## Mathematics of Data: From Theory to Computation

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### Lecture 7: Introduction to proximal-operators. Conditional gradient methods.

Laboratory for Information and Inference Systems (LIONS) École Polytechnique Fédérale de Lausanne (EPFL)

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### Outline

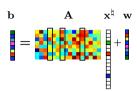
- Composite minimization
- ► Proximal gradient methods
- ► Introduction to Frank-Wolfe method

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# Recall sparse regression in generalized linear models (GLMs)

# Problem (Sparse regression in GLM)

Our goal is to estimate  $\mathbf{x}^{\natural} \in \mathbb{R}^p$  given  $\{b_i\}_{i=1}^n$  and  $\{\mathbf{a}_i\}_{i=1}^n$ , knowing that the likelihood function at  $y_i$  given  $\mathbf{a}_i$  and  $\mathbf{x}^{\natural}$  is given by  $L(\langle \mathbf{a}_i, \mathbf{x}^{\natural} \rangle, b_i)$ , and that  $\mathbf{x}^{\natural}$  is sparse.



### Optimization formulation

$$\min_{\mathbf{x} \in \mathbb{R}^p} \left\{ \underbrace{-\sum_{i=1}^n \log L(\langle \mathbf{a}_i, \mathbf{x}^{\natural} \rangle, b_i)}_{f(\mathbf{x})} + \underbrace{\rho_n \|\mathbf{x}\|_1}_{g(\mathbf{x})} \right\}$$

where  $\rho_n > 0$  is a parameter which controls the strength of sparsity regularization.

### Theorem (cf. [13] for details)

Under some technical conditions, there exists  $\{\rho_i\}_{i=1}^{\infty}$  such that with high probability, the following holds

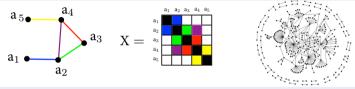
$$\parallel \mathbf{x}^{\star} - \mathbf{x}^{\natural} \parallel_2^2 = \mathcal{O}\left(\frac{s\log p}{n}\right), \quad \operatorname{supp} \mathbf{x}^{\star} = \operatorname{supp} \mathbf{x}^{\natural}.$$

$$\text{Recall ML: } \parallel \mathbf{x}_{\mathit{ML}} - \mathbf{x}^{\natural} \parallel_2^2 = \mathcal{O}\left(\frac{p}{n}\right).$$

### Sparse inverse covariance estimation

## Problem (Graphical model selection)

Given a data set  $\mathcal{D} := \{\mathbf{x}_1, \cdots, \mathbf{x}_n\}$ , where  $\mathbf{x}_i$  is a Gaussian random variable. Let  $\Sigma$  be the covariance matrix corresponding to the graphical model of the Gaussian Markov random field. Our goal is to learn a sparse precision matrix X (i.e., the inverse covariance matrix  $\Sigma^{-1}$ ) that captures the Markov random field structure.



# **Optimization formulation** [16]

$$\min_{\mathbf{X} \succ 0} \left\{ \underbrace{\operatorname{tr}(\Sigma \mathbf{X}) - \log \det(\mathbf{X})}_{f(\mathbf{x})} + \underbrace{\rho_n \|\operatorname{vec}(\mathbf{X})\|_1}_{g(\mathbf{x})} \right\}, \tag{1}$$

where  $X\succ 0$  means that X is symmetric and positive definite and  $\rho_n>0$  is a regularization parameter and vec is the vectorization operator. Let  $X^\star$  be the minimizer of (1), under some technical conditions, there exists a  $\rho_n$  such that  $\|X^\star - \Sigma^{-1}\|_2^2 = \mathcal{O}(\min\frac{1}{n}\{d^2\log p,(s+p)\log p\})$  where d is the maximum node degree.

### Composite convex minimization

## Problem (Composite convex minimization)

$$F^* := \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ F(\mathbf{x}) := f(\mathbf{x}) + g(\mathbf{x}) \right\}$$
 (2)

- ightharpoonup f and g are both proper, closed, and convex.
- ▶ The solution set  $S^* := \{ \mathbf{x}^* \in \text{dom}(F) : F(\mathbf{x}^*) = F^* \}$  is nonempty.

#### Remarks:

- $\circ$  Without loss of generality, f is smooth and g is non-smooth in the sequel.
- $\circ$  By Moreau-Rockafellar Theorem, we have  $\partial F = \partial (f+g) = \partial f + \partial g = \nabla f + \partial g$ .
- o Subgradient method attains a  $\mathcal{O}\left(\frac{1}{\sqrt{T}}\right)$  rate.
- o Without g, accelerated gradient method attains a  $\mathcal{O}\left(\frac{1}{T^2}\right)$  rate.

### Composite convex minimization

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- $\circ$  Without  $g_{\text{\tiny I}}$  accelerated gradient method attains a  $\mathcal{O}\left(\frac{1}{T^2}\right)$  rate.

Can we design algorithms that achieve a faster convergence rate for composite convex minimization?

# Designing algorithms for finding a solution $\mathbf{x}^{\star}$

# Quadratic majorizer for f

When f has L-Lipschitz continuous gradient, we have,  $\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^p$ 

$$f(\mathbf{x}) \leq f(\mathbf{y}) + \nabla f(\mathbf{y})^T (\mathbf{x} - \mathbf{y}) + \frac{L}{2} \|\mathbf{x} - \mathbf{y}\|_2^2$$

# Designing algorithms for finding a solution $\mathbf{x}^{\star}$

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### Quadratic *majorizer* for f + g

When f has L-Lipschitz continuous gradient, we have,  $\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^p$ 

$$f(\mathbf{x}) + g(\mathbf{x}) \leq f(\mathbf{y}) + \nabla f(\mathbf{y})^T (\mathbf{x} - \mathbf{y}) + \frac{L}{2} \|\mathbf{x} - \mathbf{y}\|_2^2 + g(\mathbf{x}) \coloneqq P_L(\mathbf{x}, \mathbf{y})$$

# Designing algorithms for finding a solution $\mathbf{x}^{\star}$

### Quadratic majorizer for f

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### Quadratic *majorizer* for f + g

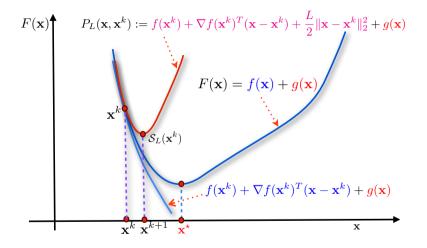
When f has L-Lipschitz continuous gradient, we have,  $\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^p$ 

$$f(\mathbf{x}) + g(\mathbf{x}) \le f(\mathbf{y}) + \nabla f(\mathbf{y})^T (\mathbf{x} - \mathbf{y}) + \frac{L}{2} ||\mathbf{x} - \mathbf{y}||_2^2 + g(\mathbf{x}) := P_L(\mathbf{x}, \mathbf{y})$$

### Majorization-minimization for f + g

$$\begin{aligned} \mathbf{x}^{k+1} &= \arg\min_{\mathbf{x} \in \mathbb{R}^p} P_L(\mathbf{x}, \mathbf{x}^k) \\ &= \arg\min_{\mathbf{x} \in \mathbb{R}^p} \left\{ g(\mathbf{x}) + \frac{L}{2} \| \mathbf{x} - \left( \mathbf{x}^k - \frac{1}{L} \nabla f(\mathbf{x}^k) \right) \|^2 \right\} \end{aligned}$$

#### Geometric illustration



### A short detour: Proximal-point operators

### Definition (Proximal operator [18])

Let  $g\in\mathcal{F}(\mathbb{R}^p)$ ,  $\mathbf{x}\in\mathbb{R}^p$  and  $\lambda>0$ . The proximal operator (or prox-operator) of g is defined as:

$$\operatorname{prox}_{\lambda g}(\mathbf{y}) \equiv \arg \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ g(\mathbf{x}) + \frac{1}{2\lambda} \|\mathbf{y} - \mathbf{x}\|_2^2 \right\}.$$
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### A short detour: Proximal-point operators

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 (3)

Remarks:

 $\circ$  The *proximal operator* of  $rac{1}{L}g$  evaluated at  $\left(\mathbf{x}^k - rac{1}{L} 
abla f(\mathbf{x}^k)
ight)$  is given by

$$\operatorname{prox}_{\frac{1}{L}g}\left(\mathbf{x}^k - \frac{1}{L}\nabla f(\mathbf{x}^k)\right) = \arg\min_{\mathbf{x} \in \mathbb{R}^p} \left\{g(\mathbf{x}) + \frac{L}{2} \|\, \mathbf{x} - \left(\mathbf{x}^k - \frac{1}{L}\nabla f(\mathbf{x}^k)\right)\|^2\right\}.$$

o This prox-operator minimizes the majorizing bound:

$$f(\mathbf{x}) + g(\mathbf{x}) \le f(\mathbf{x}^k) + \nabla f(\mathbf{x}^k)^T (\mathbf{x} - \mathbf{x}^k) + \frac{L}{2} ||\mathbf{x} - \mathbf{x}^k||_2^2 + g(\mathbf{x})$$

o Rule of thumb: Replace gradient steps with proximal gradient steps!

### **Tractable prox-operators**

### Processing non-smooth terms in (2)

- $\blacktriangleright$  We handle the nonsmooth term g in (2) using its proximal operator.
- lacktriangle However, computing proximal operator  $\mathrm{prox}_q$  of a general convex function g

$$\operatorname{prox}_g(\mathbf{y}) \equiv \arg\min_{\mathbf{x} \in \mathbb{R}^p} \left\{ g(\mathbf{x}) + \frac{1}{2} \|\mathbf{y} - \mathbf{x}\|_2^2 \right\}.$$

can be computationally demanding.

### Definition (Tractable proximity)

- ▶ Given  $g \in \mathcal{F}(\mathbb{R}^p)$ . We say that g is proximally tractable if  $\operatorname{prox}_q$  defined by (3) can be computed efficiently.
- ▶ "efficiently" = {closed form solution, low-cost computation, polynomial time}.

### Tractable prox-operators

### Example

For separable functions, the prox-operator can be efficient. When  $g(\mathbf{x}) := \|\mathbf{x}\|_1 = \sum_{i=1}^p |\mathbf{x}_i|$ , we have

$$\mathrm{prox}_{\lambda g}(\mathbf{x}) = \mathrm{sign}(\mathbf{x}) \otimes \max\{|\mathbf{x}| - \lambda, 0\}.$$

▶ Sometimes, we can compute the prox-operator via basic algebra. When  $g(\mathbf{x}) := \frac{1}{2} \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2$ , we have

$$\operatorname{prox}_{\lambda g}(\mathbf{x}) = \left(\mathbb{I} + \lambda \mathbf{A}^T \mathbf{A}\right)^{-1} \left(\mathbf{x} + \lambda \mathbf{A}^T \mathbf{b}\right).$$

ightharpoonup For the indicator functions of simple sets, e.g.,  $g(\mathbf{x}) := \delta_{\mathcal{X}}(\mathbf{x})$ , the prox-operator is the projection operator

$$\operatorname{prox}_{\lambda g}(\mathbf{x}) := \pi_{\mathcal{X}}(\mathbf{x}),$$

where  $\pi_{\mathcal{X}}(\mathbf{x})$  denotes the projection of  $\mathbf{x}$  onto  $\mathcal{X}$ . For instance, when  $\mathcal{X} = \{\mathbf{x} : \|\mathbf{x}\|_1 \leq \lambda\}$ , the projection can be obtained efficiently.

# Computational efficiency - Example

### Proximal operator of quadratic function

The **proximal operator** of a quadratic function  $g(\mathbf{x}) := \frac{1}{2} \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2$  is defined as

$$\operatorname{prox}_{\lambda g}(\mathbf{x}) := \arg \min_{\mathbf{y} \in \mathbb{R}^p} \left\{ \frac{1}{2} \|\mathbf{A}\mathbf{y} - \mathbf{b}\|_2^2 + \frac{1}{2\lambda} \|\mathbf{y} - \mathbf{x}\|_2^2 \right\}.$$
 (4)

How do we compute  $\operatorname{prox}_{\lambda q}(\mathbf{x})$ ?

The derivation: o The optimality condition implies that the solution of (4) should satisfy the following:

$$\mathbf{A}^{T}(\mathbf{A}\mathbf{y} - \mathbf{b}) + \lambda^{-1}(\mathbf{y} - \mathbf{x}) = 0.$$

• Setting  $\mathbf{y} = \operatorname{prox}_{\lambda_{\mathcal{A}}}(\mathbf{x})$ , we obtain

$$\operatorname{prox}_{\lambda g}(\mathbf{x}) = \left(\mathbb{I} + \lambda \mathbf{A}^T \mathbf{A}\right)^{-1} \left(\mathbf{x} + \lambda \mathbf{A}^T \mathbf{b}\right)$$

Remarks:

- $\circ$  The Woodbury matrix identity can be useful:  $(\mathbb{I} + \lambda \mathbf{A}^T \mathbf{A})^{-1} = \mathbb{I} \mathbf{A}^T (\lambda^{-1} \mathbb{I} + \mathbf{A} \mathbf{A}^T)^{-1} \mathbf{A}$ .
- $\circ$  When  $\mathbf{A}^T\mathbf{A}$  is efficiently diagonalizable, i.e.,  $\mathbf{A}^T\mathbf{A}:=\mathbf{U}\mathbf{\Lambda}\mathbf{U}^T$ , such that
  - $ightharpoonup \mathbf{U}$  is a unitary matrix, i.e.,  $\mathbf{U}\mathbf{U}^T = \mathbf{U}^T\mathbf{U} = \mathbb{I}$ , and  $\Lambda$  is a diagonal matrix.

### A non-exhaustive list of proximal tractability functions

Name	Function	Proximal operator	Complexity
$\ell_1$ -norm	$f(\mathbf{x}) := \ \mathbf{x}\ _1$	$\operatorname{prox}_{\lambda f}(\mathbf{x}) = \operatorname{sign}(\mathbf{x}) \otimes [ \mathbf{x}  - \lambda]_{+}$	$\mathcal{O}(p)$
$\ell_2$ -norm	$f(\mathbf{x}) := \ \mathbf{x}\ _2$	$\operatorname{prox}_{\lambda f}(\mathbf{x}) = \left[1 - \frac{\lambda}{\ \mathbf{x}\ _2}\right]_+ \mathbf{x}$	$\mathcal{O}(p)$
Support function	$f(\mathbf{x}) := \max_{\mathbf{y} \in \mathcal{C}} \mathbf{x}^T \mathbf{y}$	$\operatorname{prox}_{\lambda f}(\mathbf{x}) = \mathbf{x} - \lambda \pi_{\mathcal{C}}(\mathbf{x})$	
Box indicator	$f(\mathbf{x}) := \delta_{[\mathbf{a}, \mathbf{b}]}(\mathbf{x})$	$\operatorname{prox}_{\lambda f}(\mathbf{x}) = \pi_{[\mathbf{a}, \mathbf{b}]}(\mathbf{x})$	$\mathcal{O}(p)$
Positive semidefinite	$f(\mathbf{X}) := \delta_{\mathbb{S}^p}(\mathbf{X})$	$\mathrm{prox}_{\lambda f}(\mathbf{X}) = \mathbf{U}[\Sigma]_{+}\mathbf{U}^{T}$ , where $\mathbf{X} =$	$\mathcal{O}(p^3)$
cone indicator	+	$\mathbf{U}\Sigma\mathbf{U}^T$	
Hyperplane indicator	$f(\mathbf{x}) := \delta_{\mathcal{X}}(\mathbf{x}), \ \mathcal{X} :=$	$\operatorname{prox}_{\lambda f}(\mathbf{x}) = \pi_{\mathcal{X}}(\mathbf{x}) = \mathbf{x} +$	$\mathcal{O}(p)$
	$\{\mathbf{x} : \mathbf{a}^T \mathbf{x} = b\}$	$\left(\frac{b-\mathbf{a}^T\mathbf{x}}{\ \mathbf{a}\ _2}\right)\mathbf{a}$	
Simplex indicator	$f(\mathbf{x}) = \delta_{\mathcal{X}}(\mathbf{x}), \mathcal{X} :=$	$\operatorname{prox}_{\lambda f}(\mathbf{x}) = (\mathbf{x} - \nu 1) \text{ for some } \nu \in \mathbb{R},$	$ ilde{\mathcal{O}}(p)$
	$\{\mathbf{x} : \mathbf{x} \geq 0, 1^T \mathbf{x} = 1\}$	which can be efficiently calculated	
Convex quadratic	$f(\mathbf{x}) := \frac{1}{2} \mathbf{x}^T \mathbf{Q} \mathbf{x} + \mathbf{q}^T \mathbf{x}$	$\operatorname{prox}_{\lambda f}(\mathbf{x}) = (\lambda \mathbb{I} + \mathbf{Q})^{-1} \mathbf{x}$	$\mathcal{O}(p \log p)$
	_		$\mathcal{O}(p^3)$
Square $\ell_2$ -norm	$f(\mathbf{x}) := \frac{1}{2} \ \mathbf{x}\ _2^2$	$\operatorname{prox}_{\lambda f}(\mathbf{x}) = \frac{1}{1+\lambda}\mathbf{x}$	$\mathcal{O}(p)$
log-function	$f(\mathbf{x}) := -\log(x)$	$\operatorname{prox}_{\lambda f}(x) = \frac{1}{2}(\sqrt{x^2 + 4\lambda} + x)$	$\mathcal{O}(1)$
log det-function	$f(\mathbf{x}) := -\log \det(\mathbf{X})$	$\operatorname{prox}_{\lambda f}(\mathbf{X})$ is the log-function prox applied to the individual eigenvalues of $\mathbf{X}$	$\mathcal{O}(p^3)$

#### Remarks:

- $\circ \ \mathsf{Here} \colon [\mathbf{x}]_+ := \max\{0, \mathbf{x}\} \ \mathsf{and} \ \delta_{\mathcal{X}} \ \mathsf{is the indicator function of the convex set} \ \mathcal{X}.$
- $\circ \mathrm{\ sign}$  is the sign function,  $\mathbb{S}^p_+$  is the cone of symmetric positive semidefinite matrices.
- o For more functions, see [5, 15].

#### Solution methods

### Composite convex minimization

$$F^* := \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ F(\mathbf{x}) := f(\mathbf{x}) + g(\mathbf{x}) \right\}.$$
 (5)

### Choice of numerical solution methods

 $\circ$  Solve (5) = Find  $\mathbf{x}^k \in \mathbb{R}^p$  such that

$$F(\mathbf{x}^k) - F^* \le \varepsilon$$

for a given tolerance  $\varepsilon > 0$ .

- o Oracles: We can use one of the following configurations (oracles):
  - 1.  $\partial f(\cdot)$  and  $\partial g(\cdot)$  at any point  $\mathbf{x} \in \mathbb{R}^p$ .
  - 2.  $\nabla f(\cdot)$  and  $\operatorname{prox}_{\lambda g}(\cdot)$  at any point  $\mathbf{x} \in \mathbb{R}^p$ .
  - 3.  $\operatorname{prox}_{\lambda f}$  and  $\operatorname{prox}_{\lambda g}(\cdot)$  at any point  $\mathbf{x} \in \mathbb{R}^p$ .
  - 4.  $\nabla f(\cdot)$ , inverse of  $\nabla^2 f(\cdot)$  and  $\operatorname{prox}_{\lambda q}(\cdot)$  at any point  $\mathbf{x} \in \mathbb{R}^p$ .

Remark: Using different oracle leads to different types of algorithms.

## Proximal-gradient algorithm

### Basic proximal-gradient scheme (ISTA)

- **1.** Choose  $\mathbf{x}^0 \in \text{dom}(F)$  arbitrarily as a starting point.
- **2.** For  $k = 0, 1, \dots$ , generate a sequence  $\{\mathbf{x}^k\}_{k \geq 0}$  as:

$$\mathbf{x}^{k+1} := \operatorname{prox}_{\alpha g} \left( \mathbf{x}^k - \alpha \nabla f(\mathbf{x}^k) \right),$$

where  $\alpha := \frac{1}{L}$ .

# Proximal-gradient algorithm

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where  $\alpha := \frac{1}{L}$ .

## Theorem (Convergence of ISTA [2])

Let  $\{\mathbf{x}^k\}$  be generated by ISTA. Then:

$$F(\mathbf{x}^k) - F^* \le \frac{L_f \|\mathbf{x}^0 - \mathbf{x}^*\|_2^2}{2(k+1)}$$

The worst-case complexity to reach  $F(\mathbf{x}^k) - F^* \leq \varepsilon$  of (ISTA) is  $\mathcal{O}\left(\frac{L_f R_0^2}{\varepsilon}\right)$ , where  $R_0 := \max_{\mathbf{x}^* \in S^*} \|\mathbf{x}^0 - \mathbf{x}^*\|_2$ .

- $\circ$  Oracles:  $\operatorname{prox}_{\alpha q}(\cdot)$  and  $\nabla f(\cdot)$ .
- o Compared to the subgradient gradient method, the rate improves at the cost of prox-computation.

### Fast proximal-gradient algorithm

### Fast proximal-gradient scheme (FISTA)

- **1.** Choose  $\mathbf{x}^0 \in \text{dom}(F)$  arbitrarily as a starting point.
- **2.** Set  $\mathbf{y}^0 := \mathbf{x}^0$  and  $t_0 := 1$ ,  $\alpha := L^{-1}$ .
- **3.** For  $k=0,1,\ldots$ , generate two sequences  $\{\mathbf{x}^k\}_{k\geq 0}$  and  $\{\mathbf{y}^k\}_{k\geq 0}$  as:

$$\begin{cases} \mathbf{x}^{k+1} &:= \operatorname{prox}_{\alpha g} \left( \mathbf{y}^k - \alpha \nabla f(\mathbf{y}^k) \right), \\ t_{k+1} &:= \frac{1}{2} (1 + \sqrt{4t_k^2 + 1}), \\ \mathbf{y}^{k+1} &:= \mathbf{x}^{k+1} + \frac{t_{k-1}}{t_{k+1}} (\mathbf{x}^{k+1} - \mathbf{x}^k). \end{cases}$$

## Fast proximal-gradient algorithm

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### Theorem (Convergence of FISTA [2])

Let  $\{\mathbf{x}^k\}$  be generated by FISTA. Then:

$$F(\mathbf{x}^k) - F^* \le \frac{2L_f \|\mathbf{x}^0 - \mathbf{x}^*\|_2^2}{(k+1)^2}$$

The worst-case complexity to reach  $F(\mathbf{x}^k) - F^* \leq \varepsilon$  of (FISTA) is  $\mathcal{O}\left(R_0\sqrt{\frac{L_f}{\varepsilon}}\right)$ ,  $R_0 := \max_{\mathbf{x}^* \in \mathcal{S}^*} \|\mathbf{x}^0 - \mathbf{x}^*\|_2$ .

### Fast proximal-gradient algorithm

### Fast proximal-gradient scheme (FISTA)

- **1.** Choose  $\mathbf{x}^0 \in \text{dom}(F)$  arbitrarily as a starting point.
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Remark:

From 
$$\mathcal{O}\left(\frac{L_f R_0^2}{\epsilon}\right)$$
 to  $\mathcal{O}\left(R_0 \sqrt{\frac{L_f}{\epsilon}}\right)$  iterations at almost no additional cost!.

### Complexity per iteration

- ▶ One gradient  $\nabla f(\mathbf{y}^k)$  and one prox-operator of g;
- ▶ 8 arithmetic operations for  $t_{k+1}$  and  $\gamma_{k+1}$ ;
- ▶ 2 more vector additions, and **one** scalar-vector multiplication.

The cost per iteration is almost the same as in gradient scheme if proximal operator of g is efficient.

## Example 1: $\ell_1$ -regularized least squares

## Problem ( $\ell_1$ -regularized least squares)

Given  $\mathbf{A} \in \mathbb{R}^{n \times p}$  and  $\mathbf{b} \in \mathbb{R}^n$ , solve:

$$F^{\star} := \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ F(\mathbf{x}) := \frac{1}{2} \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2 + \lambda \|\mathbf{x}\|_1 \right\}, \tag{6}$$

where  $\lambda > 0$  is a regularization parameter.

### Complexity per iterations

- Evaluating  $\nabla f(\mathbf{x}^k) = \mathbf{A}^T(\mathbf{A}\mathbf{x}^k \mathbf{b})$  requires one  $\mathbf{A}\mathbf{x}$  and one  $\mathbf{A}^T\mathbf{y}$ .
- One soft-thresholding operator  $\operatorname{prox}_{\lambda a}(\mathbf{x}) = \operatorname{sign}(\mathbf{x}) \otimes \max\{|\mathbf{x}| \lambda, 0\}.$
- ▶ Optional: Evaluating  $L = \|\mathbf{A}^T\mathbf{A}\|$  (spectral norm) via power iterations

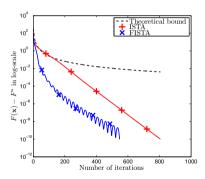
### Synthetic data generation

- $ightharpoonup \mathbf{A} := \operatorname{randn}(n,p)$  standard Gaussian  $\mathcal{N}(0,\mathbb{I})$ .
- $\mathbf{x}^{\star}$  is a k-sparse vector generated randomly.
- $\mathbf{b} := \mathbf{A} \mathbf{x}^* + \mathcal{N}(0, 10^{-3}).$

## **Example 1: Theoretical bounds vs practical performance**

#### Theoretical bounds

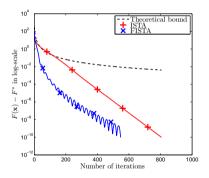
We have the following guarantees for **FISTA** :=  $\frac{2L_fR_0^2}{(k+2)^2}$  and for **ISTA** :=  $\frac{L_fR_0^2}{2(k+2)}$ . In the figure below, ISTA's practical behavior outperforms the theoretical bound for FISTA.

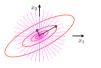


### **Example 1: Theoretical bounds vs practical performance**

### Theoretical bounds

We have the following guarantees for FISTA :=  $\frac{2L_fR_0^2}{(k+2)^2}$  and for ISTA :=  $\frac{L_fR_0^2}{2(k+2)}$ . In the figure below, ISTA's practical behavior outperforms the theoretical bound for FISTA.







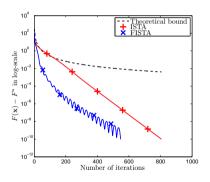
descent directions

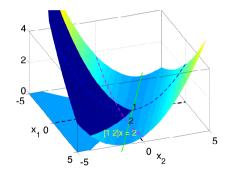
restricted descent directions

### **Example 1: Theoretical bounds vs practical performance**

#### Theoretical bounds

We have the following guarantees for FISTA :=  $\frac{2L_fR_0^2}{(k+2)^2}$  and for ISTA :=  $\frac{L_fR_0^2}{2(k+2)}$ . In the figure below, ISTA's practical behavior outperforms the theoretical bound for FISTA.





Remarks:

- $\circ$   $\ell_1\text{-regularized}$  least squares formulation has restricted strong convexity.
- The proximal-gradient method can automatically exploit this structure.

## Example 2: Sparse logistic regression

# Problem (Sparse logistic regression)

Given  $\mathbf{A} \in \mathbb{R}^{n \times p}$  and  $\mathbf{b} \in \{-1, +1\}^n$ , solve:

$$F^{\star} := \min_{\mathbf{x}, \beta} \left\{ F(\mathbf{x}) := \frac{1}{n} \sum_{j=1}^{n} \log \left( 1 + \exp \left( -\mathbf{b}_{j} (\mathbf{a}_{j}^{T} \mathbf{x} + \beta) \right) \right) + \rho \|\mathbf{x}\|_{1} \right\}.$$

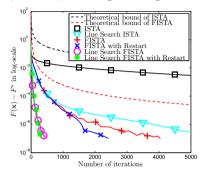
### Real data

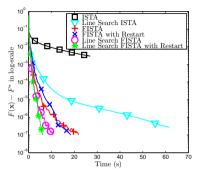
- ▶ Real data: w8a with n = 49'749 data points. p = 300 features
- ► Available at http://www.csie.ntu.edu.tw/~cjlin/libsvmtools/datasets/binary.html.

#### **Parameters**

- $\rho = 10^{-4}$ .
- Number of iterations 5000, tolerance  $10^{-7}$ .
- ▶ Ground truth: Solve problem up to  $10^{-9}$  accuracy by TFOCS to get a high accuracy approximation of  $\mathbf{x}^*$  and  $F^*$ .

### Example 2: Sparse logistic regression - numerical results





	ISTA	LS-ISTA	FISTA	FISTA-R	LS-FISTA	LS-FISTA-R
Number of iterations	5000	5000	4046	2423	447	317
CPU time (s)	26.975	61.506	21.859	18.444	10.683	6.228
Solution error $(\times 10^{-7})$	29370	2.774	1.000	0.998	0.961	0.985

# When f is strongly convex: Algorithms

### Proximal-gradient scheme (ISTA<sub> $\mu$ </sub>)

- **1.** Given  $\mathbf{x}^0 \in \mathbb{R}^p$  as a starting point.
- **2.** For  $k = 0, 1, \dots$ , generate a sequence  $\{\mathbf{x}^k\}_{k \geq 0}$  as:

$$\mathbf{x}^{k\!+\!1}\!\!:=\!\operatorname{prox}_{\alpha_k g}\!\!\left(\!\mathbf{x}^k\!-\!\alpha_k \nabla f(\mathbf{x}^k)\!\right)\!,$$

where  $\alpha_k := \frac{2}{L_f + \mu}$  is the optimal step-size.

### Fast proximal-gradient scheme (FISTA $_{\mu}$ )

- **1.** Given  $\mathbf{x}^0 \in \mathbb{R}^p$  as a starting point. Set  $\mathbf{y}^0 := \mathbf{x}^0$ .
- **2.** For  $k=0,1,\cdots$  , generate sequences  $\{\mathbf{x}^k\}_{k\geq 0}$  and  $\{\mathbf{y}^k\}_{k\geq 0}$  as:

$$\begin{cases} \mathbf{x}^{k+1} := \operatorname{prox}_{\alpha_k g} \Big( \mathbf{y}^k - \alpha_k \nabla f(\mathbf{y}^k) \Big), \\ \mathbf{y}^{k+1} := \mathbf{x}^{k+1} + \Big( \frac{\sqrt{c_f} - 1}{\sqrt{c_f} + 1} \Big) (\mathbf{x}^{k+1} - \mathbf{x}^k), \end{cases}$$

where  $c_f:=rac{L_f}{\mu}$  and  $lpha_k:=L_f^{-1}$  is the optimal step-size.

# When f is strongly convex: Convergence

### Assumption

f is strongly convex with parameter  $\mu > 0$ , i.e.,  $f \in \mathcal{F}^{1,1}_{L,\mu}(\mathbb{R}^p)$ .

**Condition number:**  $c_f := \frac{L_f}{\mu} \geq 0.$ 

# Theorem (ISTA $_{\mu}$ [14])

$$F(\mathbf{x}^k) - F^* \le \frac{L_f}{2} \left( \frac{c_f - 1}{c_f + 1} \right)^{2k} \|\mathbf{x}^0 - \mathbf{x}^*\|_2^2.$$

Convergence rate: Linear with contraction factor:  $\omega := \left(\frac{c_f-1}{c_f+1}\right)^2 = \left(\frac{L_f-\mu}{L_f+\mu}\right)^2$ .

# Theorem (**FISTA** $_{\mu}$ [14])

$$F(\mathbf{x}^k) - F^* \le \frac{L_f + \mu}{2} \left( 1 - \sqrt{\frac{\mu}{L_f}} \right)^k \|\mathbf{x}^0 - \mathbf{x}^*\|_2^2.$$

Convergence rate: Linear with contraction factor:  $\omega_f = \frac{\sqrt{L_f} - \sqrt{\mu}}{\sqrt{L_f}} < \omega$ .

# Summary of the worst-case complexities

### Comparison

Complexity	Proximal-gradient scheme	Fast proximal-gradient
		scheme
Complexity $[\mu=0]$	$\mathcal{O}\left(R_0^2 \frac{L_f}{\varepsilon}\right)$	$\mathcal{O}\left(R_0\sqrt{rac{L_f}{arepsilon}} ight)$
Per iteration	1-gradient, 1-prox, 1- $sv$ , 1-	1-gradient, 1-prox, 2-sv, 3-
	v+	v+
Complexity $[\mu > 0]$	$\mathcal{O}\left(\kappa\log(\varepsilon^{-1})\right)$	$\mathcal{O}\left(\sqrt{\kappa}\log(\varepsilon^{-1})\right)$
Per iteration	1-gradient, 1-prox, 1- $sv$ , 1-	1-gradient, 1-prox, 1-sv, 2-
	v+	v+

Here: sv = scalar-vector multiplication, v+=vector addition.

$$R_0 := \max_{\mathbf{x}^{\star} \in \mathcal{S}^{\star}} \|\mathbf{x}^0 - \mathbf{x}^{\star}\|$$
 and  $\kappa = \frac{L_f}{\mu_f}$  is the condition number.

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	v+	v+

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#### Need alternatives when

ightharpoonup computing  $\nabla f(\mathbf{x})$  is much costlier than computing  $\operatorname{prox}_g$ 

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	v+	v+

Here: sv = scalar-vector multiplication, v+=vector addition.

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#### Need alternatives when

ightharpoonup computing  $\nabla f(\mathbf{x})$  is much costlier than computing  $\operatorname{prox}_a$ 

### Software

**TFOCS** is a good software package to learn about first order methods.

### Composite minimization: Non-convex case

# Problem (Unconstrained composite minimization)

$$F^* := \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ F(\mathbf{x}) := f(\mathbf{x}) + g(\mathbf{x}) \right\}$$
 (CM)

- ▶  $g: \mathbb{R}^p \to \mathbb{R} \cup \{\infty\}$  is proper, closed, convex, and (possibly) nonsmooth.
- $f: \mathbb{R}^p \to \mathbb{R}$  is proper and closed, dom(f) is convex, and f is  $L_f$ -smooth.
- $ightharpoonup \operatorname{dom}(F) := \operatorname{dom}(f) \cap \operatorname{dom}(g) \neq \emptyset \text{ and } -\infty < F^{\star} < +\infty.$
- ▶ The solution set  $S^* := \{\mathbf{x}^* \in \text{dom}(F) : F(\mathbf{x}^*) = F^*\}$  is nonempty.

# A different quantification of convergence: Gradient mapping

# Definition (Gradient mapping)

Let  $\mathrm{prox}_g$  denote the proximal operator of g and  $\lambda>0$  some real constant. Then, the gradient mapping operator is defined as

$$\mathcal{G}_{\lambda}(\mathbf{x}) := \frac{1}{\lambda} \left( \mathbf{x} - \operatorname{prox}_{\lambda g}(\mathbf{x} - \lambda \nabla f(\mathbf{x})) \right).$$

## Properties [1]

- $\|\mathcal{G}_{\lambda}(\mathbf{x})\| = 0 \iff \mathbf{x} \text{ is a stationary point.}$
- $\qquad \qquad \textbf{Lipschitz continuity:} \ \parallel \mathcal{G}_{\frac{1}{L}}(\mathbf{x}) \mathcal{G}_{\frac{1}{L}}(\mathbf{y}) \parallel \leq (2L + L_f) \parallel \mathbf{x} \mathbf{y} \parallel$

# Why do we care about gradient mapping?

- It is the generalization of the gradient of f,  $\nabla f(\mathbf{x})$
- ▶ Recall prox-gradient update:  $\mathbf{x}^{t+1} = \text{prox}_{\lambda a}(\mathbf{x}^t \lambda \nabla f(\mathbf{x}^t))$ , which is equivalent to  $\mathbf{x}^{t+1} = \mathbf{x}^t \lambda \mathcal{G}_{\lambda}(\mathbf{x}^t)$ .
- ▶ In fact, when  $\text{prox}_q = \mathbb{I}$ , then,  $\mathcal{G}_{\lambda}(\mathbf{x}) = \frac{1}{\lambda} \left( \mathbf{x} (\mathbf{x} \lambda \nabla f(\mathbf{x})) \right) = \nabla f(\mathbf{x})$ .

## Sufficient Decrease property for proximal-gradient

## Assumption

- f is  $L_f$ -smooth.
- ightharpoonup g is proper, closed, convex, and (possibly) nonsmooth. g is proximally tractable.

$$\mathbf{x}^{k+1} := \operatorname{prox}_{\frac{1}{L}g} \left( \mathbf{x}^k - \frac{1}{L} \nabla f(\mathbf{x}^k) \right)$$

## Lemma (Sufficient decrease [1])

For any  $\mathbf{x} \in \operatorname{int}(\operatorname{dom}(f))$  and  $L \in (\frac{L_f}{2}, \infty)$ , it holds that

$$F(\mathbf{x}^{k+1}) \le F(\mathbf{x}^k) - \frac{L - \frac{L_f}{2}}{L^2} \left\| \mathcal{G}_{\frac{1}{L}}(\mathbf{x}^k) \right\|_2^2, \tag{7}$$

### Corollary

$$F(\mathbf{x}^{k+1}) \le F(\mathbf{x}^k) - \frac{1}{2L_f} \left\| \mathcal{G}_{\frac{1}{L_f}}(\mathbf{x}^k) \right\|_2^2, \quad \text{for } L = L_f$$

## Non-convex case: Convergence

#### Basic proximal-gradient scheme

- **1.** Choose  $\mathbf{x}^0 \in \text{dom}(F)$  arbitrarily as a starting point.
- **2.** For  $k=0,1,\cdots$ , generate a sequence  $\{\mathbf{x}^k\}_{k\geq 0}$  as:

$$\mathbf{x}^{k+1} := \operatorname{prox}_{\alpha g} \left( \mathbf{x}^k - \alpha \nabla f(\mathbf{x}^k) \right),\,$$

where  $\alpha:=\left(0,\frac{2}{L_f}\right)$ .

## Theorem (Convergence of proximal-gradient method: Non-convex [1])

Let  $\{\mathbf{x}^k\}$  be generated by proximal-gradient scheme above. Then, we have

$$\min_{i=0,1,\ldots,k} \|\mathcal{G}_{\alpha}(\mathbf{x}^i)\|_2^2 \leq \frac{F(\mathbf{x}^0) - F(\mathbf{x}^\star)}{M(k+1)}, \qquad \qquad \textit{where } M := \alpha^2 \left(\frac{1}{\alpha} - \frac{L_f}{2}\right)$$

- When  $\alpha = \frac{1}{L_f}$ ,  $M = \frac{1}{2L_f}$ .
- ► The worst-case complexity to reach  $\min_{i=0,1,\dots,k} \|\mathcal{G}_{\alpha}(\mathbf{x}^i)\|_2^2 \leq \varepsilon$  is  $\mathcal{O}\left(\frac{1}{\varepsilon}\right)$ .

### Stochastic convex composite minimization

## Problem (Mathematical formulation)

Consider the following composite convex minimization problem:

$$F^{\star} = \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ F(\mathbf{x}) := \mathbb{E}_{\theta}[F(\mathbf{x}, \theta)] := \mathbb{E}_{\theta}[f(\mathbf{x}, \theta) + g(\mathbf{x}, \theta)] \right\}$$

- ightharpoonup heta is a random vector whose probability distribution is supported on set  $\Theta$ .
- ▶ The solution set  $S^* := \{ \mathbf{x}^* \in \text{dom}(F) : F(\mathbf{x}^*) = F^* \}$  is nonempty.
- ▶ Oracles: (sub)gradient of  $f(\cdot, \theta)$ ,  $\nabla f(\mathbf{x}, \theta)$ , and stochastic prox operator of  $g(\cdot, \theta)$ ,  $\operatorname{prox}_{q(\cdot, \theta)}(\mathbf{x})$ .

#### Remark

- o In this setting, we replace  $\nabla f(\cdot)$  with its stochastic estimates.
- $\circ$  It is possible to replace  $\text{prox}_a(\cdot)$  with its stochastic estimate (advanced material).

## Stochastic proximal gradient method

#### Stochastic proximal gradient method (SPG)

- **1.** Choose  $\mathbf{x}^0 \in \mathbb{R}^p$  and  $(\gamma_k)_{k \in \mathbb{N}} \in ]0, +\infty[^{\mathbb{N}}]$ .
- **2.** For  $k = 0, 1, \ldots$  perform:

$$\mathbf{x}^{k+1} = \operatorname{prox}_{\gamma_k g(\cdot, \theta)}(\mathbf{x}^k - \gamma_k G(\mathbf{x}^k, \theta_k)).$$

#### **Definitions:**

- $\circ \ \mathrm{\underline{prox}}_{\lambda g(\cdot,\theta)} := \arg \min_{\mathbf{y} \in \mathbb{R}^p} \left\{ g(\mathbf{y},\theta) + \tfrac{1}{2\lambda} \| \, \mathbf{y} \mathbf{x} \, \|^2 \right\}$
- $\circ \{\theta_k\}_{k=0,1,\dots}$ : sequence of independent random variables.
- $\circ G(\mathbf{x}^k, \theta_k) \in \partial f(\mathbf{x}^k, \theta_k)$ : an unbiased estimate of the deterministic (sub)gradient:

$$\mathbb{E}[G(\mathbf{x}^k, \theta_k)] \in \partial f(\mathbf{x}^k).$$

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#### Remark

Cost of computing  $G(\mathbf{x}^k, \theta_k)$  is usually much cheaper than  $\nabla f(\mathbf{x}^k)$ .

## Convergence analysis

## Assumptions for the problem setting

- $ightharpoonup f(\cdot, \theta)$  and  $g(\cdot, \theta)$  are convex functions in the first argument, g is proximally-tractable.
- ▶ (Sub)gradients of F satisfy stochastic bounded gradient condition:  $\exists C \geq 0, B \geq 0$  such that

$$\mathbb{E}_{\theta}[\|\partial F(\mathbf{x}, \theta)\|^2] \leq B^2 + C(F(\mathbf{x}) - F(\mathbf{x}^*)).$$

 $ightharpoonup \mathbb{E}[\|\mathbf{x}^t - \mathbf{x}^\star\|^2] \le R^2 \text{ for all } t \ge 0.$ 

### Implications of the assumptions

- ▶ None of the above assumptions enforce that *f* is smooth.
- ▶ Stochastic bounded gradient condition holds with C=0 when both  $f(\cdot,\theta)$  and  $g(\cdot,\theta)$  are Lipschitz continuous.
- lacktriangle The same condition holds when  $f(\cdot,\theta)$  is  $L_f$ -smooth and  $g(\cdot,\theta)$  is Lipschitz continuous.
- ▶ However, for the upcoming theorem, we will take C > 0, which rules out the case when both functions are only Lipschitz continuous.

## Convergence analysis

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 $\mathbb{E}[\|\mathbf{x}^t - \mathbf{x}^\star\|^2] \le R^2 \text{ for all } t \ge 0.$ 

## Theorem (Ergodic convergence [12])

- ightharpoonup Assume the above assumptions hold with C>0.
- Let the sequence  $\{\mathbf{x}^k\}_{k\geq 0}$  be generated by SPG.
- Set  $\gamma_k = \frac{1}{C\sqrt{k}}$ .

#### Conclusion:

• Define  $\bar{\mathbf{x}}^k = \frac{1}{k} \sum_{i=0}^{k-1} \mathbf{x}^i$ , then

$$\mathbb{E}[F(\bar{\mathbf{x}}^k) - F(\mathbf{x}^*)] \le \frac{1}{\sqrt{k}} \left( R^2 C + \frac{B^2}{C} \right), \quad \forall k \ge 1.$$

## Revisiting a special composite structure

## A basic constrained problem setting

$$f^* := \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ f(\mathbf{x}) + \delta_{\mathcal{X}}(\mathbf{x}) \right\} := \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ f(\mathbf{x}) : \mathbf{x} \in \mathcal{X} \right\}, \tag{8}$$

#### **Assumptions**

- $\triangleright$   $\mathcal{X}$  is nonempty, convex and compact (closed and bounded) where  $\delta_{\mathcal{X}}$  is its indicator function.
- $f \in \mathcal{F}_L^{1,1}(\mathbb{R}^p)$  (i.e., convex with Lipschitz gradient).

### Recall proximal gradient algorithm

#### Basic proximal-gradient scheme (ISTA)

- **1.** Choose  $\mathbf{x}^0 \in \text{dom}(F)$  arbitrarily as a starting point.
- **2.** For  $k = 0, 1, \dots$ , generate a sequence  $\{\mathbf{x}^k\}_{k \ge 0}$  as:

$$\mathbf{x}^{k+1} := \operatorname{prox}_{\alpha g} \left( \mathbf{x}^k - \alpha \nabla f(\mathbf{x}^k) \right)$$

where  $\alpha := \frac{1}{I}$ .

ightharpoonup Prox-operator of indicator of  $\mathcal{X}$  is projection onto  $\mathcal{X}$   $\Longrightarrow$  ensures feasibility

How else can we ensure feasibility?

### Frank-Wolfe's approach - I

$$f^{\star} := \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ f(\mathbf{x}) : \mathbf{x} \in \mathcal{X} \right\},$$

## Conditional gradient method (CGM, see [10] for review)

A plausible strategy which dates back to 1956 [6]. At iteration k:

1. Consider the linear approximation of f at  $\mathbf{x}^k$ 

$$\phi_k(\mathbf{x}) := f(\mathbf{x}^k) + \nabla f(\mathbf{x}^k)^T (\mathbf{x} - \mathbf{x}^k)$$

2. Minimize this approximation within constraint set

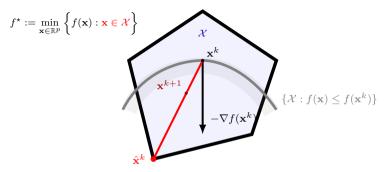
$$\hat{\mathbf{x}}^k \in \min_{\mathbf{x} \in \mathcal{X}} \phi_k(\mathbf{x}) = \min_{\mathbf{x} \in \mathcal{X}} \nabla f(\mathbf{x}^k)^T \mathbf{x}$$

3. Take a step towards  $\hat{\mathbf{x}}^k$  with step-size  $\gamma_k \in [0,1]$ 

$$\mathbf{x}^{k+1} = \mathbf{x}^k + \gamma_k (\hat{\mathbf{x}}^k - \mathbf{x}^k)$$

 $\mathbf{x}^{k+1}$  is feasible since it is convex combination of two other feasible points.

### Frank-Wolfe's approach - II



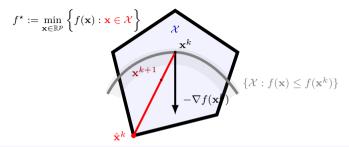
### Conditional gradient method (CGM)

- 1. Choose  $\mathbf{x}^0 \in \mathcal{X}$ .
- **2.** For  $k = 0, 1, \ldots$  perform:

$$\begin{cases} \hat{\mathbf{x}}^k &:= \arg\min_{\mathbf{x} \in \mathcal{X}} \nabla f(\mathbf{x}^k)^T \mathbf{x} \\ \mathbf{x}^{k+1} &:= (1-\gamma_k) \mathbf{x}^k + \gamma_k \hat{\mathbf{x}}^k, \end{cases}$$

where  $\gamma_k := \frac{2}{k+2}$ .

#### On the linear minimization oracle



### Definition (Linear minimization oracle)

Let  $\mathcal{X}$  be a convex, closed and bounded set. Then, the linear minimization oracle of  $\mathcal{X}$  ( $lmo_{\mathcal{X}}$ ) returns a vector  $\hat{\mathbf{x}}$  such that

$$lmo_{\mathcal{X}}(\mathbf{x}) := \hat{\mathbf{x}} \in \arg\min_{\mathbf{y} \in \mathcal{X}} \mathbf{x}^{T} \mathbf{y}$$
(9)

- $ightharpoonup \operatorname{lmo}_{\mathcal{X}}$  returns an extreme point of  $\mathcal{X}$ .
- $ightharpoonup ext{lmo}_{\mathcal{X}}$  is arguably cheaper than projection.
- ▶  $lmo_{\mathcal{X}}$  is not single valued, note  $\in$  in the definition.

## Convergence guarantees of CGM

## Problem setting

$$f^* := \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ f(\mathbf{x}) : \mathbf{x} \in \mathcal{X} \right\},$$

#### **Assumptions**

- $\triangleright$   $\mathcal{X}$  is nonempty, convex, closed and bounded.
- $f \in \mathcal{F}^{1,1}_L(\mathbb{R}^p)$  (i.e., convex with Lipschitz gradient).

#### **Theorem**

Under assumptions listed above, CGM with step size  $\gamma_k = \frac{2}{k+2}$  satisfies

$$f(\mathbf{x}^k) - f(\mathbf{x}^*) \le \frac{4LD_{\mathcal{X}}^2}{k+1}$$
 (10)

where  $D_{\mathcal{X}} := \max_{\mathbf{x}, \mathbf{y} \in \mathcal{X}} \|\mathbf{x} - \mathbf{y}\|_2$  is diameter of constraint set.

## \*Convergence guarantees of CGM: A faster rate

### Problem setting

$$f^* := \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ f(\mathbf{x}) : \mathbf{x} \in \mathcal{X} \right\},$$

#### **Assumptions**

- $\triangleright$   $\mathcal{X}$  is nonempty,  $\alpha$ -strongly convex, closed and bounded.
- $f \in \mathcal{F}^{1,1}_{L,\mu}(\mathbb{R}^p)$  (i.e., strongly convex with Lipschitz gradient).

## Definition ( $\alpha$ -strongly convex set) [7]

A convex set  $\mathcal{X} \in \mathbb{R}^{p \times p}$  is  $\alpha$ -strongly convex with respect to  $\|\cdot\|$  if for any  $\mathbf{x}, \mathbf{y} \in \mathcal{X}$ , any  $\gamma \in [0,1]$  and any vector  $\mathbf{z} \in \mathbb{R}^{p \times p}$  such that  $\|\mathbf{z}\| = 1$ , it holds that

$$\gamma \mathbf{x} + (1 - \gamma)\mathbf{y} + \gamma (1 - \gamma)\frac{\alpha}{2} \|\mathbf{x} - \mathbf{y}\|^2 \mathbf{z} \in \mathcal{X}$$

That is, for any  $\mathbf{x}, \mathbf{y} \in \mathcal{X}$ , the ball centered at  $\gamma \mathbf{x} + (1 - \gamma) \mathbf{y}$  with radius  $\gamma (1 - \gamma) \frac{\alpha}{2} \| \mathbf{x} - \mathbf{y} \|^2$  is contained in  $\mathcal{X}$ .

## \*CGM for strongly convex objective + strongly convex set

#### Conditional gradient method - CGM2

- **1.** Choose  $\mathbf{x}^0 \in \mathcal{X}$ .
- **2.** For  $k = 0, 1, \ldots$  perform:

$$\begin{cases} \hat{\mathbf{x}}^k &:= \arg\min_{\mathbf{x} \in \mathcal{X}} \nabla f(\mathbf{x}^k)^T \mathbf{x} \\ \gamma_k &:= \arg\min_{\gamma \in [0,1]} \gamma \left\langle \hat{\mathbf{x}}^k - \mathbf{x}^k, \nabla f(\mathbf{x}^k) \right\rangle + \gamma^2 \frac{L}{2} \| \hat{\mathbf{x}}^k - \mathbf{x}^k \|^2 \\ \mathbf{x}^{k+1} &:= (1 - \gamma_k) \mathbf{x}^k + \gamma_k \hat{\mathbf{x}}^k, \end{cases}$$

### Theorem ([7])

Under assumptions listed previously, CGM2 satisfies

$$f(\mathbf{x}^k) - f(\mathbf{x}^*) = \mathcal{O}\left(\frac{1}{k^2}\right).$$
 (11)

## **Example:** lmo of nuclear-norm bal

Consider 
$$\delta_{\mathcal{X}}, \text{ the indicator of nuclear-norm ball } \mathcal{X} := \left\{ \mathbf{X} : \mathbf{X} \in \mathbb{R}^{p \times p}, \,\, \|\, \mathbf{X} \,\|_* \leq \alpha \right\}$$

### lmo of nuclear-norm ball

$$\mathrm{lmo}_{\mathcal{X}}(\mathbf{X}) := \hat{\mathbf{X}} \in \mathrm{arg} \min_{\mathbf{Y} \in \mathcal{X}} \ \langle \mathbf{Y}, \mathbf{X} \rangle$$

This can be computed as follows:

- lacktriangledown Compute top singular vectors of  $\mathbf{X} \implies (\mathbf{u}_1, \sigma_1, \mathbf{v}_1) = \mathtt{svds}(\mathbf{X}, 1)$ .
- $lackbox{
  ightharpoonup}$  Form the rank-1 output  $\implies$   $\mathbf{X} = -\mathbf{u}_1 lpha \mathbf{v}_1^T$

We can efficiently approximate top singular vectors by power method!

## Proximal gradient vs. Frank-Wolfe

#### **Definitions:**

- ▶ Here: sv = scalar-vector multiplication, v+=vector addition.
- $ightharpoonup R_0 := \max_{\mathbf{x}^{\star} \in \mathcal{S}^{\star}} \|\mathbf{x}^0 \mathbf{x}^{\star}\|$  is the maximum initial distance.
- $D_{\mathcal{X}} := \max_{\mathbf{x}, \mathbf{y} \in \mathcal{X}} \|\mathbf{x} \mathbf{y}\|_2$  is diameter of constraint set  $\mathcal{X}$ .

Algorithm	Proximal-gradient scheme	Frank-Wolfe method
Rate	$\mathcal{O}\left(\frac{L_f R_0^2}{k}\right)$	$\mathcal{O}\left(rac{L_f D_{\mathcal{X}}^2}{k} ight)$
Complexity	$\mathcal{O}\left(R_0^2 \frac{L_f}{\varepsilon}\right)$	$\mathcal{O}\left(D_{\mathcal{X}}^{2} \frac{L_{f}}{\varepsilon}\right)$
Per iteration	$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$

How do prox operator and lmo compare in practice?

### An example with matrices

#### **Problem Definition**

$$\min_{\mathbf{X} \in \mathbb{R}^{p \times p}} f(\mathbf{X}) + g(\mathbf{X})$$

- ▶ Define  $g(\mathbf{X}) = \delta_{\mathcal{X}}(\mathbf{X})$ , where  $\mathcal{X} := \left\{ \mathbf{X} : \mathbf{X} \in \mathbb{R}^{p \times p}, \ \|\mathbf{X}\|_* \leq \alpha \right\}$  is nuclear norm ball.
- ► This problem is equivalent to:

$$\min_{\mathbf{X} \in \mathcal{X}} f(\mathbf{X})$$

### Observations

- ▶  $\operatorname{prox}_q = \pi_{\mathcal{X}}$ . Projection requires full SVD,  $\mathcal{O}(p^3)$ .
- $\blacktriangleright$  lmo computes (approximately) top singular vectors, roughly in  $\approx \mathcal{O}(p^2)$  with Lanczos algorithm.

### **Example: Phase retrieval**

#### Phase retrieval

Aim: Recover signal  $\mathbf{x}^{\natural} \in \mathbb{C}^p$  from the measurements  $\mathbf{b} \in \mathbb{R}^n$ :

$$b_i = \left| \langle \mathbf{a}_i, \mathbf{x}^{\dagger} \rangle \right|^2 + \omega_i.$$

 $(\mathbf{a}_i \in \mathbb{C}^p \text{ are known measurement vectors, } \omega_i \text{ models noise}).$ 

ullet Non-linear measurements o **non-convex** maximum likelihood estimators.

## PhaseLift [4]

Phase retrieval can be solved as a convex matrix completion problem, following a combination of

- ightharpoonup semidefinite relaxation  $(\mathbf{x}^{\natural}\mathbf{x}^{\natural}^{H} = \mathbf{X}^{\natural})$
- ightharpoonup convex relaxation  $(rank o || \cdot ||_*)$

albeit in terms of the lifted variable  $\mathbf{X} \in \mathbb{C}^{p \times p}$ .

### Example: Phase retrieval - II

### Problem formulation

We solve the following PhaseLift variant:

$$f^* := \min_{\mathbf{X} \in \mathbb{C}^{p \times p}} \left\{ \frac{1}{2} \| \mathcal{A}(\mathbf{X}) - \mathbf{b} \|_2^2 : \| \mathbf{X} \|_* \le \kappa, \quad \mathbf{X} \ge 0 \right\}.$$
 (12)

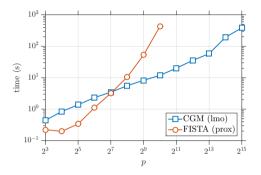
### Experimental setup [19]

Coded diffraction pattern measurements,  $\mathbf{b} = [\mathbf{b}_1, \dots, \mathbf{b}_L]$  with L = 20 different masks

$$\mathbf{b}_\ell = |\mathtt{fft}(\mathbf{d}_\ell^H \odot \mathbf{x}^{
atural})|^2$$

- $ightarrow \odot$  denotes Hadamard product;  $|\cdot|^2$  applies element-wise
- ightarrow  $\mathbf{d}_{\ell}$  are randomly generated octonary masks (distributions as proposed in [4])
- $\rightarrow$  Parametric choices:  $\lambda^0 = \mathbf{0}^n$ ;  $\epsilon = 10^{-2}$ ;  $\kappa = \text{mean}(\mathbf{b})$ .

### Example: Phase retrieval - III



## Test with synthetic data: Prox vs sharp

- $\rightarrow {\sf Synthetic\ data:\ } {\bf x}^{\natural} = {\tt randn}(p,1) + i \cdot {\tt randn}(p,1).$
- ightarrow Stopping criteria:  $\frac{\|\mathbf{x}^{\natural} \mathbf{x}^k\|_2}{\|\mathbf{x}^{\natural}\|_2} \leq 10^{-2}$ .
- $\rightarrow$  Averaged over 10 Monte-Carlo iterations.

Note that the problem is  $p \times p$  dimensional!

### A basic constrained non-convex problem

### Problem setting

$$f^{\star} := \min_{\mathbf{x} \in \mathbb{R}^p} \bigg\{ f(\mathbf{x}) : \mathbf{x} \in \mathcal{X} \bigg\},$$

#### **Assumptions**

- $\triangleright$   $\mathcal{X}$  is nonempty, convex, closed and bounded.
- ▶ f has *L*-Lipschitz continuous gradients, but it is **non-convex**.

### Stationary point

Due to constraints,  $\|\nabla f(\mathbf{x}^*)\| = 0$  may not hold!

Frank-Wolfe gap: Following measure, known as FW-gap, generalizes the definition of stationary point for constrained problems:

$$g_{FW}(\mathbf{x}) := \max_{\mathbf{y} \in \mathcal{X}} (\mathbf{x} - \mathbf{y})^T \nabla f(\mathbf{x})$$

- $ightharpoonup g_{FW}(\mathbf{x}) \geq 0$  for all  $\mathbf{x} \in \mathcal{X}$ .
- $\mathbf{x} \in \mathcal{X}$  is a stationary point if and only if  $g_{FW}(\mathbf{x}) = 0$ .

## CGM for non-convex problems

#### CGM for non-convex problems

- **1.** Choose  $\mathbf{x}^0 \in \mathcal{X}$ , K > 0 total number of iterations.
- **2.** For k = 0, 1, ..., K 1 perform:

$$\begin{cases} \hat{\mathbf{x}}^k &:= \text{Imo}_{\mathcal{X}}(\nabla f(\mathbf{x}^k)) \\ \mathbf{x}^{k+1} &:= (1 - \gamma_k) \mathbf{x}^k + \gamma_k \hat{\mathbf{x}}^k, \end{cases}$$

where  $\gamma_k := \frac{1}{\sqrt{K+1}}$ .

#### **Theorem**

Denote  $\bar{\mathbf{x}}$  chosen uniformly random from  $\{\mathbf{x}^1, \mathbf{x}^2, \dots, \mathbf{x}^K\}$ . Then, CGM satisfies

$$\min_{k=1,2,\ldots,K} g_{FW}(\mathbf{x}^k) \leq \mathbb{E}[g_{FW}(\bar{\mathbf{x}})] \leq \frac{1}{\sqrt{K}} \left( f(\mathbf{x}^0) - f^* + \frac{LD^2}{2} \right).$$

\* There exist stochastic CGM methods for non-convex problems. See [17] for details.

## A basic constrained stochastic problem

## Problem setting (Stochastic)

$$f^* := \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ \mathbb{E}[f(\mathbf{x}, \theta)] : \mathbf{x} \in \mathcal{X} \right\},\tag{13}$$

#### **Assumptions**

- ightharpoonup heta is a random vector whose probability distribution is supported on set  $\Theta$
- $\triangleright$   $\mathcal{X}$  is nonempty, convex, closed and bounded.
- $f(\cdot,\theta) \in \mathcal{F}_L^{1,1}(\mathbb{R}^p)$  for all  $\theta$  (i.e., convex with Lipschitz gradient).

## Example (Finite-sum model)

$$\mathbb{E}[f(\mathbf{x}, \theta)] = \frac{1}{n} \sum_{j=1}^{n} f_j(\mathbf{x})$$

- $i = \theta$  is a drawn uniformly from  $\Theta = \{1, 2, \dots, n\}$
- $f_j \in \mathcal{F}_L^{1,1}(\mathbb{R}^p)$  for all j (i.e., convex with Lipschitz gradient).

### Stochastic conditional gradient method

### Stochastic conditional gradient method (SFW)

- **1.** Choose  $\mathbf{x}^0 \in \mathcal{X}$ .
- **2.** For  $k = 0, 1, \ldots$  perform:

$$\begin{cases} \hat{\mathbf{x}}^k &:= \text{Imo}_{\mathcal{X}}(\tilde{\nabla}f(\mathbf{x}^k, \theta_k)) \\ \mathbf{x}^{k+1} &:= (1 - \gamma_k)\mathbf{x}^k + \gamma_k \hat{\mathbf{x}}^k, \end{cases}$$

where  $\gamma_k:=\frac{2}{k+2}$ , and  $\tilde{\nabla} f$  is an unbiased estimator of  $\nabla f$ .

## Theorem [9]

Assume that the following variance condition holds

$$\mathbb{E}\|\nabla f(\mathbf{x}^k) - \tilde{\nabla} f(\mathbf{x}^k, \theta_k)\|^2 \le \left(\frac{LD}{k+1}\right)^2. \tag{*}$$

Then, the iterates of SFW satisfies

$$\mathbb{E}[f(\mathbf{x}^k, \theta)] - f^* \le \frac{4LD^2}{k+1}.$$

 $(\star) \rightarrow SFW$  requires decreasing variance!



## Stochastic conditional gradient method

#### Stochastic conditional gradient method (SFW)

- 1. Choose  $\mathbf{x}^0 \in \mathcal{X}$ .
- **2.** For  $k = 0, 1, \ldots$  perform:

$$\begin{cases} \hat{\mathbf{x}}^k &:= \text{Imo}_{\mathcal{X}}(\tilde{\nabla}f(\mathbf{x}^k, \theta_k)) \\ \mathbf{x}^{k+1} &:= (1 - \gamma_k)\mathbf{x}^k + \gamma_k \hat{\mathbf{x}}^k, \end{cases}$$

where  $\gamma_k := \frac{2}{k+2}$ , and  $\tilde{\nabla} f$  is an unbiased estimator of  $\nabla f$ .

## Example (Finite-sum model)

$$\mathbb{E}[f(\mathbf{x}, \theta)] = \frac{1}{n} \sum_{j=1}^{n} f_j(\mathbf{x})$$

Assume  $f_j$  is G-Lipschitz continuous for all j. Suppose that  $\mathcal{S}_k$  is a random sampling (with replacement) from  $\Theta=\{1,2,\ldots,n\}$ . Then,

$$\tilde{\nabla} f(\mathbf{x}^k, \theta_k) := \frac{1}{|\mathcal{S}_k|} \sum_{j \in \mathcal{S}_k} f_j(\mathbf{x}^k) \quad \implies \quad \mathbb{E} \| \nabla f(\mathbf{x}) - \tilde{\nabla} f(\mathbf{x}, \theta_k) \|^2 \le \frac{G^2}{|\mathcal{S}_k|}.$$

Hence, by choosing  $|S_k| = (\frac{G(k+1)}{LD})^2$  we satisfy the variance condition for SFW.

## Wrap up!

 $\circ$  Monday: Transition from variance reduction to deep learning...

## \*Expanding on prox operator and optimality condition

#### Notes

- ▶ By definition,  $g(\mathbf{y}) + \frac{1}{2\lambda} \|\mathbf{y} \mathbf{x}\|^2$  attains its minimum when  $\mathbf{y} = \text{prox}_{\lambda q}(\mathbf{x})$ .
- ▶ One can see that  $g(y) + \frac{1}{2\lambda} ||y x||^2$  is convex, and prox operator computes its minimizer over  $\mathbb{R}^p$ .
- As a result, subdifferential of  $g(y) + \frac{1}{2\lambda} ||y x||^2$  at the minimizer  $(y = prox_{\lambda g}(x))$  should include 0.
- ▶ Hence,  $0 \in \partial g(\operatorname{prox}_{\lambda g}(\mathbf{x})) + \frac{1}{\lambda} \left(\operatorname{prox}_{\lambda g}(\mathbf{x}) \mathbf{x}\right)$ .
- ▶ After rearranging the above inclusion we obtain:  $\mathbf{x} \in \lambda \partial g(\operatorname{prox}_{\lambda g}(\mathbf{x})) + \operatorname{prox}_{\lambda g}(\mathbf{x})$
- We can rewrite the RHS as a single function:  $\lambda \partial g(\operatorname{prox}_{\lambda g}(\mathbf{x})) + \operatorname{prox}_{\lambda g}(\mathbf{x}) = (\lambda \partial g + \mathbb{I})(\operatorname{prox}_{\lambda g}(\mathbf{x}))$
- ► The inclusion becomes:  $\mathbf{x} \in (\lambda \partial g + \mathbb{I})(\text{prox}_{\lambda g}(\mathbf{x}))$ .
- Finally, we compute the inverse of  $(\lambda \partial g + \mathbb{I})(\cdot)$  to conclude:  $\operatorname{prox}_{\lambda g}(\mathbf{x}) = (\lambda \partial g + \mathbb{I})^{-1}(\mathbf{x})$ .
- o In the literature,  $(\lambda \partial g + \mathbb{I})^{-1}$  is called the resolvent of the subdifferential of g with parameter  $\lambda$ .
- o This is just a technical term that stands for proximal operator of  $\lambda g$ , as we have defined in this course.

## \*A short detour: Basic properties of prox-operator

## Property (Basic properties of prox-operator)

- 1.  $\operatorname{prox}_g(\mathbf{x})$  is well-defined and single-valued (i.e., the prox-operator (3) has a unique solution since  $g(\cdot) + \frac{1}{2} \|\cdot \mathbf{x}\|_2^2$  is strongly convex).
- 2. Optimality condition:

$$\mathbf{x} \in \operatorname{prox}_g(\mathbf{x}) + \partial g(\operatorname{prox}_g(\mathbf{x})), \ \mathbf{x} \in \mathbb{R}^p.$$

3.  $\mathbf{x}^*$  is a fixed point of  $\operatorname{prox}_q(\cdot)$ :

$$0 \in \partial g(\mathbf{x}^*) \Leftrightarrow \mathbf{x}^* = \operatorname{prox}_g(\mathbf{x}^*).$$

4. Nonexpansiveness:

$$\|\operatorname{prox}_g(\mathbf{x}) - \operatorname{prox}_g(\tilde{\mathbf{x}})\|_2 \le \|\mathbf{x} - \tilde{\mathbf{x}}\|_2, \ \forall \mathbf{x}, \tilde{\mathbf{x}} \in \mathbb{R}^p.$$

Note: An operator is called *non-expansive* if it is L-Lipschitz continuous with L=1.

## \*Adaptive Restart

# It is possible the preserve $\mathcal{O}(\frac{1}{k^2})$ convergence guarantee!

One needs to slightly modify the algorithm as below.

#### Generalized fast proximal-gradient scheme

- **1.** Choose  $\mathbf{x}^0 = \mathbf{x}^{-1} \in \text{dom}(F)$  arbitrarily as a starting point.
- **2.** Set  $\theta_0 = \theta_{-1} = 1$ ,  $\lambda := L_f^{-1}$
- **3.** For  $k=0,1,\ldots$ , generate two sequences  $\{\mathbf{x}^k\}_{k\geq 0}$  and  $\{\mathbf{y}^k\}_{k\geq 0}$  as:

$$\begin{cases} \mathbf{y}^{k} := \mathbf{x}^{k} + \theta_{k}(\theta_{k-1}^{-1} - 1)(\mathbf{x}^{k} - \mathbf{x}^{k-1}) \\ \mathbf{x}^{k+1} := \operatorname{prox}_{\lambda g} \left( \mathbf{y}^{k} - \lambda \nabla f(\mathbf{y}^{k}) \right), \\ \text{if restart test holds} \\ \theta_{k-1} = \theta_{k} = 1 \\ \mathbf{y}^{k} = \mathbf{x}^{k} \\ \mathbf{x}^{k+1} := \operatorname{prox}_{\lambda g} \left( \mathbf{y}^{k} - \lambda \nabla f(\mathbf{y}^{k}) \right) \end{cases}$$
(14)

## $\theta_k$ is chosen so that it satisfies

$$\theta_{k+1} = \frac{\sqrt{\theta_k^4 + 4\theta_k^2} - \theta_k^2}{2} < \frac{2}{k+3}$$

## \*Adaptive Restart: Guarantee

## Theorem (Global complexity [8])

The sequence  $\{\mathbf{x}^k\}_{k\geq 0}$  generated by the modified algorithm satisfies

$$F(\mathbf{x}^k) - F^* \le \frac{2L_f}{(k+2)^2} \left( R_0^2 + \sum_{k_i \le k} \left( \|\mathbf{x}^* - \mathbf{x}^{k_i}\|_2^2 - \|\mathbf{x}^* - \mathbf{z}^{k_i}\|_2^2 \right) \right) \ \forall k \ge 0.$$
 (15)

where  $R_0 := \min_{\mathbf{x}^{\star} \in \mathcal{S}^{\star}} \|\mathbf{x}^0 - \mathbf{x}^{\star}\|$ ,  $\mathbf{z}^k = \mathbf{x}^{k-1} + \theta_{k-1}^{-1}(\mathbf{x}^k - \mathbf{x}^{k-1})$  and  $k_i, i = 1...$  are the iterations for which the restart test holds.

# Various restarts tests that might coincide with $\|\mathbf{x}^* - \mathbf{x}^{k_i}\|_2^2 \leq \|\mathbf{x}^* - \mathbf{z}^{k_i}\|_2^2$

- Exact non-monotonicity test:  $F(\mathbf{x}^{k+1}) F(\mathbf{x}^k) > 0$
- Non-monotonicity test:  $\langle (L_F(\mathbf{y}^{k-1} \mathbf{x}^k), \mathbf{x}^{k+1} \frac{1}{2}(\mathbf{x}^k + y^{k-1}) \rangle > 0$  (implies exact non-monotonicity and it is advantageous when function evaluations are expensive)
- ▶ Gradient-mapping based test:  $\langle (L_f(\mathbf{y}^k \mathbf{x}^{k+1}), \mathbf{x}^{k+1} \mathbf{x}^k) > 0$

## \*Recall: Composite convex minimization

## Problem (Unconstrained composite convex minimization)

$$F^* := \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ F(\mathbf{x}) := f(\mathbf{x}) + g(\mathbf{x}) \right\}$$

- ightharpoonup f and g are both proper, closed, and convex.
- $ightharpoonup \operatorname{dom}(F) := \operatorname{dom}(f) \cap \operatorname{dom}(g) \neq \emptyset \text{ and } -\infty < F^{\star} < +\infty.$
- ▶ The solution set  $S^* := \{ \mathbf{x}^* \in \text{dom}(F) : F(\mathbf{x}^*) = F^* \}$  is nonempty.

## \*Recall: Composite convex minimization guarantees

Proximal gradient method(ISTA) vs. fast proximal gradient method (FISTA)

### Assumptions, step sizes and convergence rates

Proximal gradient method:

$$f \in \mathcal{F}_L^{1,1}, \quad \alpha = \frac{1}{L} \qquad \qquad F(\mathbf{x}^k) - F(\mathbf{x}^\star) \le \epsilon, \quad \mathcal{O}\Big(\frac{1}{\epsilon}\Big).$$

Fast proximal gradient method:

$$f \in \mathcal{F}_L^{1,1}, \quad \alpha = \frac{1}{L} \qquad \qquad F(\mathbf{x}^k) - F(\mathbf{x}^\star) \le \epsilon, \quad \mathcal{O}\bigg(\frac{1}{\sqrt{\epsilon}}\bigg).$$

### \*Recall: Composite convex minimization guarantees

Proximal gradient method(ISTA) vs. fast proximal gradient method (FISTA)

### Assumptions, step sizes and convergence rates

Proximal gradient method:

$$f \in \mathcal{F}_L^{1,1}, \quad \alpha = \frac{1}{L} \qquad \qquad F(\mathbf{x}^k) - F(\mathbf{x}^\star) \le \epsilon, \quad \mathcal{O}\Big(\frac{1}{\epsilon}\Big).$$

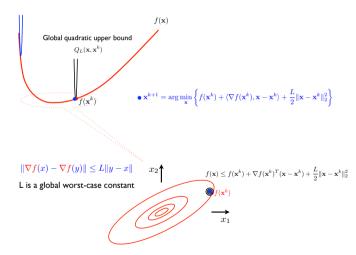
Fast proximal gradient method:

$$f \in \mathcal{F}_L^{1,1}, \quad \alpha = \frac{1}{L}$$
 
$$F(\mathbf{x}^k) - F(\mathbf{x}^\star) \le \epsilon, \quad \mathcal{O}\left(\frac{1}{\sqrt{\epsilon}}\right).$$

- $\circ$  We require  $\alpha_k$  to be a function of L.
- o It may not be possible to know exactly the Lipschitz constant. Line-search?
- $\circ$  Adaptation to local geometry  $\rightarrow$  may lead to larger steps.

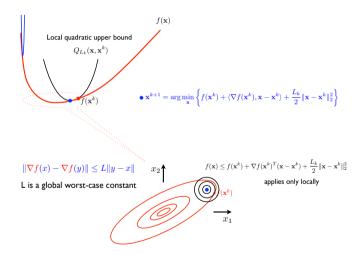
## \*How can we better adapt to the local geometry?

#### Non-adaptive:



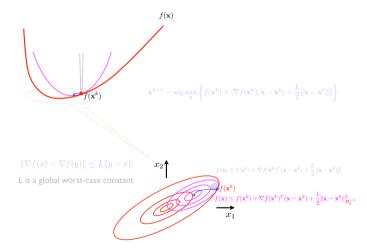
## \*How can we better adapt to the local geometry?

#### Line-search:



## \*How can we better adapt to the local geometry?

#### Variable metric:



# \*The idea of the proximal-Newton method

### Assumptions A.2

Assume that  $f \in \mathcal{F}^{2,1}_{L,\mu}(\mathbb{R}^p)$  and  $g \in \mathcal{F}_{\mathrm{prox}}(\mathbb{R}^p)$ .

### \*Proximal-Newton update

ightharpoonup Similar to classical newton, proximal-newton suggests the following update scheme using second order Taylor series expansion near  $\mathbf{x}_k$ .

$$\mathbf{x}^{k+1} := \arg\min_{\mathbf{x} \in \mathbb{R}^p} \left\{ \underbrace{\frac{1}{2} (\mathbf{x} - \mathbf{x}^k)^T \nabla^2 f(\mathbf{x}^k) (\mathbf{x} - \mathbf{x}^k) + \nabla f(\mathbf{x}^k)^T (\mathbf{x} - \mathbf{x}^k)}_{\text{2nd-order Taylor expansion near } \mathbf{x}^k} + g(\mathbf{x}) \right\}. \tag{16}$$

# \*The proximal-Newton-type algorithm

### Proximal-Newton algorithm (PNA)

- **1.** Given  $\mathbf{x}^0 \in \mathbb{R}^p$  as a starting point.
- **2.** For  $k = 0, 1, \dots$ , perform the following steps:
- 2.1. Evaluate an SDP matrix  $\mathbf{H}_k \approx \nabla^2 f(\mathbf{x}^k)$  and  $\nabla f(\mathbf{x}^k)$ .
- $\text{2.2. Compute } \mathbf{d}^k := \operatorname*{prox}_{\mathbf{H}_k^{-1}g} \bigg( \mathbf{x}^k \mathbf{H}_k^{-1} \nabla f(\mathbf{x}^k) \bigg) \mathbf{x}^k.$
- 2.3. Update  $\mathbf{x}^{k+1} := \mathbf{x}^k + \alpha_k \mathbf{d}^k$ .

# \*The proximal-Newton-type algorithm

#### Proximal-Newton algorithm (PNA)

- **1.** Given  $\mathbf{x}^0 \in \mathbb{R}^p$  as a starting point.
- **2.** For  $k = 0, 1, \dots$ , perform the following steps:
- 2.1. Evaluate an SDP matrix  $\mathbf{H}_k \approx \nabla^2 f(\mathbf{x}^k)$  and  $\nabla f(\mathbf{x}^k)$ .
- $\text{2.2. Compute } \mathbf{d}^k := \operatorname*{prox}_{\mathbf{H}_k^{-1}g} \bigg( \mathbf{x}^k \mathbf{H}_k^{-1} \nabla f(\mathbf{x}^k) \bigg) \mathbf{x}^k.$
- ${\color{red}\mathbf{2.3.}} \ \, \mathsf{Update} \,\, \mathbf{x}^{k+1} := \mathbf{x}^k + \alpha_k \mathbf{d}^k.$

### Remark

- $ightharpoonup \mathbf{H}_k \equiv 
  abla^2 f(\mathbf{x}^k) \Longrightarrow \mathsf{proximal-Newton algorithm}.$
- $lackbox{ iny }\mathbf{H}_kpprox 
  abla^2 f(\mathbf{x}^k) \Longrightarrow \mathsf{proximal-quasi-Newton}$  algorithm.
- ► A generalized prox-operator:  $\operatorname{prox}_{\mathbf{H}_k^{-1}g}\Big(\mathbf{x}^k + \mathbf{H}_k^{-1}\nabla f(\mathbf{x}^k)\Big).$

# \*Convergence analysis

# Theorem (Global convergence [11])

Assume generalized-prox subproblem is solved exactly for the algorithm and there exists  $\mu>0$  such that  $\mathbf{H}_k\succeq\mu\mathbb{I}$  for all  $k\geq0$ . Then;

 $\{\mathbf{x}^k\}_{k\geq 0}$  globally converges to a solution  $\mathbf{x}^{\star}$  of (2).

# \*Convergence analysis

# Theorem (Global convergence [11])

Assume generalized-prox subproblem is solved exactly for the algorithm and there exists  $\mu>0$  such that  $\mathbf{H}_k\succeq\mu\mathbb{I}$  for all  $k\geq0$ . Then;

$$\{\mathbf{x}^k\}_{k\geq 0}$$
 globally converges to a solution  $\mathbf{x}^*$  of (2).

# Theorem (Local convergence [11])

Assume generalized-prox subproblem is solved exactly for the algorithm there exists  $0 < \mu \le L_2 < +\infty$  such that  $\mu \mathbb{I} \prec \mathbf{H}_k \prec L_2 \mathbb{I}$  for all sufficiently large k. Then;

- If  $\mathbf{H}_k \equiv \nabla^2 f(\mathbf{x}^k)$ , then  $\alpha_k = 1$  for k sufficiently large (full-step).
- ▶ If  $\mathbf{H}_k \equiv \nabla^2 f(\mathbf{x}^k)$ , then  $\{\mathbf{x}^k\}$  locally converges to  $\mathbf{x}^*$  at a quadratic rate.
- ► If **H**<sub>k</sub> satisfies the Dennis-Moré condition:

$$\lim_{k \to +\infty} \frac{\|(\mathbf{H}_k - \nabla^2 f(\mathbf{x}^*))(\mathbf{x}^{k+1} - \mathbf{x}^k)\|}{\|\mathbf{x}^{k+1} - \mathbf{x}^k\|} = 0,$$
(17)

then  $\{x^k\}$  locally converges to  $x^*$  at a super linear rate.



\*How to compute the approximation  $H_k$ ?

# How to update $\mathbf{H}_k$ ?

Matrix  $\mathbf{H}_k$  can be updated by using low-rank updates.

**BFGS** update: maintain the Dennis-Moré condition and  $\mathbf{H}_k \succ 0$ .

$$\mathbf{H}_{k+1} := \mathbf{H}_k + \frac{\mathbf{y}_k \mathbf{y}_k^T}{\mathbf{s}_k^T \mathbf{y}_k} - \frac{\mathbf{H}_k \mathbf{s}_k \mathbf{s}_k^T \mathbf{H}_k}{\mathbf{s}_k^T \mathbf{H}_k \mathbf{s}_k}, \quad \mathbf{H}_0 := \gamma \mathbb{I}, \ (\gamma > 0).$$

where  $\mathbf{y}_k := \nabla f(\mathbf{x}^{k+1}) - \nabla f(\mathbf{x}^k)$  and  $\mathbf{s}_k := \mathbf{x}^{k+1} - \mathbf{x}^k$ .

Diagonal+Rank-1 [3]: computing PN direction d<sup>k</sup> is in polynomial time, but it does not maintain the Dennis-Moré condition:

$$\mathbf{H}_k := \mathbf{D}_k + \mathbf{u}_k \mathbf{u}_k^T, \ \mathbf{u}_k := \frac{\mathbf{s}_k - \mathbf{H}_0 \mathbf{y}_k}{\sqrt{(\mathbf{s}_k - \mathbf{H}_0 \mathbf{y}_k)^T \mathbf{y}_k}},$$

where  $\mathbf{D}_k$  is a positive diagonal matrix.

#### \*Pros and cons

#### Pros

- ► Fast local convergence rate (super-linear or quadratic)
- ▶ Numerical robustness under the inexactness/noise ([11]).
- ▶ Well-suited for problems with many data points but few parameters. For example,

$$F^* := \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ \sum_{j=1}^n \ell_j(\mathbf{a}_j^T \mathbf{x} + b_j) + g(\mathbf{x}) \right\},$$

where  $\ell_j$  is twice continuously differentiable and convex,  $g \in \mathcal{F}_{prox}$ ,  $p \ll n$ .

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#### Cons

- Expensive iteration compared to proximal-gradient methods.
- ▶ Global convergence rate may be worse than accelerated proximal-gradient methods.
- Requires a good initial point to get fast local convergence.
- ▶ Requires strict conditions for global/local convergence analysis.

## \*Example 1: Sparse logistic regression

# Problem (Sparse logistic regression)

Given a sample vector  $\mathbf{a} \in \mathbb{R}^p$  and a binary class label vector  $\mathbf{b} \in \{-1, +1\}^n$ . The conditional probability of a label b given  $\mathbf{a}$  is defined as:

$$\mathbb{P}(b|\mathbf{a}, \mathbf{x}, \mu) = \frac{1}{1 + e^{-b(\mathbf{x}^T \mathbf{a} + \mu)}},$$

where  $\mathbf{x} \in \mathbb{R}^p$  is a weight vector,  $\mu$  is called the intercept.

Goal: Find a sparse-weight vector x via the maximum likelihood principle.

### **Optimization formulation**

$$\min_{\mathbf{x} \in \mathbb{R}^p} \left\{ \underbrace{\frac{1}{n} \sum_{i=1}^n L(b_i(\mathbf{a}_i^T \mathbf{x} + \mu))}_{f(\mathbf{x})} + \underbrace{\rho \|\mathbf{x}\|_1}_{g(\mathbf{x})} \right\}, \tag{18}$$

where  $\mathbf{a}_i$  is the *i*-th row of data matrix  $\mathbf{A}$  in  $\mathbb{R}^{n\times p}$ ,  $\rho>0$  is a regularization parameter, and  $\ell$  is the logistic loss function  $\ell(\tau):=\log(1+e^{-\tau})$ .

\*Example: Sparse logistic regression

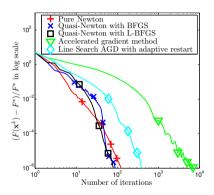
#### Real data

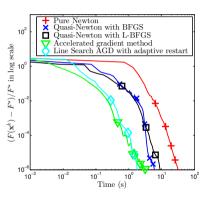
- ▶ Real data: w2a with n = 3470 data points, p = 300 features
- ► Available at http://www.csie.ntu.edu.tw/~cjlin/libsvmtools/datasets/binary.html.

#### **Parameters**

- ▶ Tolerance  $10^{-6}$ .
- ▶ L-BFGS memory m = 50.
- ▶ Ground truth: Get a high accuracy approximation of  $x^*$  and  $f^*$  by TFOCS with tolerance  $10^{-12}$ .

## \*Example: Sparse logistic regression-Numerical results





# \*Example 2: $\ell_1$ -regularized least squares

# Problem ( $\ell_1$ -regularized least squares)

Given  $\mathbf{A} \in \mathbb{R}^{n \times p}$  and  $\mathbf{b} \in \mathbb{R}^n$ , solve:

$$F^{\star} := \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ F(\mathbf{x}) := \frac{1}{2} \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2 + \rho \|\mathbf{x}\|_1 \right\},$$

where  $\rho > 0$  is a regularization parameter.

# Complexity per iterations

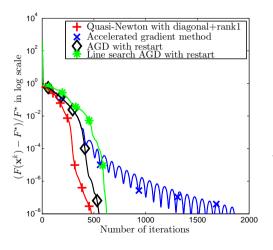
- ▶ Evaluating  $\nabla f(\mathbf{x}^k) = \mathbf{A}^T(\mathbf{A}\mathbf{x}^k \mathbf{b})$  requires one  $\mathbf{A}\mathbf{x}$  and one  $\mathbf{A}^T\mathbf{y}$ .
- ▶ One soft-thresholding operator  $\operatorname{prox}_{\lambda g}(\mathbf{x}) = \operatorname{sign}(\mathbf{x}) \otimes \max\{|\mathbf{x}| \rho, 0\}.$
- ▶ Optional: Evaluating  $L = \|\mathbf{A}^T \mathbf{A}\|$  (spectral norm) via power iterations (e.g., 20 iterations, each iteration requires one  $\mathbf{A}\mathbf{x}$  and one  $\mathbf{A}^T \mathbf{y}$ ).

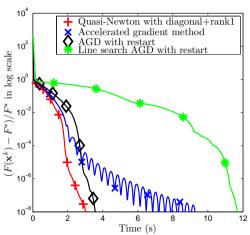
# Synthetic data generation

- $ightharpoonup \mathbf{A} := \operatorname{randn}(n,p)$  standard Gaussian  $\mathcal{N}(0,\mathbb{I})$ .
- $ightharpoonup \mathbf{x}^*$  is a s-sparse vector generated randomly.
- $\mathbf{b} := \mathbf{A} \mathbf{x}^* + \mathcal{N}(0, 10^{-3}).$

# \*Example 2: $\ell_1$ -regularized least squares - Numerical results - Trial 1

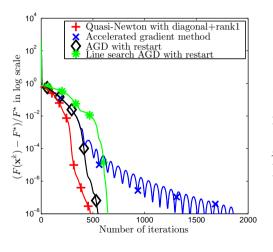
Parameters:  $n = 750, p = 2000, s = 200, \rho = 1$ 

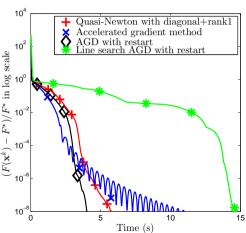




# \*Example 2: $\ell_1$ -regularized least squares - Numerical results - Trial 2

Parameters:  $n = 750, p = 2000, s = 200, \rho = 1$ 





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