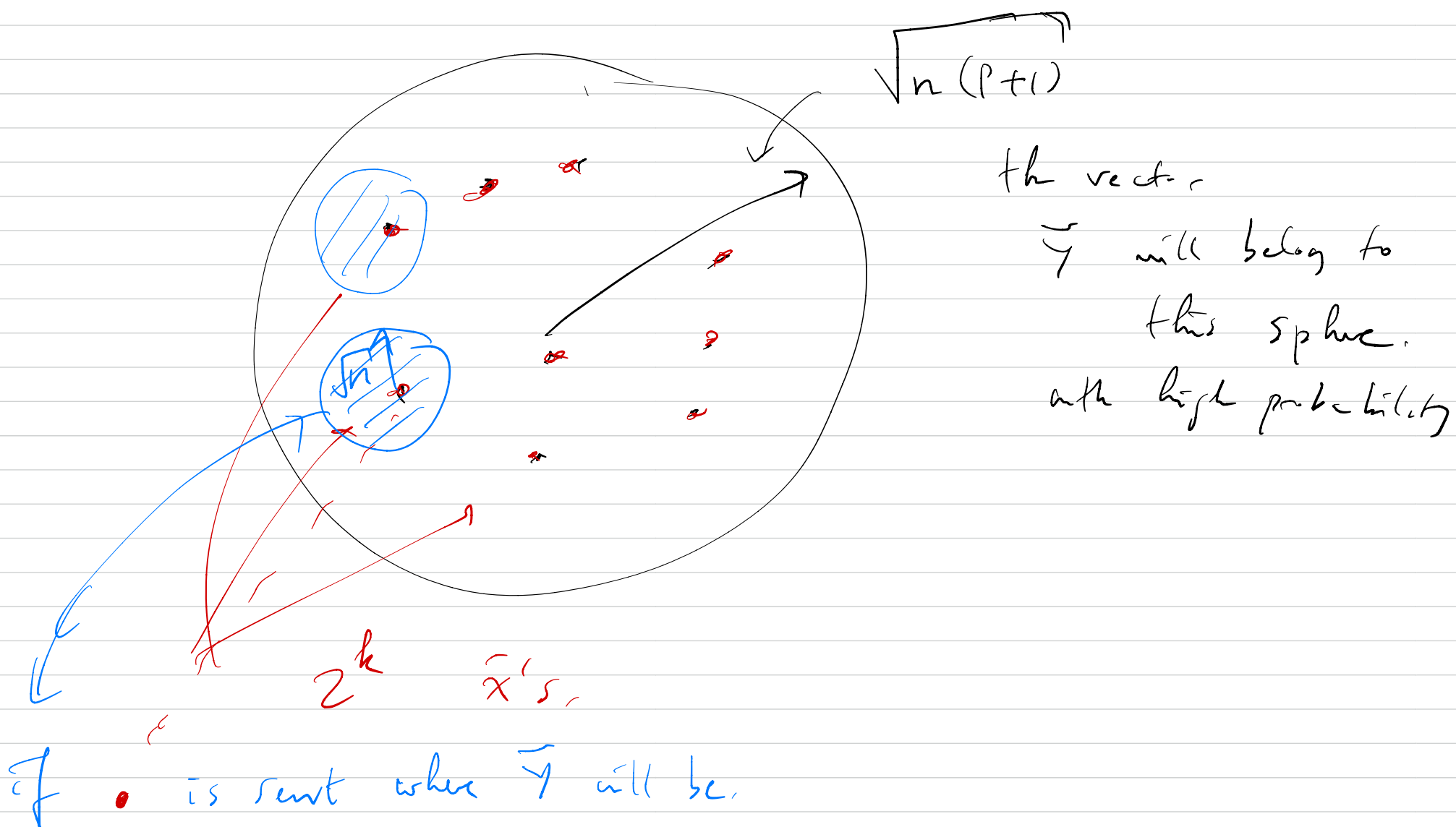
$$\Pr(\widehat{b_1 \dots b_k} \neq b_1 \dots b_k) \text{ small } P_{\text{error}}$$

Primary question:

given  $R, P, P_{\text{error}}$ , is there an enc/dec

with these parameters?

Recall from some weeks ago



Reliable communication ( $P_{\text{error}}$  is small) requires that

Blue spheres are ~ disjoint.

$$\Rightarrow 2^k (\sqrt{n})^n \leq (\sqrt{(1+P)n})^n$$

$$\Rightarrow 2^R \leq \sqrt{1+P} \Rightarrow R \leq \frac{1}{2} \log_2(1+P)$$

Ex:  $P \leq 1 \quad \sigma^2 = 10$

$$R \leq \frac{1}{2} \log_2\left(1 + \frac{1}{10}\right) \approx \frac{1}{14} \approx 0.07$$

$$n = 60000 \Rightarrow k \leq \underline{\underline{4000}}$$

Question: how far is " $\frac{1}{2} \log(1+P)$ " from the true value of best possible R.

Ans: for any  $P_e > 0$  (no matter how small) any  $R < \frac{1}{2} \log(1+P)$  (but arbitrarily close to  $\frac{1}{2} \log(1+P)$ ) for any P.

there exists an enc/dec s.t.

$$\left( \begin{array}{c} k \\ n \end{array} \right) > R \text{ \& } P_{\text{error}} < P_e \}$$

$$\text{and } \frac{1}{n} \sum x_i^2 \leq P.$$

- More "fine" question: do there exist low-complexity (enc/dec) s that approach  $\frac{1}{2} \log(1+P)$  limit?

- Ans: yes

- "Constructive" methods to design the enc/dec.

- Remarks: if we restrict ourselves to convolutional encoders & viterbi decoder, the " $\frac{1}{2} \log(1+P)$ "

needs to be replaced by  $\frac{1}{2} \log(1 + \frac{P}{2})$  (??)

- even with this it is not so clear how to choose

