

Solutions for Statistical analysis of network data – Sheet 11

Here, 1 correspond to graph (a), 2 to (b) and 3 to (c). We list a few definitions that are helpful for any English longbowman that does not practice on Sundays:

- The *order* of a hypergraph is the number of nodes it has.
- In hypergraphs, if you have a hyperedge that connects a set of vertices, and another hyperedge that connects a subset of these vertices, the latter is said to be "*included*" in the former.
- A *multiple* edge occurs when there are two or more hyperedges connecting the same set of vertices.
- A hypergraph is *simple* if it has no multiple and induced edges.
- Nodes are *adjacent* if they are part of the same edge.
- The *degree* of a vertex is the cardinality of the set of edges it's belonging to.
- The *size* of an edge is its cardinality, i.e. its number of nodes.
- A hyper graph is *regular* if all its nodes have the same degree.
- A hyper graph is *uniform* if all its edges have the same cardinality.
- The *rank* of a hypergraph is the biggest edge cardinality in the hypergraph.

- Order = 6, $n_{\text{edges}} = 11$.
 - Included edges are edges that are a strict subset of another edge. Thus, edge $\{2, 6\}$ in edge $\{2, 4, 6\}$ (top circle); edges $\{3, 4\}$ and $\{5, 6\}$ in edge $\{3, 4, 5, 6\}$ (right-hand side circle); edge $\{1, 5\}$ in edge $\{1, 3, 5\}$ (bottom circle); edges $\{3, 4\}$ and $\{1, 2\}$ in edge $\{1, 2, 3, 4\}$ (lefthand side circle).
 - A multiple edge occurs when there are two or more hyperedges connecting the same set of vertices. Node groups that are encircled, that is, $\{2, 4, 6\}$ (top circle), $\{3, 4, 5, 6\}$ (right-hand side circle), $\{1, 3, 5\}$ (bottom circle), and $\{1, 2, 3, 4\}$ (left-hand side circle).
 - Graph (a) is not a simple graph: it has included edges as seen above as well as multiple edges.
 - Not adjacent: 1 and 6, 2 and 5. All other pairs are adjacent.
 - Vertex 1: $d_1 = 5$, neighbourhood: $\{2, 3, 4, 5\}$
 - Vertex 2: $d_2 = 5$, neighbourhood: $\{1, 3, 4, 6\}$
 - Vertex 3: $d_3 = 5$, neighbourhood: $\{1, 2, 4, 5, 6\}$
 - Vertex 4: $d_4 = 5$, neighbourhood: $\{1, 2, 3, 5, 6\}$
 - Vertex 5: $d_5 = 4$, neighbourhood: $\{1, 3, 4, 6\}$
 - Vertex 6: $d_6 = 4$, neighbourhood: $\{2, 3, 4, 5\}$
 - Edge list: $\{1, 2\}, \{1, 3\}, \{1, 5\}, \{2, 4\}, \{2, 6\}, \{3, 4\}, \{5, 6\}, \{2, 4, 6\}, \{1, 3, 5\}, \{3, 4, 5, 6\}$, and $\{1, 2, 3, 4\}$. The last two have cardinality 4, the two before them have cardinality 3, the others have cardinality 2.
 - $\Delta(\mathcal{H}_a) = 5$.
 - Graph (a) is not regular (not all degrees are equal).
 - Graph (a) is not uniform (not all edges have the same cardinality).
 - The rank of graph correspond to the biggest cardinality of an edge in the hypergraph. Here, in (a), it is $r(\mathcal{H}_a) = 4$.
 - Isolated nodes cannot be found in either of the graphs (a), (b) and (c). Graph (c) contains a pendant node, the other two does not.
 - No singleton or empty edges can be found in any of the three graphs.
- Order = 5, $n_{\text{edges}} = 8$.

- (b) Edge $\{1, 2\}$ in edge $\{1, 2, 3\}$ and $\{1, 2, 5\}$; $\{1, 3\}$ and $\{2, 3\}$ are in edge $\{1, 2, 3\}$, $\{3, 4\}$ and $\{4, 5\}$ are in $\{3, 4, 5\}$
- (c) Node groups that are encircled.
- (d) Graph (a) is not a simple graph: it has included edges as seen above as well as multiple edges.
- (e) Not adjacent: 1 and 4, 2 and 4. All other pairs are adjacent.
- (f)
- Vertex 1: $d_1 = 4$, neighbourhood: $\{2, 3, 5\}$
 - Vertex 2: $d_2 = 4$, neighbourhood: $\{1, 3, 5\}$
 - Vertex 3: $d_3 = 5$, neighbourhood: $\{1, 2, 4, 5\}$
 - Vertex 4: $d_4 = 3$, neighbourhood: $\{3, 5\}$
 - Vertex 5: $d_5 = 3$, neighbourhood: $\{1, 2, 3, 4\}$
- (g) Take the cardinality (i.e. number of nodes of question (b)).
- (h) $\Delta(\mathcal{H}_b) = 5$.
- (i) Graph (b) is not regular (not all degrees are equal).
- (j) Graph (b) is not uniform (not all edges have the same cardinality).
- (k) The rank of graph (b) is $r(\mathcal{H}_a) = 3$.
- (l) Isolated nodes cannot be found in either of the graphs (a), (b) and (c). Graph (c) contains a pendant node, the other two does not.
- (m) No singleton or empty edges can be found in any of the three graphs.
3. (a) Order = 4, $n_{\text{edges}} = 4$.
- (b) Edges $\{2, 3\}$ and $\{3, 4\}$ in edge $\{2, 3, 4\}$.
- (c) Node groups that are encircled.
- (d) Graph (c) is not a simple graph: it has included edges as seen above as well as multiple edges.
- (e) Not adjacent: 1 and 4. All other pairs are adjacent.
- (f)
- Vertex 1: $d_1 = 1$, neighbourhood: $\{2\}$
 - Vertex 2: $d_2 = 3$, neighbourhood: $\{1, 3, 4\}$
 - Vertex 3: $d_3 = 3$, neighbourhood: $\{2, 4\}$
 - Vertex 4: $d_4 = 2$, neighbourhood: $\{2, 3\}$
- (g) Edge list: $\{1, 2\}, \{2, 3\}, \{3, 4\}, \{2, 3, 4\}$ and the cardinality of each.
- (h) $\Delta(\mathcal{H}_c) = 3$.
- (i) Graph (c) is not regular (not all degrees are equal).
- (j) Graph (c) is not uniform (not all edges have the same cardinality).
- (k) The rank of graph (c) is $r(\mathcal{H}_a) = 3$.
- (l) Isolated nodes cannot be found in either of the graphs (a), (b) and (c). Graph (c) contains a pendant node, the other two does not.
- (m) No singleton or empty edges can be found in any of the three graphs.
- (n) See Figure 1. for the two graphs. The edge set of the first is $e_1 = v_1v_2, e_2 = v_2v_3, e_3 = v_3v_4, e_4 = v_2v_3v_4$. The edge set of the second is $f_1 = w_1w_2, f_2 = w_2w_3, f_3 = w_3w_4, f_4 = w_1w_2w_3$. An isomorphism between the two is $\Phi(v_1) = w_4, \Phi(v_2) = w_3, \Phi(v_3) = w_2, \Phi(v_4) = w_1$. Since $\Phi(e_1) = w_3w_4 = f_3, \Phi(e_2) = w_2w_3 = f_2, \Phi(e_3) = w_1w_2 = f_1, \Phi(e_4) = w_1w_2w_3 = f_4$, and this constitutes a bijection, the two hypergraphs are isomorphic.
- (o) The incidence matrix of the graph on the left-hand side of Figure1, using the same edge labeling as in point (xv), is

$$I = \begin{array}{c|cccc} & e_1 & e_2 & e_3 & e_4 \\ \hline v_1 & 1 & 0 & 0 & 0 \\ v_2 & 1 & 1 & 0 & 1 \\ v_3 & 0 & 1 & 1 & 1 \\ v_4 & 0 & 0 & 1 & 1 \end{array}$$

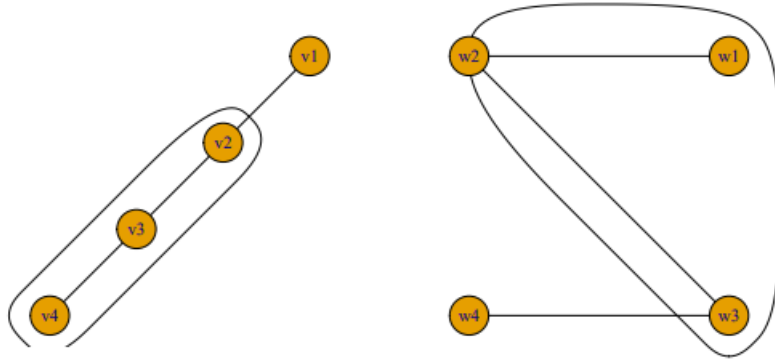


Figure 1: Isomorphism to different labelling.

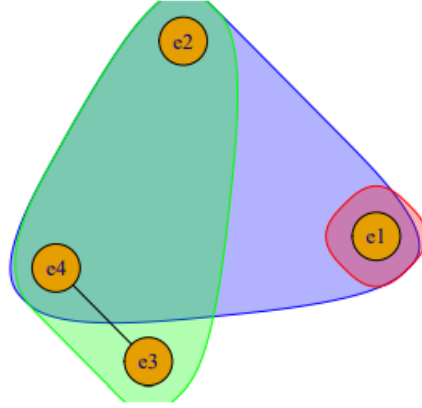


Figure 2: Dual of graph (c)

Its dual is the graph whose incidence matrix is the transpose of the above, that is,

$$I = \begin{array}{c|cccc} & v_1 & v_2 & v_3 & v_4 \\ \hline e_1 & 1 & 1 & 0 & 0 \\ e_2 & 0 & 1 & 1 & 0 \\ e_3 & 0 & 0 & 1 & 1 \\ e_4 & 0 & 1 & 1 & 1 \end{array}$$

The resulting graph is shown in Figure 2. The former edges now become vertices (see labels on vertices). There is an edge now with a single node (corresponding to the former node v_1 , which was a pendant.