

## Serie 5

### Optimal transport, Fall semester

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**Exercise 5.1.** Let  $x_1, x_2, y_1, y_2 \in \mathbb{R}^d$ ,  $x_1 \neq x_2$  and let

$$\mu = \frac{1}{2}\delta_{x_1} + \frac{1}{2}\delta_{x_2} \quad \text{and} \quad \nu = \frac{1}{2}\delta_{y_1} + \frac{1}{2}\delta_{y_2}.$$

- (i) Describe all maps transporting  $\mu$  to  $\nu$ ; that is, such that  $T_{\#}\mu = \nu$ .
- (ii) Describe all couplings of  $\mu$  and  $\nu$ ; that is  $\gamma \in \mathcal{P}(X \times Y)$  such that  $(\pi_X)_{\#}\gamma = \mu$  and  $(\pi_Y)_{\#}\gamma = \nu$ .
- (iii) Prove that, for any choice of continuous cost  $c: \mathbb{R}^d \times \mathbb{R}^d \rightarrow \mathbb{R}$ , there exists an optimal transport map (i.e., the optimal coupling has a map structure).
- (iv) Assuming that  $x_1, Tx_1, x_2$  and  $Tx_2$  are not colinear, observe that for the linear cost  $c(x, y) = |x - y|$ , the corresponding optimal transport map does not cross trajectories.

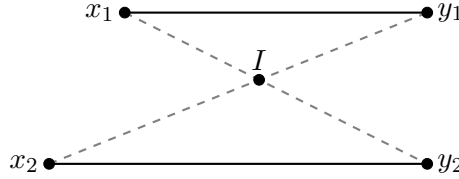


Figure 1: Crossing trajectories

#### **Solution:**

- (i) We have

$$\frac{1}{2}\varphi(y_1) + \frac{1}{2}\varphi(y_2) = \int \nu\varphi = \int T_{\#}\mu\varphi = \frac{1}{2}\varphi(Tx_1) + \frac{1}{2}\varphi(Tx_2)$$

for all  $\varphi \in C_c^\infty(\mathbb{R}^d)$ . Suppose now  $Tx_1 \notin \{y_1, y_2\}$ ; by taking  $\varphi$  with  $\varphi(Tx_1) > 0$ ,  $\varphi(y_1) = \varphi(y_2) = 0$  and  $\varphi(Tx_2) = 0$  if  $Tx_2 \neq Tx_1$ , we reach a contradiction. Therefore,  $Tx_1 \in \{y_1, y_2\}$ . Then, we have  $\varphi(\bar{y}) = \varphi(Tx_2)$  for some  $\bar{y} \in \{y_1, y_2\}$  and for all  $\varphi$ , so  $\bar{y} = Tx_2$ . Thus, we have two possibilities:

$$\begin{cases} Tx_1 = y_1 \\ Tx_2 = y_2 \end{cases} \quad \text{or} \quad \begin{cases} Tx_1 = y_2 \\ Tx_2 = y_1 \end{cases}.$$

- (ii) Notice that  $\text{supp } \gamma \subseteq \bigcup_{i,j} (x_i, y_j)$ . Indeed, let  $(x, y) \notin \bigcup_{i,j} (x_i, y_j)$ . Assume without loss of generality that  $x$  is different from  $x_1$  and  $x_2$ . There is a neighborhood  $N$  of  $x$  containing

neither  $x_1$  nor  $x_2$  and hence  $\gamma(N \times \mathbb{R}^d) = 0$ . This proves that  $(x, y) \in (\mathbb{R}^d \times \mathbb{R}^d) \setminus \text{supp } \gamma$ . In particular,

$$\gamma = \alpha_{11}\delta_{x_1y_1} + \alpha_{12}\delta_{x_1y_2} + \alpha_{21}\delta_{x_2y_1} + \alpha_{22}\delta_{x_2y_2},$$

and  $\gamma$  can be represented by the matrix

$$A_\gamma = \begin{pmatrix} \alpha_{11} & \alpha_{12} \\ \alpha_{21} & \alpha_{22} \end{pmatrix}.$$

Notice that we must have  $(\pi_X)_\# \gamma = \mu$ , so since  $(\pi_X)_\# \gamma = (\alpha_{11} + \alpha_{12})\delta_{x_1} + (\alpha_{21} + \alpha_{22})\delta_{x_2} = \mu$  this implies

$$\begin{cases} \alpha_{11} + \alpha_{12} = 1/2 \\ \alpha_{21} + \alpha_{22} = 1/2. \end{cases}$$

Similarly, since we have  $(\pi_Y)_\# \gamma = \nu$ ,

$$\begin{cases} \alpha_{11} + \alpha_{21} = 1/2 \\ \alpha_{12} + \alpha_{22} = 1/2. \end{cases}$$

In particular, we can take  $\alpha_{11} = \alpha$ , and then  $\alpha_{12} = \alpha_{21} = 1/2 - \alpha$  and  $\alpha_{22} = \alpha$ , that is

$$A_\gamma = \begin{pmatrix} \alpha & \frac{1}{2} - \alpha \\ \frac{1}{2} - \alpha & \alpha \end{pmatrix} \quad \text{for some } \alpha \in [0, 1/2],$$

and all such possibilities of  $A_\gamma$  describe a coupling between  $\mu$  and  $\nu$ .

(iii) Let us consider

$$\min_{\gamma \in \Gamma(\mu, \nu)} \int_{\mathbb{R}^d \times \mathbb{R}^d} c(x, y) d\gamma(x, y).$$

Notice that it is well-defined, because  $\gamma$  is supported on  $\bigcup_{i,j} (x_i, y_j)$  and  $c$  is bounded from below on this finite set. Notice also that by the previous part, the set of couplings is given by  $\Gamma(\mu, \nu) = \bigcup_{\alpha \in [0, 1/2]} \gamma_\alpha$ , with  $\gamma_\alpha$  represented by the matrix

$$A_{\gamma_\alpha} = \begin{pmatrix} \alpha & \frac{1}{2} - \alpha \\ \frac{1}{2} - \alpha & \alpha \end{pmatrix}, \quad \gamma_\alpha = \alpha\delta_{x_1y_1} + \left(\frac{1}{2} - \alpha\right)\delta_{x_1y_2} + \left(\frac{1}{2} - \alpha\right)\delta_{x_2y_1} + \alpha\delta_{x_2y_2}.$$

Thus, we want

$$\min_{\alpha \in [0, 1/2]} \int_{\mathbb{R}^d \times \mathbb{R}^d} c(x, y) d\gamma_\alpha(x, y).$$

Let us split

$$A_\alpha := A_{\gamma_\alpha} = 2\alpha A_{1/2} + (1 - 2\alpha)A_0, \quad \text{so} \quad \gamma_\alpha = 2\alpha\gamma_{1/2} + (1 - 2\alpha)\gamma_0,$$

and

$$\text{cost}(\gamma_\alpha) = \int_{\mathbb{R}^d \times \mathbb{R}^d} c(x, y) d\gamma_\alpha(x, y) = 2\alpha \text{cost}(\gamma_{1/2}) + (1 - 2\alpha) \text{cost}(\gamma_0).$$

In particular, since  $2\alpha \in [0, 1]$ ,  $\text{cost}(\gamma_\alpha) \geq \min\{\text{cost}(\gamma_0), \text{cost}(\gamma_{1/2})\}$  for all  $\alpha \in [0, 1/2]$  and

$\gamma_0$  or  $\gamma_{1/2}$  is an optimal coupling, which has a map structure ( $T_1$  or  $T_2$  from (i) ).

- (iv) We need to show that the segments  $\overline{x_1 T x_1}$  and  $\overline{x_2 T x_2}$  do not cross. Since  $\{T x_1, T x_2\} = \{y_1, y_2\}$  if they were to cross,  $x_1, x_2, y_1, y_2$  would be coplanar, thus it is enough to study the case  $d = 2$ . Notice that the cost to bring  $x_1$  to  $T x_1$  is simply  $|x_1 - T x_1|$ , thus the result follows by the triangular inequality. Indeed, if the trajectories were to cross at a point  $I$  (see figure 1 below), then have

$$|x_1 - y_1| + |x_2 - y_2| \leq |x_1 - I| + |I - y_1| + |x_2 - I| + |I - y_2| = |x_1 - y_2| + |x_2 - y_1|$$

where equality holds only if  $x_1, x_2, y_1$  and  $y_2$  are all colinear, a contradiction. This implies that the trajectories do not cross.

**Exercise 5.2.** Say if the following sentences are true or false. If they are true, prove it, if they are false, provide a counterexample. The statements below all refer to the quadratic cost.

- (i) Let  $\varphi : \mathbb{R}^d \rightarrow \mathbb{R}$  be a convex function, then  $\varphi$  is differentiable  $\mathcal{L}^d$ -a.e. in  $\mathbb{R}^d$  and we call  $N \subset \mathbb{R}^d$  the Lebesgue measure zero set where  $\varphi$  is not differentiable. For any  $x \in N$  take an element  $y_x \in \partial\varphi(x)$ , and define the map  $T : \mathbb{R}^d \rightarrow \mathbb{R}^d$  as follows:

$$T(x) = \begin{cases} \nabla\varphi(x) & \text{if } x \in \mathbb{R}^d \setminus N, \\ y_x & \text{if } x \in N. \end{cases}$$

Then, given  $\mu \ll \mathcal{L}^d$ , the map  $T$  is optimal from  $\mu$  to  $T_{\#}\mu$ .

- (ii) If  $T : \mathbb{R} \rightarrow \mathbb{R}$  is an optimal map between  $\mu_1$  and  $\mu_2$  (i.e.  $T_{\#}\mu_1 = \mu_2$ ) and  $S : \mathbb{R} \rightarrow \mathbb{R}$  is optimal between  $\mu_2$  and  $\mu_3$ , then  $S \circ T$  is optimal between  $\mu_1$  and  $\mu_3$ .
- (iii) The same as before but in general dimension  $d \geq 2$ , namely: if  $T : \mathbb{R}^d \rightarrow \mathbb{R}^d$  is an optimal map between  $\mu_1$  and  $\mu_2$  (i.e.  $T_{\#}\mu_1 = \mu_2$ ) and  $S : \mathbb{R}^d \rightarrow \mathbb{R}^d$  is optimal between  $\mu_2$  and  $\mu_3$ , then  $S \circ T$  is optimal between  $\mu_1$  and  $\mu_3$ .

**Solution:**

- (i) True. The  $\mathcal{L}^d$ -a.e. differentiability follows from the Rademacher Theorem. Then if we define  $\gamma := (Id, T)_{\#}\mu$ , then by construction we have that  $\text{supp}\gamma \subset \partial\varphi$ , from which thanks to Theorem 2.4.3 (see also Remark 2.4.4) we conclude that  $\gamma$  is optimal because it is contained in a c-cyclically monotone set.

Notice that if we defined for instance, for all  $x \in N$ ,  $T(x) = y$  for  $y \notin \partial\varphi$  we should have proved also that  $\text{supp}(\gamma) = \text{supp}(\gamma) \cap (\mathbb{R}^d \setminus N) \times \mathbb{R}^d \subset \partial\varphi$  where the first equality is not trivial.

- (ii) It is true provided that  $\mu_3$  is non-atomic.

Suppose that  $\mu_3$  is non-atomic. Then, since there exist a transport map from  $\mu_2$  to  $\mu_3$ ,  $\mu_2$  is non-atomic too. With the same argument we get that  $\mu_1$  is non-atomic. As  $T$  and  $S$  are optimal maps, they have to coincide with the monotone rearrangements from  $\mu_1$  to  $\mu_2$  and from  $\mu_2$  to  $\mu_3$  respectively. In particular,  $S \circ T$  is monotone, and of course is a transport map from  $\mu_1$  to  $\mu_3$ . This means that  $S \circ T$  is optimal from  $\mu_1$  to  $\mu_3$ .

Now we provide a counterexample (suggested by Berk Ceylan) in the case in which  $\mu_3$  has atoms. Take

$$\begin{aligned}\mu_1 &= \frac{1}{6}\delta_1 + \frac{1}{6}\delta_2 + \frac{1}{6}\delta_3 + \frac{1}{6}\delta_4 + \frac{1}{6}\delta_5 + \frac{1}{6}\delta_6, \\ \mu_2 &= \frac{1}{3}\delta_1 + \frac{1}{3}\delta_2 + \frac{1}{6}\delta_5 + \frac{1}{6}\delta_6, \\ \mu_3 &= \frac{1}{2}\delta_1 + \frac{1}{2}\delta_6.\end{aligned}$$

It can be verified that the following maps are optimal for the Monge problem from  $\mu_1$  to  $\mu_2$  and from  $\mu_2$  to  $\mu_3$  respectively:

$$T(x) = \begin{cases} 1 & \text{if } x = 1, 2, \\ 2 & \text{if } x = 3, 4, \\ 5 & \text{if } x = 5, \\ 6 & \text{if } x = 6. \end{cases}, \quad S(x) = \begin{cases} 1 & \text{if } x = 1, 5, \\ 6 & \text{if } x = 2, 6. \end{cases}$$

However, the map

$$S \circ T(x) = \begin{cases} 1 & \text{if } x = 1, 2, 5, \\ 6 & \text{if } x = 3, 4, 6 \end{cases}$$

is clearly non optimal from  $\mu_1$  to  $\mu_3$ . The problem here is that when there are atoms, in general Monge's optimal maps are not minimizers for the Kantorovich problem. In particular, optimal maps can be non monotone.

- (iii) False. We provide a counterexample for Dirac deltas, the exercise can be generalized to absolutely continuous measures.

Let  $\mu_1 = \frac{\delta_{(1,0)} + \delta_{(-1,0)}}{2}$ ,  $\mu_2 = \frac{\delta_{(-1,0)} + \delta_{(1,4)}}{2}$ ,  $\mu_3 = \frac{\delta_{(-1,4)} + \delta_{(1,0)}}{2}$ . Using point (i) of Exercise 5.1 we can directly prove that  $T_1$  defined as

$$T_1(-1,0) = (-1,0) \quad T_1(1,0) = (1,4)$$

and extended in whatever way is an optimal map from  $\mu_1$  to  $\mu_2$ , as well as

$$T_2(-1,0) = (1,0) \quad T_2(1,4) = (-1,4)$$

and extended in whatever way is an optimal map from  $\mu_2$  and  $\mu_3$ , but  $T_2 \circ T_1$  is not an optimal map from  $\mu_1$  to  $\mu_3$  because

$$T_2 \circ T_1(-1,0) = (1,0) \quad T_2 \circ T_1(1,0) = (-1,4),$$

indeed the optimal map is  $T_3$  from  $\mu_1$  to  $\mu_3$  is defined as

$$T_3(1, 0) = (1, 0) \quad T_3(-1, 0) = (-1, 4)$$

and extended in whatever way.

**Exercise 5.3** (Birkhoff - Von Neumann Theorem). A  $(n \times n)$ -matrix  $A \in \mathcal{M}(n, \mathbb{R})$  with nonnegative entries is said to be:

- a *doubly-stochastic matrix* if  $\sum_{i=1}^n A_{ij} = 1$  for any  $j = 1, \dots, n$ , and  $\sum_{j=1}^n A_{ij} = 1$  for any  $i = 1, \dots, n$ .
- a *permutation matrix* if there is a permutation  $\sigma : \{1, \dots, n\} \rightarrow \{1, \dots, n\}$  such that  $A_{i\sigma(i)} = 1$  and  $A_{ij} = 0$  if  $j \neq \sigma(i)$ .

Prove that any doubly-stochastic matrix can be written as a finite convex combination of permutation matrices.

*Hints:* Here is a guideline through a possible proof of the result:

- Use Hall's marriage Theorem<sup>1</sup> to prove that given a doubly-stochastic matrix  $A$ , there exists a permutation  $\sigma \in S_n$  such that  $A_{i\sigma(i)} > 0$  for any  $i = 1, \dots, n$ . Deduce that there exists a permutation matrix  $P$  and  $\lambda > 0$  such that  $A_{ij} \geq \lambda P_{ij}$ ,  $\forall i, j \in \{1, \dots, n\}$ .
- Let us now prove the result by induction on the number of non-zero entries  $k$  of  $A$ . Start by proving that  $k \geq n$  and that the result holds for  $k = n$ .
- Let now  $k > n$ . Consider the permutation  $P$  and  $\lambda$  given in the first bullet above, and define

$$A' = \frac{1}{1 - \lambda}(A - \lambda P).$$

Show that  $A'$  is doubly-stochastic with at most  $k - 1$  non-zero entries.

- Deduce, by induction, that  $A$  is a convex combination of permutation matrices.

**Solution:** The solution is decomposed into four steps based on the hints.

*Step 1: (Application of Hall's marriage theorem)* Let us begin with the following lemma.

**Lemma 1.** Given a doubly-stochastic matrix  $A$ , there is a permutation  $\sigma \in S_n$  such that  $A_{i\sigma(i)} > 0$  for any  $i = 1, \dots, n$ .

*Proof.* Let us construct a bipartite graph as follows: the graph consists of  $2n$  vertices labeled by  $\{1_r, \dots, n_r\}$  and  $\{1_c, \dots, n_c\}$  (the indexes  $r, c$  stand for row and column). Then, we say that there is an edge between  $i_r$  and  $j_c$  if and only if  $A_{ij} > 0$ . We denote the presence of an edge with  $i_r \sim j_c$ . The first step of the proof consists in showing that such a bipartite graph admits a

<sup>1</sup>See [https://en.wikipedia.org/wiki/Hall%27s\\_marriage\\_theorem#Graph\\_theoretic\\_formulation](https://en.wikipedia.org/wiki/Hall%27s_marriage_theorem#Graph_theoretic_formulation).

perfect matching (i.e., there is a permutation  $\sigma : \{1, \dots, n\} \rightarrow \{1, \dots, n\}$  such that  $i_r \sim \sigma(i)_c$  for any  $i = 1, \dots, n$ ). In order to do so, we want to apply Hall's marriage theorem. Given a subset  $S \subset \{1, \dots, n\}$ , let  $T$  be the subset defined as

$$T = \{t \in \{1, \dots, n\} : s_r \sim t_c \text{ for at least one } s \in S\}$$

Exploiting the fact that the matrix  $A$  is doubly-stochastic and the definition of  $T$ , we obtain

$$\#S = \sum_{s \in S} \sum_{j=1}^n A_{sj} = \sum_{s \in S} \sum_{t \in T} A_{st} \leq \sum_{i=1}^n \sum_{t \in T} A_{it} = \sum_{t \in T} \sum_{i=1}^n A_{it} = \#T$$

Since we can choose  $S$  arbitrarily, the inequality  $\#S \leq \#T$  is exactly the hypothesis necessary to apply Hall's marriage theorem and deduce the existence of a perfect matching. Hence, by definition of perfect matching, there is a permutation  $\sigma$  such that  $i_r \sim \sigma(i)_c$  for any  $i = 1, \dots, n$ . This last fact is equivalent to the desired statement.  $\square$

*Step 2:* We can now prove the statement of the theorem by induction on the number of nonzero entries of the matrix  $A$ .

Since  $A$  is doubly-stochastic, it is easy to see that it must have at least  $n$  nonzero entries. Moreover, if it has exactly  $n$  nonzero entries, then it must be already a permutation matrix.

*Step 3:* Let us assume that the number of nonzero entries of  $A$  is  $k > n$ . Let  $P^\sigma$  be the permutation matrix induced by the permutation  $\sigma$  (that is,  $P_{i\sigma(i)}^\sigma = 1$  for all  $i$ , and  $P_{ij}^\sigma = 0$  if  $j \neq \sigma(i)$ ) whose existence is provided by the lemma. Let  $\lambda > 0$  be the maximum value such that  $\lambda P^\sigma \leq A$  (the inequality must be understood entry-wise, namely  $\lambda P_{ij}^\sigma \leq A_{ij}$  for all  $i, j$ ). Notice that, since  $A$  is doubly-stochastic, each entry of  $A$  is bounded by 1 and therefore  $\lambda \leq 1$ . Also, it must be  $\lambda < 1$ , as otherwise  $A$  would have exactly  $n$  nonzero entries.

Let  $A' := \frac{1}{1-\lambda} (A - \lambda P^\sigma)$ . Since  $\lambda P^\sigma \leq A$ , all entries of  $A'$  are nonnegative. Moreover, thanks to the choice of  $\lambda$ , the matrix  $A'$  has at most  $k - 1$  nonzero entries. Finally, one can easily check that  $A'$  is doubly-stochastic.

*Step 4:* By the inductive hypothesis, we are able to write  $A'$  as a convex combination of permutation matrices

$$A' = \sum_{i \in I} \lambda_i P^{\sigma_i}, \quad \lambda_i \geq 0, \quad \sum_{i \in I} \lambda_i = 1,$$

where  $I$  is a finite set of indices and  $P^{\sigma_i}$  are permutation matrices (induced by the permutations  $\sigma_i$ ). From the definition of  $A'$ , it follows that

$$A = \lambda P^\sigma + \sum_{i \in I} \lambda_i (1 - \lambda) P^{\sigma_i},$$

thus  $A$  is a convex combination of permutation matrices.

**Exercise 5.4** (Discrete optimal transport). Given two families  $\{x_1, \dots, x_n\}$  and  $\{y_1, \dots, y_n\}$  of points in  $\mathbb{R}^d$ , let  $\mu := \frac{1}{n} \sum_{i=1}^n \delta_{x_i}$  and  $\nu := \frac{1}{n} \sum_{i=1}^n \delta_{y_i}$ . Prove that, for *any* choice of a continuous cost  $c : \mathbb{R}^d \times \mathbb{R}^d \rightarrow \mathbb{R}$ , there exists an optimal transport map from  $\mu$  to  $\nu$ .

*Hint:* Use Exercise 5.3 or Kantorovich duality.

**Solution:** We present two different solutions of this exercise, the first one uses Birkhoff-Von Neumann's theorem, whereas the second one borrows some ideas from the duality theory.

Note that  $c$  is bounded below on the finite set  $\{(x_i, y_j)\}_{1 \leq i, j \leq n}$ , so it follows by Theorem 2.3.2 and Remark 2.3.3 that there exists an optimal coupling  $\gamma \in \Gamma(\mu, \nu)$ .

Let  $A_\gamma$  be the  $n \times n$  matrix defined as  $A_{ij} := \gamma(\{x_i, y_j\})$ . From the marginal constraints on  $\gamma$ , it follows that  $nA$  is a doubly-stochastic matrix. Hence, applying Exercise 5.3, we can express  $nA$  as a convex combination of permutation matrices

$$nA = \sum_{k \in I} \lambda_k P^{\sigma_k},$$

where  $I$  is a finite set of indices,  $\sum_{k \in I} \lambda_k = 1$ , and  $P^{\sigma_k}$  are permutation matrices (induced by the permutations  $\sigma_k$ ).

Let us define the cost of an  $n \times n$  matrix  $B$  as

$$\mathcal{C}(B) := \sum_{i,j=1}^n B_{ij} c(x_i, y_j).$$

By definition, the cost  $\mathcal{C}$  is linear and it holds

$$\int_{\mathbb{R}^d \times \mathbb{R}^d} c(x, y) d\gamma(x, y) = \mathcal{C}(A) = \frac{1}{n} \sum_{k \in I} \lambda_k \mathcal{C}(P^{\sigma_k}) \geq \frac{1}{n} \min_{k \in I} \mathcal{C}(P^{\sigma_k}).$$

Hence, there is a permutation  $\sigma_k$  such that

$$\frac{1}{n} \sum_{i=1}^n c(x_i, y_{\sigma_k(i)}) = \frac{1}{n} \mathcal{C}(P^{\sigma_k}) \leq \int_{\mathbb{R}^d \times \mathbb{R}^d} c(x, y) d\gamma(x, y),$$

and therefore the map  $T : \mathbb{R}^d \rightarrow \mathbb{R}^d$  such that  $T(x_i) = y_{\sigma_k(i)}$  is optimal.

**Solution:** Without loss of generality (up to relabeling the indices of the points  $y_i$ ), we can assume that the trivial permutation has the minimum cost among all permutations, that is

$$\sum_{i=1}^n c(x_i, y_i) \leq \sum_{i=1}^n c(x_i, y_{\sigma(i)}) \quad (1)$$

for any permutation  $\sigma : \{1, \dots, n\} \rightarrow \{1, \dots, n\}$ . Under this assumption, we want to prove that the map  $T(x_i) := y_i$  is optimal, in the sense that the coupling induced by it is optimal (in the Kantorovich sense).

We now want to use Theorem 2.6.6. In fact, we only need to use the inequality

$$\inf_{\gamma \in \Gamma(\mu, \nu)} \int_{X \times Y} c d\gamma \geq \sup_{\varphi(x) + \psi(y) + c(x, y) \geq 0} \int -\varphi d\mu + \int -\psi d\nu,$$

which follows immediately from the marginal condition (see the proof of (2.12)). In any case, it

suffices to construct two functions  $\varphi : \{x_1, \dots, x_n\} \rightarrow \mathbb{R}$  and  $\psi : \{y_1, \dots, y_n\} \rightarrow \mathbb{R}$  such that

$$\varphi(x_i) + \psi(y_j) + c(x_i, y_j) \geq 0 \quad \text{for any } 1 \leq i, j \leq n, \quad (2)$$

$$\sum_{i=1}^n c(x_i, y_i) = \sum_{i=1}^n -\varphi(x_i) - \psi(y_i). \quad (3)$$

Indeed, from these equations we get

$$\int_{\mathbb{R}^d} c(x, T(x)) d\mu = \frac{1}{n} \sum_{i=1}^n c(x_i, y_i) \stackrel{(3)}{=} \int_{\mathbb{R}^d} -\varphi d\mu + \int_{\mathbb{R}^d} -\psi d\nu \stackrel{(2)}{\leq} \inf_{\gamma \in \Gamma(\mu, \nu)} \int_{\mathbb{R}^d \times \mathbb{R}^d} c d\gamma,$$

so the optimality of  $T$  follows.

To prove (2)-(3), we claim that it suffices to construct a function  $\varphi$  such that

$$\varphi(x_j) - \varphi(x_i) \leq c(x_i, y_j) - c(x_j, y_j) =: b_{ij} \quad \text{for any } 1 \leq i, j \leq n. \quad (4)$$

Indeed, if the above bound holds, then (2)-(3) hold with the function  $\psi$  defined as  $\psi(y_i) := -c(x_i, y_i) - \varphi(x_i)$  for any  $1 \leq i \leq n$ .

So, it remains only to construct a function  $\varphi$  such that (4) holds. To do this, let us consider the weighted oriented complete graph with vertices  $\{1, \dots, n\}$  such that the weight of the edge  $i \rightarrow j$  is  $b_{ij}$ , and denote with  $d(i, j)$  the distance (i.e. the infimum of the sum of the weights of a path from the first vertex to the second one) between vertex  $i$  and vertex  $j$  (notice the similarity between this approach and the proof of Rockafellar's theorem, Theorem 2.5.2).

Let us check that, for any  $1 \leq i, j \leq n$ , it holds  $d(i, j) > -\infty$ . Since the graph consists of finitely many points, one can note that the distance between two vertices can be  $-\infty$  if and only if there is a simple loop (that is, a closed path that visits each vertex at most once)  $i_1, i_2, \dots, i_k$  with negative length, that is

$$b_{i_1 i_2} + b_{i_2 i_3} + \dots + b_{i_k i_1} < 0. \quad (5)$$

To rule out this possibility, we have to use the optimality condition (1) (that we have never used until now). Let  $\bar{\sigma}$  be the permutation such that  $\bar{\sigma}(i_1) = i_2, \bar{\sigma}(i_2) = i_3, \dots, \bar{\sigma}(i_k) = i_1$ , and  $\bar{\sigma}(i) = i$  for all other values of  $i$ . Applying 1 with  $\sigma = \bar{\sigma}$ , we get

$$0 \leq \sum_{i=1}^n c(x_i, y_{\bar{\sigma}(i)}) - c(x_{\bar{\sigma}(i)}, y_{\bar{\sigma}(i)}) = \sum_{i=1}^n b_{i\bar{\sigma}(i)} = b_{i_1 i_2} + b_{i_2 i_3} + \dots + b_{i_k i_1},$$

which shows that (5) cannot hold.

Hence, we have proven that the distance  $d(i, j)$  is finite for every  $1 \leq i, j \leq n$ . We now observe that, even if this notion of distance on a graph might be negative and not symmetric, it still satisfies the triangle inequality. Therefore we have

$$d(1, j) \leq d(1, i) + d(i, j) \leq d(1, i) + b_{ij} \quad \forall i, j.$$

Hence, if we set  $\varphi(x_i) := d(1, i)$  then the desired inequality (4) holds, concluding the proof.