

Lecture 9

The Laplacian matrix and consensus in continuous time

Textbook , §6, §7

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Setup

• $G = (V, E, w)$ is a weighted digraph, $V = \{1, \dots, n\}$

▶ A is the adjacency matrix

▶ $D_{\text{out}} = \begin{bmatrix} d^{\text{out}}(1) & & 0 \\ & \ddots & \\ 0 & & d^{\text{out}}(n) \end{bmatrix}$ is the out-degree matrix

▶ Standing assumption: $w_{ij} > 0$ if $(i, j) \in E$

$$d^{\text{out}}(v) = \sum_{(v, j) \in E} w_{v, j}$$



Definition

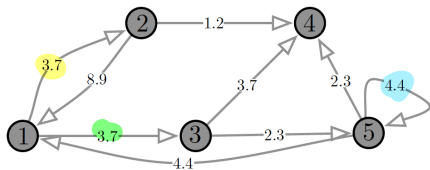
The Laplacian matrix of G is $L = D_{\text{out}} - A$

Remark

$$L_{ij} = \begin{cases} -A_{ij} & i \neq j \\ \sum_{h=1, h \neq i}^n A_{ih} & i = j \end{cases} \rightarrow \text{Sum of the off-diagonal entries of } A$$

sh. v. r. ↓

Example



$$L = \begin{bmatrix} 7.4 & -3.7 & -3.7 & 0 & 0 \\ -8.9 & 10.1 & 0 & -1.2 & 0 \\ 0 & 0 & 6.0 & -3.7 & -2.3 \\ 0 & 0 & 0 & 0 & 0 \\ -4.4 & 0 & 0 & -2.3 & 6.7 \end{bmatrix}$$

Properties of L

- Off-diagonal elements are ≤ 0 . $L_{ij} \geq 0$
- $L_{ij} = 0$ only if i is a sink (up to a self-loop)
- Zero row-sums: $L\mathbb{1}_n = 0$
 - ▶ $\lambda = 0$ is an eigenvalue of L and $\mathbb{1}_n$ is the associated eigenvector
- Self-loops are invisible from L
- L symmetric $\Leftrightarrow A$ symmetric

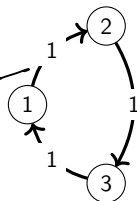
Balanced graphs and undirected graphs

Definition: The digraph G is balanced if $d^{\text{out}}(v) = d^{\text{in}}(v), \forall v \in V$

G balanced $\nRightarrow A$ symmetric

$$A = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix}$$

*not symmetric
but G balanced*



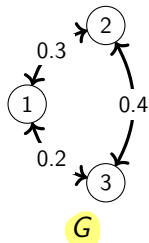
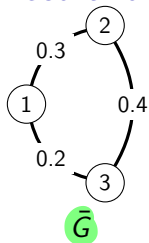
Definition: A digraph $G = (V, E, w)$ and an undirected, weighted graph $\bar{G} = (\bar{V}, \bar{E}, \bar{w})$ are associated with each other if $V = \bar{V}$ and

- $(i, j) \in \bar{E} \Rightarrow (i, j) \text{ and } (j, i) \in E$
- G has no self loops
- $\bar{w}_{ij} > 0 \Rightarrow w_{ij} = w_{ji} = \bar{w}_{ij}$

Proposition: Let G be associated to \bar{G} and let A and \bar{A} be the respective adjacency matrices. Then

- $A = \bar{A}$ and they are symmetric
- G is balanced

Balanced and undirected graphs



\downarrow

$$\bar{A} = \begin{bmatrix} 0 & 0.3 & 0.2 \\ 0.3 & 0 & 0.4 \\ 0.2 & 0.4 & 0 \end{bmatrix}$$

\downarrow

$$A = \begin{bmatrix} 0 & 0.3 & 0.2 \\ 0.3 & 0 & 0.4 \\ 0.2 & 0.4 & 0 \end{bmatrix}$$

\Rightarrow Same adjacency matrices \rightarrow Laplacians are the same

Definition: The Laplacian matrix of an undirected weighted graph $\bar{G} = (V, \bar{E}, \bar{w})$ is the Laplacian matrix of the associated digraph.

Balanced and undirected graphs

Remark

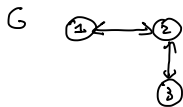
\bar{G} undirected and connected \Rightarrow the associated digraph G is strongly connected (because of bidirectional edges)

Proposition

If the digraph G has a Globally Reachable Node (GRN) and is associated to an undirected graph \bar{G} , then

- (a) G is strongly connected
- (b) \bar{G} is connected

Sketch of the proof of (a)



If ① is a GRN in G , then one can go from any $v \in V$ to ① following the edges in G . But since we also have the opposite edges (because G is associated to \bar{G}) one can also go from ① to v .

Fix $i, j \in V$, $i \neq j$. Then there is a path $i \dots \text{①}$ and a path $\text{①} \dots j \rightarrow$ there is a path from i to j

Back to the general case

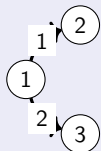
Definition: (Laplacian without using graphs) $L \in \mathbb{R}^{n \times n}$, $n \geq 2$ is a Laplacian matrix if

- all row sums are zero
- diagonal entries are ≥ 0
- off-diagonal entries are ≤ 0

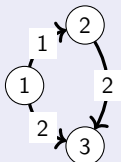
Remark: To every L one can associate a unique digraph, up to self-loops

Example:

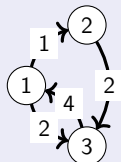
$$L = \begin{bmatrix} 3 & -1 & -2 \\ 0 & 2 & -2 \\ -4 & 0 & 4 \end{bmatrix}$$



From first row
(off-diagonal entries)



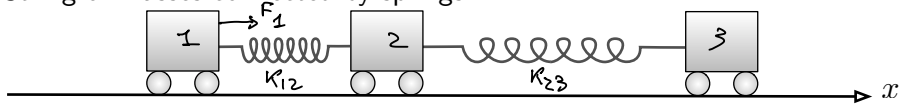
From second row
(and first too)



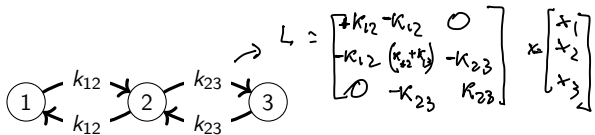
From third row

Examples of Laplacians in physical systems

String of masses connected by springs



- Associated graph



- Mass i subject to the elastic force: $F_i = \sum_{j \in \mathcal{N}^{out}(i)} k_{ij}(x_j - x_i) = -(Lx)_i$
- Total elastic energy: $E = \frac{1}{2} \sum_{(i,j) \in \mathcal{E}} k_{ij}(x_i - x_j)^2 = \frac{1}{2} x^T L x$
- Dynamics of the mass i : $M_i \ddot{x}_i = -(Lx)_i$, $M_i > 0$

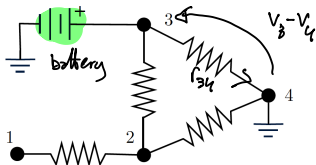
▶ Collective dynamics setting $x = [x_1, \dots, x_n]^T$, $M_i = 1$, $i = 1, \dots, n$

$$\ddot{x} = -Lx$$

$$\begin{aligned} F_1 &= k_{12}(x_2 - x_1) \\ F_2 &= k_{23}(x_3 - x_2) + k_{12}(x_1 - x_2) \\ &= k_{12}x_1 - (k_{12} + k_{23})x_2 + k_{23}x_3 \end{aligned}$$

2h L12 ↓

Electrical network of resistors



- Current from i to j

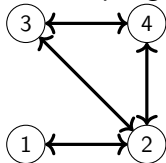
$$C_{i \rightarrow j} = \frac{V_i - V_j}{r_{ij}}$$

$$= a_{ij}(V_i - V_j)$$

$\hookrightarrow \frac{1}{r_{ij}}$

- r_{ij} = resistance, r_{ij}

- Associated coupling graph



Weights a_{ij}



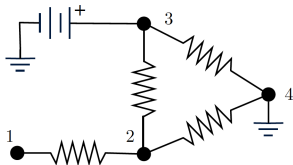
Adjacency matrix A

- Current injected into i : $c_i = - \sum_{j \in \mathcal{N}^{out}(i)} a_{ij}(V_i - V_j) = -(LV)_i$

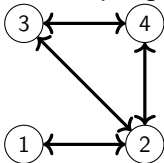
▶ Setting $V = [V_1, \dots, V_n]$ and $c = [c_1, \dots, c_n]$ one has

$$c = -LV$$

Electrical network of resistors



- Associated coupling graph



- Current from i to j

$$C_{i \rightarrow j} = \frac{V_i - V_j}{r_{ij}} \\ = a_{ij}(V_i - V_j)$$

- r_{ij} = resistance,

Weights a_{ij}



Adjacency matrix A

- Power dissipated by a resistor: $C_{i \rightarrow j}(V_i - V_j)$
- Total dissipated power: $P = \sum_{(i,j) \in \mathcal{E}} a_{ij}(V_i - V_j)^2 = V^T L V$

Properties of Laplacian matrices

G : weighted digraph with n nodes $\rightarrow L$: Laplacian matrix, A : Adjacency matrix

Remark: G is balanced $\Leftrightarrow D_{out} = D_{in}$, where

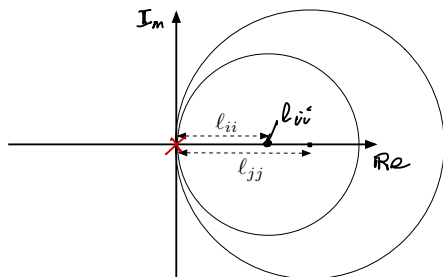
$$D_{in} = \begin{bmatrix} d^{in}(1) & & 0 \\ & \ddots & \\ 0 & & d^{in}(n) \end{bmatrix} \text{ is the in-degree matrix}$$

Lemma (zero column sum)

G is balanced $\Leftrightarrow \mathbb{1}_n^T L = [0, \dots, 0]$

Proof: At home

Properties of Laplacian matrices



Lemma (spectrum of L)

The **nonzero** eigenvalues of L have **strictly-positive real part**

Proof: Recall the standing assumption that elements a_{ij} of the adjacency matrix are nonnegative. In row i of L , $l_{ii} = \sum_{j \neq i, j=1}^n a_{ij} \geq 0$, $l_{ij} = -a_{ij} \leq 0$

- Gersgorin Disks Theorem: $\text{Spec}(L) \subset \cup_{i=1}^n D_i$
- $D_i = B(l_{ii}, r_i)$, $r_i = \sum_{j=1, j \neq i}^n |l_{ij}| = \sum_{j=1, j \neq i}^n a_{ij} = l_{ii}$
↳ closed balls

Properties of Laplacian matrices

Remark

- G undirected $\Rightarrow L = L^T \Rightarrow$ real eigenvalues. From the Lemma:

$$0 = \lambda_1 \leq \lambda_2 \leq \dots \leq \lambda_n$$

- ▶ λ_2 is called the Fiedler eigenvalue
- If $\lambda = 0$ is simple, its eigenspace is the consensus subspace $\alpha \mathbb{1}_n, \alpha \in \mathbb{R}$. When $\lambda = 0$ is simple?

Theorem

Let d be the number of sinks in the condensation graph $C(G)$. Then, $\text{rank}(L) = n - d$.

Remark

$$\rightarrow \text{rank}(L) = n - \overbrace{\dim(\text{Ker}(L))}^d$$

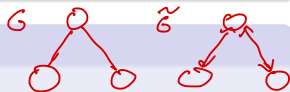
Recall the rank/nullity theorem: $\text{rank}(L) + \dim(\text{Ker}(L)) = n$, This shows that d is the geometric multiplicity of $\lambda = 0$.

Properties of Laplacian matrices

- **Theorem 1:** $d = 1 \Leftrightarrow G$ has a GRN (globally reachable node) \Leftrightarrow the eigenvalue $\lambda = 0$ is simple
- **Corollary 2:** Assume G is undirected. Then $\lambda = 0$ is simple $\Leftrightarrow G$ is connected.

Proof of Corollary 2

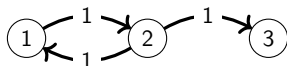
Proof: Let \tilde{G} be the digraph associated to G



- \Rightarrow • $\lambda = 0$ is simple $\stackrel{\text{Thm 1}}{\Rightarrow} d = 1$ (number of sinks of $C(\tilde{G})$) $\Rightarrow \tilde{G}$ is strongly connected $\Rightarrow G$ connected.
- \Leftarrow • G connected $\Rightarrow \tilde{G}$ strongly connected and $C(\tilde{G})$ has a single sink.

Therefore, from Thm 1, $\lambda = 0$ is simple

Example



- G has a globally reachable node
 $\Rightarrow \lambda = 0$ is a simple eigenvalue of L
- All other eigenvalues have positive real parts

Check:

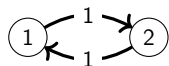
$$A = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 0 & 0 \end{bmatrix} \quad L = D^{\text{out}} - A = \begin{bmatrix} 1 & -1 & 0 \\ -1 & 2 & -1 \\ 0 & 0 & 0 \end{bmatrix} \quad \text{Spec}(L) = \{0, 0.382, 2.61\}$$

Computation of $\text{Ker}(L)$

$$L \begin{bmatrix} v_1 \\ v_2 \\ v_3 \end{bmatrix} = 0 \quad \begin{cases} v_1 - v_2 = 0 \\ 2v_2 - v_1 - v_3 = 0 \end{cases} \rightarrow \begin{cases} v_1 = v_2 \\ v_2 = v_3 \end{cases}$$

$$\begin{bmatrix} v_1 \\ v_2 \\ v_3 \end{bmatrix} = \begin{bmatrix} \alpha \\ \alpha \\ \alpha \end{bmatrix}, \quad \alpha \in \mathbb{R}$$

Example



$C(G)$



- H_1 is the subgraph induced by nodes $\{1, 2\}$ and $H_2 = (\{3\}, \emptyset)$
- From the theorem, since $C(G)$ has two sinks $\rightarrow \downarrow = \mathcal{Z}$
 - ▶ $\text{Rank}(L) = 3 - 2 = 1 \rightarrow \text{dim}(\text{Ker}(L)) = 2$

Check:

$$A = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix} \quad L = D^{\text{out}} - A = \begin{bmatrix} 1 & -1 & 0 \\ -1 & 1 & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

$$\text{Spec}(L) = \{2, 0, 0\}$$

$$L \begin{bmatrix} v_1 \\ v_2 \\ v_3 \end{bmatrix} = 0 \quad \begin{cases} v_1 - v_2 = 0 \\ -v_1 + v_2 = 0 \end{cases} \rightarrow \begin{bmatrix} v_1 \\ v_2 \\ v_3 \end{bmatrix} = \begin{bmatrix} \alpha \\ \alpha \\ \beta \end{bmatrix}, \quad \alpha, \beta \in \mathbb{R}$$

The Laplacian flow

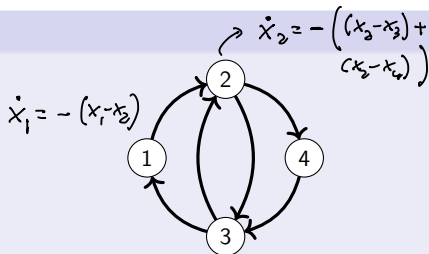
$$\dot{x} = -Lx \quad (1)$$

$x(t) \in \mathbb{R}^n$, $L \in \mathbb{R}^{n \times n}$ Laplacian matrix

Why is it interesting?

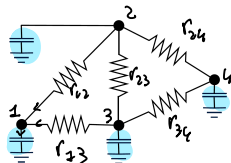
Agent dynamics: from (1),
assuming unit weights,

$$\dot{x}_i = \underbrace{(-Lx)_i}_{\sum_{j \in \mathcal{N}^{\text{out}}(i)} (x_i - x_j)}$$



- Agent i receives information only from its out neighbors
 - ▶ partial communication, distributed computations
- For $t \rightarrow +\infty$, do we have $x(t) \rightarrow$ **consensus state**?
 - ▶ if yes, **when average consensus** is achieved?
 - ▶ emergent behavior!

A physical example: RC network



- $c = \begin{bmatrix} C_1 \\ \vdots \\ C_n \end{bmatrix}$ = injected current in the nodes

- $V = \begin{bmatrix} V_1 \\ \vdots \\ V_n \end{bmatrix}$ = voltages at the nodes

- We have seen that $\dot{c} = -LV$, where $L =$ Laplacian of the coupling graph with weights $a_{ij} = \frac{1}{r_{ij}}$
- If C_1, \dots, C_n are the capacitances, then

$$C_i \dot{V}_i = c_i$$

The collective model, (for $C_i = 1$, $i = 1, \dots, n$)

$$\dot{V} = c \rightarrow \dot{V} = -LV$$

Problem: $\hat{=} V(t)$

- Will $x(t)$ converge to a consensus point $\alpha \mathbb{1}_n$ for some $\alpha \in \mathbb{R}$?
- Will $x(t)$ converge to average consensus, i.e. $\alpha = \langle x(0) \rangle$? ↓

Analysis of $\dot{x} = -Lx$

Lemma (equilibria)

If G contains a GRN, then all and only equilibria of $\dot{x} = -Lx$ are the states $\bar{x} = \alpha \mathbb{1}_n, \alpha \in \mathbb{R}$

Proof: $0 = -L\bar{x} \Leftrightarrow \bar{x}$ is the eigenvector with $\lambda = 0$. But

- $\alpha \mathbb{1}_n$ are eigenvectors of $\lambda = 0$
- $\lambda = 0$ is simple if G has a GRN
 - ▶ $\alpha \mathbb{1}_n$ are the only eigenvectors for $\lambda = 0$

} Theorem 1

Problem:

Is the set of equilibria attractive?

- Analysis of $x(t) = e^{-Lt}x(0)$

Analysis of $\dot{x} = -Lx \rightarrow x(t) = e^{-Lt}x(0)$

Theorem (Consensus with a GRN)

If G has a GRN, then

- 1 $\lim_{t \rightarrow +\infty} e^{-Lt} = \mathbb{1}_n w^T$, where w is the left eigenvector of L with $\lambda = 0$ verifying $w^T \mathbb{1}_n = 1$
- 2 $w_i > 0$ and $w_j > 0$ if and only if the node i is globally reachable
- 3 the solution $\dot{x}(t) = -Lx(t)$ verifies

$$\lim_{t \rightarrow +\infty} x(t) = (w^T x(0)) \mathbb{1}_n$$

- 4 if, in addition, G is balanced, then
 - (a) G is strongly connected,
 - (b) $\mathbb{1}_n^T L = [0 \dots 0]$,
 - (c) $w = \frac{1}{n} \mathbb{1}_n$, and
 - (d) $\lim_{t \rightarrow +\infty} x(t) = \left(\frac{1}{n} \mathbb{1}_n^T x(0) \right) \mathbb{1}_n$

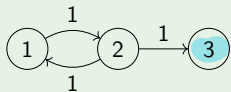
average of the initial conditions

Analysis of $\dot{x} = -Lx$

Remarks:

- (3) is consensus
- (4-d) is average consensus
- G is strongly connected \Leftrightarrow every node is globally reachable $\Leftrightarrow w_i > 0 \quad \forall i = 1, \dots, n$

Example (consensus with a leader)



$$\dot{x}_i = - \sum_{j \in \mathcal{N}^{\text{out}}(i)} (x_i - x_j)$$

From the theorem

- $w^T = [0 \ 0 \ w_3]$ $w_3 > 0$ because only node 3 is GR

- $x(t) \rightarrow (w^T x(0)) \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} = w_3 x_3(0) \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix}$

$\hookrightarrow w^T \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} = 1 \Rightarrow w_3 = 1$

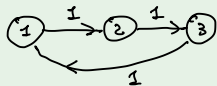
Consensus not influenced by agents 1 and 2

Check

$\dot{x} = -Lx$ gives

$$\begin{bmatrix} \dot{x}_1 \\ \dot{x}_2 \\ \dot{x}_3 \end{bmatrix} = - \begin{bmatrix} 1 & -1 & 0 \\ -1 & 2 & -1 \\ 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} \rightarrow \begin{cases} \dot{x}_1 = x_2 - x_1 \\ \dot{x}_2 = (x_3 - x_2) + (x_1 - x_2) \\ x_3(t) = x_3(0) \end{cases}$$

Example (average consensus)



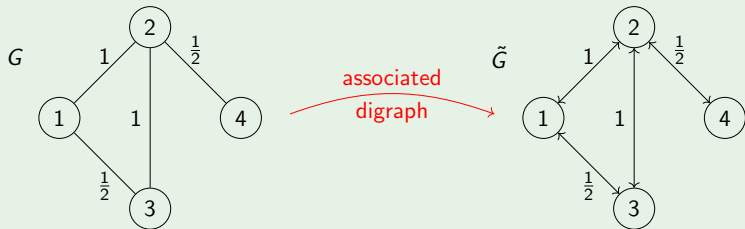
$$\dot{x}_i = - \sum_{j \in N^{\text{out}}(i)} (x_i - x_j)$$

From the Theorem

- G is strongly connected \rightarrow all nodes are GR $\rightarrow w > 0$
- G is balanced $\rightarrow w = \frac{1}{3} \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix}$

$$\hookrightarrow x(t) \rightarrow \left(\frac{1}{3} [1 \ 1 \ 1] \begin{bmatrix} x_1(0) \\ x_2(0) \\ x_3(0) \end{bmatrix} \right) \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} \text{ as } t \rightarrow \infty \rightarrow \text{AVERAGE CONSENSUS}$$

Example (undirected graph)



G connected $\Rightarrow \tilde{G}$ is strongly connected and **balanced**

For $\dot{x} = -Lx$ one has $x(t) \rightarrow \left(\frac{1}{n} \mathbb{1}_n^T x(0)\right) \mathbb{1}_n$, i.e. **average consensus**

Design of balanced digraphs

Problem

Given a strongly connected and weighted digraph $G = (V, E, a)$, how to re-define the weights in order to obtain $\tilde{G} = (V, E, \tilde{a})$ that is balanced?

From the previous Theorem, there is a left eigenvector of L (the Laplacian associated to G) with zero eigenvalue verifying $w^T \mathbb{1}_n = 1$ and $w \succ 0$. So we have

$$L \mathbb{1}_n = 0 \text{ and } w^T L = 0^T$$

Define $L_{\text{res}} = \text{diag}(w)L$

We have $L_{\text{res}} \mathbb{1}_n = \text{diag}(w)L \mathbb{1}_n = 0$ and

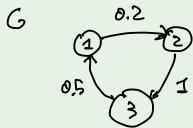
$$\mathbb{1}_n^T L_{\text{res}} = \mathbb{1}_n^T \text{diag}(w)L = [1 \quad \dots \quad 1] \begin{bmatrix} w_1 & & & \\ & \ddots & & \\ & & & w_n \end{bmatrix} L = w^T L = 0$$

Remark

- L_{res} is a Laplacian matrix (zero row sum, positive diagonal entries, non-positive non-diagonal entries)
- L_{res} is the Laplacian of \tilde{G} with weights $\tilde{a}_{ij} = w_i a_{ij}$
↪ By construction \tilde{G} is balanced!

because $w > 0$

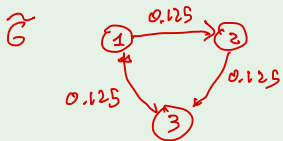
Example



$$L = D_{out} - A = \begin{bmatrix} 0.2 & -0.2 & 0 \\ 0 & 1 & -1 \\ -0.5 & 0 & 0.5 \end{bmatrix}$$

$$\text{ Matlab } \rightarrow w = \begin{bmatrix} 0.625 \\ 0.125 \\ 0.25 \end{bmatrix}$$

$$L_{res} = \text{diag}(w) L = \begin{bmatrix} 0.125 & -0.125 & 0 \\ 0 & 0.125 & -0.125 \\ -0.125 & 0 & 0.125 \end{bmatrix} \quad \text{which gives}$$

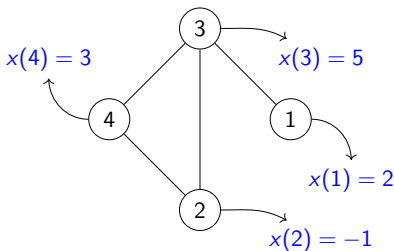


→ BALANCED

zh 413 ↓

Laplacian flow for undirected, connected graphs - intuition

- $x = [x_1, \dots, x_n]$ defines a function $x(v) \rightarrow \mathbb{R}$ on the node set V



Partial "derivative"

$$\partial_j x(i) = x(j) - x(i) \quad \text{if } j \in \mathcal{N}(i) \quad \partial_4 x(2) = x(4) - x(2)$$

Properties

- $\partial_j x(j) = 0$
 - $\partial_j x(i) = -\partial_i x(j)$
 - $\partial_j^2 x(i) = \underbrace{\partial_j x(j)}_{=0} - \partial_j x(i) = x(i) - x(j)$
- $\partial_3 (\partial_5 x(i)) = \partial_3 (x(i_5) - x(i_3)) = \partial_5 x(i_5) - \partial_5 x(i_3)$

Laplacian operators

In calculus

$$\Delta f(\xi) \stackrel{\text{def}}{=} \sum_{i=1}^n \frac{\partial^2 f(\xi)}{\partial \xi_i^2} \quad \text{for } f : \mathbb{R}^n \rightarrow \mathbb{R}$$

$\xi = \begin{bmatrix} \xi_1 \\ \vdots \\ \xi_n \end{bmatrix}$

On graphs

$$\Delta x(i) \stackrel{\text{def}}{=} - \sum_{j \in \mathcal{N}(i)} \partial_j^2(x(i)) = \sum_{j \in \mathcal{N}(i)} (x(j) - x(i))$$

$\partial_j^2 x(i) = -\partial_j x(i)$

L : Laplacian of the graph

Remark

$$\Delta x(i) = -(Lx)_i$$

Analogies

- In calculus $\Delta f(\xi) = 0$ if $f(\cdot)$ is constant
- On graphs $\Delta x(i) = 0$ if $x(\cdot)$ is constant (that is $x = \alpha \mathbb{1}_n$, $\alpha \in \mathbb{R}$)

- In calculus

$$\dot{f}(t, \xi) = \Delta f(t, \xi) \quad \text{THE HEAT EQUATION}$$

- ▶ evolution of the temperature f at time t and in point ξ of an isolated room
 $\Leftrightarrow f(t, \xi) \rightarrow \bar{f}(\xi)$ as $t \rightarrow +\infty$ where \bar{f} constant in space (heat diffusion). Moreover $\bar{f}(\xi) = \langle f(0, \xi) \rangle$, where $f(0, \xi)$ is the initial temperature in the point ξ

- On connected graphs

$$\dot{x}(t, v) = \Delta x(t, v) \Leftrightarrow \dot{x}(t) = -Lx(t)$$

gives $x(t) \rightarrow \bar{x}$ as $t \rightarrow +\infty$ where \bar{x} is constant over the graph nodes ("space") and $\bar{x} = \langle x(0) \rangle$

ALL ANALOGIES hold because Δ and $-L$ have a similar eigenstructure which qualifies $-L$ as a diffusion operator on graphs

Take home messages

- Graph Laplacians are the key for analyzing consensus algorithms in continuous time
 - ▶ continuous-time networks abstract real networks for very small sampling times
 - ▶ Laplacians naturally appear in physical models of electric and mechanical systems
- Consensus theorem for networks with a GRN
 - ▶ generalizations to time-varying graphs exist
- The Laplacian matrix is a diffusion operator on graphs