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## Quantum Information and Quantum Computing, Problem set 7

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### Problem 1 : Shor's factoring algorithm and code

Our first goal is to implement on Qiskit the simplest instance of the period-finding code, so to factor the number  $N = 15$ . Assume that in step 3 of Shor's algorithm we have randomly chosen  $x = 4$  among the co-prime numbers of  $N$ , i.e. such that  $\gcd(x, N) = 1$ . After this simple case, you are going to design a quantum circuit for the more complex case  $N = 21$ .

1. Devise a quantum circuit that executes the modular exponentiation  $f(z) = x^z \bmod 15$ , with  $x = 4$ . We will assume that  $z$  is a 2-qubit number. This means that we choose  $t = 2$  as the number of qubits in the first register of the quantum period finding algorithm. This assumption is justified if we expect to find a small period, which is the case here.
2. Implement on Qiskit the period-finding algorithm with the modular exponentiation code that you just found. Show that it finds a factor of  $N = 15$ .
3. Now do the same for  $N = 21$  and  $x = 11$ . For this case it is difficult to devise a modular exponentiation circuit. Therefore, we will assume that modular multiplication is available as a primitive and only solve the problem on paper.
  - What is the period of  $x = 11$  for  $N = 21$ ? How many qubits  $t$  are needed in the phase register to correctly execute Shor's algorithm? How many were necessary in the case  $N = 15$ ,  $x = 4$ ?
  - What is the added algorithmic difficulty of this second factoring case?
  - Write the quantum circuit of Shor's algorithm for  $N = 21$  and  $x = 11$  using the one-qubit semiclassical QFT (the one with mid-circuit measurements). Assume that you have the controlled modular multiplication operations  $c - U_{2^k}$  available as block operations. Here  $U_{2^k}$  is the permutation unitary that achieves  $|y\rangle \rightarrow |yx^{2^k} \bmod N\rangle$ . Notice that the  $U_{2^k}$  are permutation matrices because  $x$  and  $N$  are coprimes. Also, remember from class that to correctly define these unitaries for Shor's factoring algorithm, you have to set them to the identity for  $y \geq N$ .
  - Optional: If you are skilled in Qiskit, define the  $U_{2^k}$  and  $c - U_{2^k}$  using the Qiskit class `qiskit.circuit.library.PermutationGate`.